

CLASSIFICATION OF 3-SEMINETS WITH AT MOST 5 POINTS.

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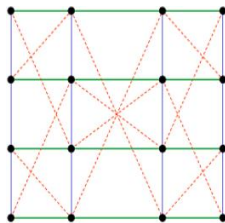
VIII JMDA, Almería 2012

NETS.

$\mathfrak{P} \equiv$ Finite set of points.

$\mathcal{L} \equiv$ Subsets (*lines*) of \mathfrak{P} s.t. \exists partition of \mathcal{L} into $k \geq 3$ *parallel* classes:

$L_1, \dots, L_k.$



$(\mathfrak{P}, L_1, \dots, L_k)$ is a **k -net** [Bruck, 1963] if:

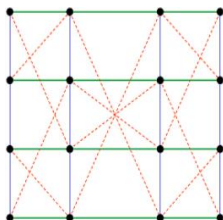
- ▶ Any two lines of different classes intersect in **exactly one** point.
- ▶ Every point belongs to **exactly one** line of each class.

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\equiv

1	2	3	4
2	4	1	3
3	1	4	2
4	3	2	1

- ▶ All parallel classes have the same number of lines.
- ▶ Every line contains the same number of points (**order**).
- ▶ Two k -nets are *isomorphic* if there exists a bijection between their sets of points which preserves collinearity in each parallel class.
- ▶ Any 3-net of order n is uniquely identified with a *Latin square* of the same order.

CLASSIFICATION OF LATIN SQUARES.

- ▶ **Orthogonal representation:** $O(P) = \{(\text{row}, \text{column}, \text{symbol})\}$.
- ▶ **Classification:**
 1. **Isotopism:** Permutations of rows, columns and symbols.
 2. **Parastrophism:**

$$\pi \in S_3 \rightarrow O(P^\pi) = \{(I_{\pi(1)}, I_{\pi(2)}, I_{\pi(3)}) \mid (I_1, I_2, I_3) \in O(P)\}.$$
 3. **Paratopism (main classes):** Composition of isotopism and parastrophism.
- ▶ Isotopic LS \equiv Isomorphic 3-nets.
- ▶ Paratopic LS \equiv Isomorphic 3-nets after relabeling parallel classes.

n	$ LS_n $	IC	MC
1	1	1	1
2	2	1	1
3	12	1	1
4	576	2	2
5	161280	2	2
6	812851200	22	12
7	61479419904000	564	147
8	108776032459082956800	1676267	283657
9	5524751496156892842531225600	115618721533	19270853541
10	9982437658213039871725064756920320000	208904371354363006	34817397894749939
11	776966836171770144107444346734230682311065600000	12216177315369229261482540	2036029552582883134196099

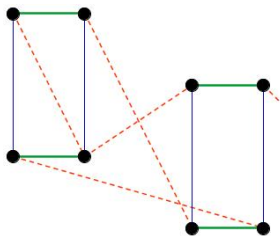
[McKay and Wanless, 2005; Hulpke et al., 2011]

SEMINETS.

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$(\mathfrak{P}, L_1, \dots, L_k)$ is a ***k*-semi-net** [Ušan, 1977] if:

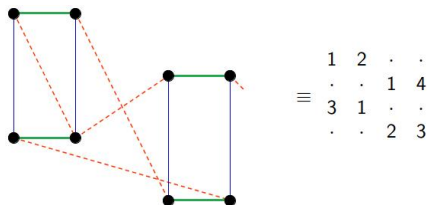
- ▶ Any two lines of different classes intersect in **at most one** point.
- ▶ Every point belongs to **exactly one** line of each class.

SEMINETS.

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$L_1, \dots, L_k.$



- ▶ Parallel classes can have different number of lines: r, s and n .
- ▶ Lines can contain different number of points.
- ▶ It can contain **skew** lines: Non-parallel lines without common points.
- ▶ Any 3-semi-net with parallel classes of r, s and n lines can be uniquely identified with an $r \times s$ **partial Latin rectangle** based on n symbols.

