# Proposed experiment for the quantum "Guess My Number" protocol 

Adán Cabello* and Antonio J. López-Tarrida ${ }^{\dagger}$<br>Departamento de Física Aplicada II, Universidad de Sevilla, 41012 Sevilla, Spain<br>(Received 21 April 2004; published 4 February 2005)


#### Abstract

An experimental realization of the entanglement-assisted "Guess My Number" protocol for the reduction of communication complexity, introduced by Steane and van Dam, would require producing and detecting threequbit GHZ states with an efficiency $\eta>0.70$, which would require single photon detectors of efficiency $\sigma$ $>0.89$. We propose a modification of the protocol which can be translated into a real experiment using present-day technology. In the proposed experiment, the quantum reduction of the multiparty communication complexity would require an efficiency $\eta>0.05$, achievable with detectors of $\sigma>0.47$, for four parties, and $\eta>0.17(\sigma>0.55)$ for three parties.


DOI: 10.1103/PhysRevA.71.020301
PACS number(s): 03.67.Hk, 02.50.Le, 03.65.Ud, 03.65.Ta

One of the most impressive applications of quantum resources for information processing is the reduction of the communication complexity required for certain computations [1-5]. Let us suppose that two or more separated parties need to compute a function of a number of inputs distributed among them. Using the best classical strategy, this would require a certain minimum amount of classical communication to be transmitted between the parties. However, if the parties initially shared some entangled states, then the amount of classical communication required for the computation would be a great deal smaller than if no entanglement were present. The quantum advantage usually grows with the number of parties involved [2]. Entanglement-assisted reduction of classical communication complexity has numerous potential applications in computer networks, VLSI circuits, and data structures [6].

A particularly attractive, thought-provoking, and stimulating way to show the quantum advantage was proposed by Steane and van Dam as a method for always winning the television contest "Guess My Number" (GMN) [5]. A team of three contestants (Alice, Bob, and Charlie), each of them isolated in a booth, is given an integer number $n=n_{A}+n_{B}$ $+n_{C}$ of apples (where $n_{j}=0,1 / 2,1$, or $3 / 2$ ). One of the contestants must guess whether the number is odd or even just by receiving one bit from the other two contestants. The best classical strategy would allow the contestants to win in $75 \%$ of the cases. However, they can win in $100 \%$ of the cases if they initially share three-qubit Greenberger-HorneZeilinger (GHZ) states [7,8]. The same game can be played with four contestants and the quantum versus classical advantage is the same: $100 \%$ vs $75 \%$. Steane and van Dam stressed that "A laboratory demonstration of entanglementenhanced communication would be (...) a landmark in quantum physics and quantum information science" [5]. So far, however, the requirements for an experimental implementation of the quantum GMN protocol have impeded further progress. Some progress has been reported on simpler schemes of quantum reduction of classical communication

[^0]complexity. For instance, Xue et al. presented an experiment on quantum reduction of two-party communication complexity based on two-qubit entanglement [9]. Galvão proposed a protocol requiring only one qubit and a detection efficiency $\sigma>0.33$ [10]. More recently, Brukner, Żukowski, and Zeilinger have introduced a quantum reduction of two-party communication complexity based on the entanglement between two qutrits [11].

The main obstacle for an experimental realization of the quantum GMN protocol is the high detection efficiency required. The required setup would consist of a source of GHZ states, single qubit operations, and single qubit detectors. If we define the overall efficiency $\eta$ as the number of threequbit (or four-qubit) joint detections corresponding to GHZ states, divided by the number of three-qubit (or four-qubit) systems emitted by the source, then, assuming that when no joint detection occurs the probability of winning the game is only $1 / 2$, the experimental probability of winning the GMN game using GHZ states is

$$
\begin{equation*}
P_{\exp }(\eta)=\eta+(1-\eta) \frac{1}{2} \tag{1}
\end{equation*}
$$

Therefore, the quantum advantage could be detected if an overall efficiency $\eta>0.50$ could be achieved. In the threequbit case, the experiment would require threefold coincidences between detectors so that each individual detector should have an efficiency $\sigma=0.79$ (since $\sigma=\eta^{1 / c}, c$ being the number of qubits). Moreover, in order to obtain an experimental quantum probability of winning the GMN game $10 \%$ higher than the best classical probability, we would need $\eta$ $>0.70$, which would require detectors of efficiency $\sigma$ $>0.89$.