PARTIAL LATIN RECTANGLES.

- ▶ An $r \times s$ partial Latin rectangle based on a set of n symbols is an $r \times s$ array in which each cell is either empty or contains one element chosen from a set of symbols, $[n] = \{1, 2, \dots, n\}$, s.t. each symbol occurs **at most** once in each row and in each column.
- ▶ The number of filled cells is its **size**. Their positions determine the **shape**:

$$P = \begin{pmatrix} 1 & \cdot & 3 & 4 & 6 \\ 2 & \cdot & 5 & \cdot & 4 \\ \cdot & \cdot & 4 & 5 & 1 \\ \cdot & \cdot & 2 & \cdot & 3 \end{pmatrix} \in \mathcal{PLR}_{4 \times 5, 12}^6 \rightarrow \text{Sh}(P) = \begin{pmatrix} 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \end{pmatrix}$$

- ▶ $r = s = n \rightarrow$ Partial Latin square.
- ▶ Size $r \cdot s \rightarrow$ Latin rectangle (square if $r = s$).
- ▶ **Applications:** Algebra (quasigroups), Experimental Designs, Cryptography.

PARTIAL LATIN RECTANGLES.

$|PLR_{r \times s}^n|$ upper bounded for $r = s = n$ [Ghandehari, 2005].

$|IC|$ and $|MC|$ lower bounded for $r = s = n \leq 6$ [Adams, 2003].

Order n	Size m	$ PLS_{n,m} $
1	1	1
	2	8
2	1	16
	2	8
	3	2
	4	27
	5	270
	6	1,278
3	4	3,078
	5	3,834
	6	2,412
	7	756
	8	108
	9	12
	10	64
4	1	64
	2	1,728
	3	25,920
	4	239,760
	5	1,437,696
	6	5,728,896
	7	15,326,208
	8	27,534,816
	9	32,971,008
	10	25,941,504
	11	13,153,536
	12	4,215,744
	13	847,872
	14	110,592
	15	9,216
	16	576

Some exact values have recently been obtained by applying Gröbner bases in an equivalent planar assignment problem:

$$\begin{cases} \sum_{k \in [n]} x_{ijk} \leq 1, \forall i, j \in [n], \\ \sum_{j \in [n]} x_{ijk} \leq 1, \forall i, k \in [n], \\ \sum_{i \in [n]} x_{ijk} \leq 1, \forall j, k \in [n], \\ \sum_{i,j,k \in [n]} x_{ijk} = m, \\ x_{ijk} \in \{0, 1\}, \forall i, j, k \in [n], \end{cases}$$

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TYPE AND STRUCTURE OF A PLR.

$$\sum_{i,j,k \in [n]} x_{ijk} = m \rightarrow \begin{cases} \sum_{j,k \in [n]} x_{ijk} = T_1(P, i), \leftarrow \text{Rows.} \\ \sum_{i,k \in [n]} x_{ijk} = T_2(P, j), \leftarrow \text{Columns.} \\ \sum_{i,j \in [n]} x_{ijk} = T_3(P, k), \leftarrow \text{Symbols.} \end{cases}$$

$$P = \begin{pmatrix} 1 & \cdot & 3 & 4 & 6 \\ 2 & \cdot & 5 & \cdot & 4 \\ \cdot & \cdot & 4 & 5 & 1 \\ \cdot & \cdot & 2 & \cdot & 3 \end{pmatrix} \in \mathcal{PLR}_{4 \times 5, 12}^6 \rightarrow \text{Sh}(P) = \left(\begin{array}{c|ccccc} & 2 & 0 & 4 & 2 & 4 \\ \hline 4 & 1 & 0 & 1 & 1 & 1 \\ 3 & 1 & 0 & 1 & 0 & 1 \\ 3 & 0 & 0 & 1 & 1 & 1 \\ 2 & 0 & 0 & 1 & 0 & 1 \end{array} \right)$$

Row type: $T_1(P) = (4, 3, 3, 2)$

Column type: $T_2(P) = (2, 0, 4, 2, 4)$.