Quantum optics provides the best way to produce qubits in a GHZ state and distribute them to various spacetime regions. However, the first experiments producing threephoton polarization-entangled GHZ states [12,13] did not satisfy the demands of the GMN protocol, because only a tiny fraction of the ensemble of photon triplets was detected [5]. Further experiments producing four-photon GHZ states [14] yield a fourfold coincidence with a success probability 4 times higher than that of previous three-photon experiments. Moreover, recent experiments [15] report a fourfold coinci-
dence rate 2 orders of magnitude brighter than in [14]. We shall show that in the very near future this technology could allow an experimental demonstration of a quantum reduction of a genuine three or four-party communication complexity. In this paper, we introduce a modified version of the quantum GMN protocol which is experimentally feasible with current technology. We shall describe a quantum reduction of three-party (four-party) communication complexity in which the quantum advantage is clear, provided we can produce three (four) qubits in a GHZ state and detect them all separately with an overall efficiency $\eta>0.17$, which would require detectors of efficiency $\sigma>0.55 \quad(\eta>0.05$ and $\sigma$ $>0.47$, for four qubits). The main goal of this proposal is to note that the absence of perfect sources and detectors does not prevent us from performing an experimental demonstration of a quantum reduction of a genuine multiparty communication complexity and to stimulate experimental work along these lines.

The modified GMN game preserves all the essential features of the original game, but includes rules that relax the detection requirements to experimentally show the quantum advantage. The modified GMN game features one referee (and a fourth contestant in the $c=4$ version). We shall discuss in detail the four-party version of the modified protocol; similar rules apply to the three-party version. During the game, each of four contestants (Alice, Bob, Charlie, and David) is isolated in a booth. Before the game starts, they can take anything they want with them into the booths, but once they are in, they will not be able to communicate with each other or with anybody else, save for the referee. Once they are in the booths, the referee distributes among them a randomly chosen integer number $n$ of apples in four portions, $n=n_{A}+n_{B}+n_{C}+n_{D}$, such that $n_{j}=0,1 / 2,1$, or $3 / 2$. Then, the referee asks each and everyone whether or not they are ready to play the game; if all contestants say yes, then the referee asks Bob, Charlie, and David to give him a bit. Then, the referee adds (modulo 2) the three bits, and hands the result over to Alice. The team wins if Alice ascertains whether the total number of distributed apples is even or odd. If any contestant refuses to play the game, then the referee distributes a new number $n^{\prime}=n_{A}^{\prime}+n_{B}^{\prime}+n_{C}^{\prime}+n_{D}^{\prime}$ of apples and asks the four contestants again whether or not they are ready to play the game, etc. If the referee distributes $N$ rounds of apples, then the contestants are forced to play the game for at least $r$ rounds (hereafter referred as "the played rounds"). The contestants know $p=r / N$ before the game starts. In addition, the referee must ensure that each of the 128 possible variations of apples (see Table I) occurs with the same frequency in the played rounds.

In the modified GMN game, if the referee forces the contestants to play in $r=p N$ of the $N$ rounds, the contestants can refuse to play between the first and the $N-r$ round, but then they are forced to play in the remaining $r$ rounds. If they decide to play without being forced to do so then, every time they play, they will postpone in one round the moment they have to play compulsorily. The maximum classical probability of winning is obtained by combining two strategies. The first one applies in the rounds in which they play without being forced to do so, and can be designed in a way such that the contestants know when they must play and success is

TABLE I. The 19 integer combinations of $0,1 / 2,1$, and $3 / 2$, and their corresponding 128 variations. In 64 of them $n_{i}+n_{j}+n_{k}$ $+n_{l}$ is an odd number while in the other 64 it is an even number.

| $\left\{n_{i}, n_{j}, n_{k}, n_{l}\right\}$ | $n_{i}+n_{j}+n_{k}+n_{l}$ | Number of variations |
| :---: | :---: | :---: |
| $\{0,0,0,0\}$ | 0 | 1 |
| $\{0,0,0,1\}$ | 1 | 4 |
| $\{0,0,1 / 2,1 / 2\}$ | 1 | 6 |
| $\{0,0,1 / 2,3 / 2\}$ | 2 | 12 |
| $\{0,0,1,1\}$ | 2 | 6 |
| $\{0,1 / 2,1 / 2,1\}$ | 2 | 12 |
| $\{1 / 2,1 / 2,1 / 2,1 / 2\}$ | 2 | 1 |
| $\{0,0,3 / 2,3 / 2\}$ | 3 | 6 |
| $\{0,1 / 2,1,3 / 2\}$ | 3 | 24 |
| $\{0,1,1,1\}$ | 3 | 4 |
| $\{1 / 2,1 / 2,1 / 2,3 / 2\}$ | 3 | 4 |
| $\{1 / 2,1 / 2,1,1\}$ | 3 | 6 |
| $\{0,1,3 / 2,3 / 2\}$ | 4 | 12 |
| $\{1 / 2,1 / 2,3 / 2,3 / 2\}$ | 4 | 6 |
| $\{1 / 2,1,1,3 / 2\}$ | 4 | 12 |
| $\{1,1,1,1\}$ | 4 | 1 |
| $\{1 / 2,3 / 2,3 / 2,3 / 2\}$ | 5 | 4 |
| $\{1,1,3 / 2,3 / 2\}$ | 5 | 6 |
| $\{3 / 2,3 / 2,3 / 2,3 / 2\}$ | 6 | 1 |

guaranteed when they do play (this happens, at best, once in every 32 rounds, on average, if $c=4$, and once in every 8 rounds, on average, if $c=3$ ). The second strategy applies when they are forced to play. It could be any of the best classical strategies of the original GMN game, giving a probability of success of 3/4 (for instance, each contestant would give the referee a bit value 0 if she/he had received $n_{j}=0$ or $1 / 2$, or a bit value 1 if she/he had received $n_{j}=1$ or $3 / 2$ ). From all this follows that, for the modified GMN game, the best classical strategies (of which there are several) give the following maximum probability of winning for $c=3$ or $c=4$ contestants being forced to play in at least $p$ of the rounds,

$$
\begin{align*}
P_{C}(c, p)= & \lim _{N \rightarrow \infty}\left[\frac{(1-\mu)^{N-p N}}{4 p N} \sum_{j=0}^{p N}(j+3 p N) \mu^{j}\right. \\
& \left.\times\binom{ N-p N+j-1}{j}+\sum_{j=p N+1}^{N}(1-\mu)^{N-j} \mu^{j}\binom{N}{j}\right], \tag{2}
\end{align*}
$$

where

$$
\begin{equation*}
\mu=\frac{8}{2^{2 c}} . \tag{3}
\end{equation*}
$$

For $c=3$ and $c=4$, this probability is represented as a function of $p$ in Fig. 1. In addition, Fig. 1 contains numerical simulations of the probability that the team with $c=3$ and $c$ $=4$ wins when using the best classical strategy for games of $N=100$ rounds.


FIG. 1. Exact and numerical simulations of the probability of the contestants winning the modified GMN game using the best classical strategy, as a function of the minimum percentage of rounds the referee forces them to play, for the three-contestant game (black squares) and four-contestant game (white squares). In the numerical simulations the referee distributes $N=100$ rounds. The exact probabilities are given by Eq. (2). Interestingly, for $c=3$ contestants forced to play in at least $p=0.17$ of the rounds, the best classical probability of winning is only $P_{C}=0.92$. For $c$ $=4$ contestants forced to play in at least $p=0.05$ of the rounds, the best classical probability of winning is only $P_{C}=0.90$.

Note that, in both cases, if the referee forces the team to play all the rounds, the probability of winning by using the best classical strategy is $3 / 4$ while, if the referee forces them to play in at least one of every 100 rounds, then the probability of success using the best classical strategy is approximately 1 .

Let us now see what the probabilities of winning are when using the best entanglement-assisted strategy. The contestants will always win if they use the following method.
(1) Each contestant carries a qubit belonging to a fourqubit system initially prepared in the GHZ state

$$
\begin{equation*}
|\mathrm{GHZ}\rangle=\frac{1}{\sqrt{2}}(|\overline{0} \overline{0} \overline{0} \overline{0}\rangle+|\overline{1} \overline{1} \overline{1} \overline{1}\rangle) \tag{4}
\end{equation*}
$$

where $\quad|\overline{0} \overline{0} \overline{0} \overline{0}\rangle=|\overline{0}\rangle \otimes|\overline{0}\rangle \otimes|\overline{0}\rangle \otimes|\overline{0}\rangle$, where $\quad|\overline{0}\rangle=(1 / \sqrt{2})(|0\rangle$ $+|1\rangle)$ and $|\overline{1}\rangle=(1 / \sqrt{2})(|0\rangle-|1\rangle)$.
(2) Each contestant $j$ applies to her/his qubit the rotation

$$
\begin{equation*}
R\left(n_{j}\right)=|\overline{0}\rangle\langle\overline{0}|+e^{i n_{j} \pi}|\overline{1}\rangle\langle\overline{1}|, \tag{5}
\end{equation*}
$$

where $n_{j}$ is her/his number of apples.
(3) Then, each contestant measures her/his qubit in the computational basis $\{|0\rangle,|1\rangle\}$.
(4) If, due to the inefficiency of the detectors, a contestant does not obtain a result, then she/he will tell the referee that she/he will not play the game, and the referee will therefore abort that round. Note that, in the aborted rounds, Alice does not receive any bits from the referee. If all contestants consent to play that round, then Bob, Charlie, and David will give their outcomes to the referee, who will add them up, and give the result to Alice.

In this case Alice can give the correct answer with probability 1 because state (4) has the following property: for any $n_{A}+n_{B}+n_{C}+n_{D}$ integer (where $n_{j}=0,1 / 2,1$, or $3 / 2$ ),

$$
\begin{align*}
& R\left(n_{A}\right) \otimes R\left(n_{B}\right) \otimes R\left(n_{C}\right) \otimes R\left(n_{D}\right)|\mathrm{GHZ}\rangle \\
& \quad= \begin{cases}|\mathrm{GHZ}\rangle & \text { if } n_{A}+n_{B}+n_{C}+n_{D} \text { is even }, \\
\left|\mathrm{GHZ}^{\perp}\right\rangle & \text { if } n_{A}+n_{B}+n_{C}+n_{D} \text { is odd, }\end{cases} \tag{6}
\end{align*}
$$

where $|\mathrm{GHZ}\rangle$ and $\left|\mathrm{GHZ}^{\perp}\right\rangle$ can be reliably distinguished by local measurements in the computational basis:

$$
\begin{align*}
|\mathrm{GHZ}\rangle= & \frac{1}{2 \sqrt{2}}(|0000\rangle+|0011\rangle+|0101\rangle+|0110\rangle+|1001\rangle \\
& +|1010\rangle+|1100\rangle+|1111\rangle)  \tag{7}\\
\left|\mathrm{GHZ}^{\perp}\right\rangle= & \frac{1}{2 \sqrt{2}}(|0001\rangle+|0010\rangle+|0100\rangle+|0111\rangle+|1000\rangle \\
& +|1011\rangle+|1101\rangle+|1110\rangle) \tag{8}
\end{align*}
$$

Assuming that, when all four contestants obtain a result, this corresponds to a GHZ state (i.e., assuming that any error in the preparation is negligible), then having an experimental efficiency $\eta$ allows the team to play the modified GMN game with $p=\eta$. Now let us go back to the probabilities illustrated in Fig. 1. In the first place we shall compare the experimental requirements for the original GMN game with three qubits to those of the modified protocol. The most important point is that, while in the original GMN protocol the difference between the quantum and classical probabilities of winning could be detected only if the experimental setup has an overall efficiency $\eta>0.50$ (that is, a single qubit detection efficiency $\sigma=0.79$ ), in the modified protocol the difference between the quantum and classical probabilities can be detected for almost any efficiency. Moreover, as seen above, to obtain an experimental quantum probability of winning $10 \%$ higher than the best classical probability in the original GMN protocol, the setup would need to have $\eta>0.70$ (that is, a single qubit detection efficiency $\sigma>0.89$ ). However, to obtain a difference between the quantum and classical probabilities of winning higher than $7.7 \%$ in the modified proto-

TABLE II. Examples of single photon detection efficiency requirements for the modified GMN protocol. $c$ is the number of parties, $P_{Q}-P_{C}$ is the difference between the quantum and classical probabilities of winning, $\eta$ is the number of joint detections divided by the number of systems emitted by the source, and $\sigma$ is the corresponding single photon detection efficiency.

| $c$ | $P_{Q}-P_{C}$ | $\eta$ | $\sigma$ |
| :---: | :---: | :---: | :---: |
| 3 | 0.250 | 1 | 1 |
| 4 | 0.250 | 1 | 1 |
| 3 | $>0.214$ | $>0.50$ | $>0.79$ |
| 4 | $>0.218$ | $>0.20$ | $>0.67$ |
| 3 | $>0.107$ | $>0.20$ | $>0.58$ |
| 4 | $>0.177$ | $>0.10$ | $>0.56$ |
| 3 | $>0.077$ | $>0.17$ | $>0.55$ |
| 4 | $>0.097$ | $>0.05$ | $>0.47$ |

col, the setup would only require $\eta>0.17$ (that is, detectors of efficiency $\sigma>0.55$ ). On the other hand, since sources of four-photon GHZ states (4) are currently available [14,15], then it is interesting to note that, for $c=4$ contestants and an experimental setup with an overall efficiency $\eta>0.05$ (that is, with detectors of efficiency $\sigma>0.47$ ), it would be possible to obtain a difference between the quantum and classical probabilities higher than $9.7 \%$. Photodetectors of $\sigma>0.47$ are currently available. Other examples for different values of $\sigma$ can be found in Table II. An interesting advantage of all these experiments is that the expected quantum probabilities are 1 , which implies that the error of the experimental results, given by the standard deviation $\sqrt{P(1-P) / r}$, where $r$ is the number of coincidences (i.e., played rounds), should be very low.

The proposed experiment would consist of a source emit-
ting three (or four) polarization-entangled photons in a GHZ state generated in a parametric-down conversion process [14,15], coupled into three (four) single mode optical fibers which distribute the photons to different regions, where each photon suffers a randomly chosen rotation of the type (5), and a linear polarization measurement (typically the horizontal and vertical states represent the computational basis). If all photons are detected, then two (three) of the contestants send their result to the referee who adds them up and sends the result to the third (fourth) contestant, who adds it to her result and gives the answer.

To sum up, while testing the advantage of the original quantum GMN protocol involving three parties would require detectors of an efficiency at least $\sigma>0.79$ (or $\sigma$ $>0.89$ to obtain an experimental quantum probability of winning $10 \%$ higher than the best classical probability), we have introduced a modified quantum GMN protocol involving three or four parties and preserving all the essential features of the original one, but with the remarkable property that the quantum vs classical advantage is detectable for any $\sigma$. To be specific, $\sigma>0.55$ would allow us to obtain an experimental quantum probability of winning at least $7.7 \%$ higher than the best classical probability in the three-party case, and $\sigma>0.47$ would allow us to obtain an experimental quantum probability of winning at least $9.7 \%$ higher than the best classical probability in the four-party case. Our hope is that this proposal will stimulate experimental work to detect the quantum reduction of a genuine multiparty communication complexity.

The authors thank M. Bourennane, E. F. Galvão, C. Serra, and H . Weinfurter for useful discussions and comments. This work was supported by the Spanish Ministerio de Ciencia y Tecnología Project BFM2002-02815 and the Junta de Andalucía Project FQM-239.
[1] R. Cleve and H. Buhrman, Phys. Rev. A 56, 1201 (1997).
[2] H. Buhrman, R. Cleve, and A. Wigderson, in Proceedings of the 30th Annual ACM Symposium on the Theory of Computing (ACM Press, New York, 1998), p. 63.
[3] H. Buhrman, W. van Dam, P. Høyer, and A. Tapp, Phys. Rev. A 60, 2737 (1999).
[4] R. Raz, in Proceedings of the 31st Annual ACM Symposium on the Theory of Computing (ACM Press, New York, 1999), p. 358.
[5] A. M. Steane and W. van Dam, Phys. Today 53(2), 35 (2000).
[6] E. Kushilevitz and N. Nisan, Communication Complexity (Cambridge University Press, Cambridge, England, 1997).
[7] D. M. Greenberger, M. A. Horne, and A. Zeilinger, in Bell's Theorem, Quantum Theory, and Conceptions of the Universe, edited by M. Kafatos (Kluwer Academic, Dordrecht, 1989), p. 69 .
[8] D. M. Greenberger, M. A. Horne, A. Shimony, and A. Zeilinger, Am. J. Phys. 58, 1131 (1990).
[9] P. Xue, Y.-F. Huang, Y.-S. Zhang, C.-F. Li, and G.-C. Guo, Phys. Rev. A 64, 032304 (2001).
[10] E. F. Galvão, Phys. Rev. A 65, 012318 (2002).
[11] Č. Brukner, M. Żukowski, and A. Zeilinger, Phys. Rev. Lett. 89, 197901 (2002).
[12] D. Bouwmeester, J.-W. Pan, M. Daniell, H. Weinfurter, and A. Zeilinger, Phys. Rev. Lett. 82, 1345 (1999).
[13] J.-W. Pan, D. Bouwmeester, M. Daniell, H. Weinfurter, and A. Zeilinger, Nature (London) 403, 515 (2000).
[14] J.-W. Pan, M. Daniell, S. Gasparoni, G. Weihs, and A. Zeilinger, Phys. Rev. Lett. 86, 4435 (2001).
[15] Z. Zhao, T. Yang, Y.-A. Chen, A.-N. Zhang, M. Żukowski, and J.-W. Pan, Phys. Rev. Lett. 91, 180401 (2003).


[^0]:    *Electronic address: adan@us.es
    ${ }^{\dagger}$ Electronic address: tarrida@us.es