Symbol type: $T_3(P) = (2, 2, 2, 3, 2, 1)$.

TYPE AND STRUCTURE OF A PLR.

- ▶ $\mathcal{T}_{l,w} = \{(t_1, \dots, t_l) \text{ of weight } \sum_{i \in [l]} t_i = w, \text{ s.t. } t_i \in \mathbb{N}\}$.
- ▶ **Structure** of $T \in \mathcal{T}_{l,w}$: $\text{st}(T) = w^{\lambda_w^T} \dots 1^{\lambda_1^T}$, where λ_i^T is the number of occurrences of i in T .
- ▶ $\mathcal{Z}_{l,w}$: Set of possible structures of length l and weight w .

$$P = \begin{pmatrix} 1 & \cdot & 3 & 4 & 6 \\ 2 & \cdot & 5 & \cdot & 4 \\ \cdot & \cdot & 4 & 5 & 1 \\ \cdot & \cdot & 2 & \cdot & 3 \end{pmatrix} \in \mathcal{PLR}_{4 \times 5, 12}^6 \rightarrow \text{Sh}(P) = \left(\begin{array}{c|cccccc} & 2 & 0 & 4 & 2 & 4 \\ \hline 4 & 1 & 0 & 1 & 1 & 1 \\ 3 & 1 & 0 & 1 & 0 & 1 \\ 3 & 0 & 0 & 1 & 1 & 1 \\ 2 & 0 & 0 & 1 & 0 & 1 \end{array} \right)$$

Row type: $T_1(P) = (4, 3, 3, 2) \rightarrow \text{st}(T_1(P)) = 43^22$.

Column type: $T_2(P) = (2, 0, 4, 2, 4) \rightarrow \text{st}(T_2(P)) = 4^22^2$.

Symbol type: $T_3(P) = (2, 2, 2, 3, 2, 1) \rightarrow \text{st}(T_3(P)) = 32^41$.

TYPE AND STRUCTURE OF A PLR.

- ▶ Given $R \in \mathcal{T}_{r,w}$, $C \in \mathcal{T}_{s,w}$ and $S \in \mathcal{T}_{n,w}$:

$$\mathcal{PLR}_{(R,C)}^n = \{P \in \mathcal{PLR}_{r \times s}^n : T_1(P) = R \text{ and } T_2(P) = C\}.$$

$$\mathcal{PLR}_{(R,C,S)} = \{P \in \mathcal{PLR}_{r \times s}^n : T_1(P) = R, T_2(P) = C \text{ and } T_3(P) = S\}.$$

LEMMA

$|\mathcal{PLR}_{(R,C)}^n|$ and $|\mathcal{PLR}_{(R,C,S)}|$ only depend on the structures of R, C, S .

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LEMMA

$|\mathcal{PLR}_{(R,C)}^n|$ and $|\mathcal{PLR}_{(R,C,S)}|$ only depend on the structures of R, C, S .

- ▶ Given $z_1 \in \mathcal{Z}_{l_1,w}$, $z_2 \in \mathcal{Z}_{l_2,w}$ and $z_3 \in \mathcal{Z}_{l_3,w}$:

$$\Delta_{r \times s}^n(z_1, z_2) = |\mathcal{PLR}_{(R,C)}^n|, \forall R \in \mathcal{T}_{r,w}, C \in \mathcal{T}_{s,w}, \text{ s.t. } \text{st}(R) = z_1, \text{st}(C) = z_2.$$

$$\Delta_{r \times s}^n(z_1, z_2, z_3) = |\mathcal{PLR}_{(R,C,S)}^n|, \forall R \in \mathcal{T}_{r,w}, C \in \mathcal{T}_{s,w}, S \in \mathcal{T}_{n,w}, \text{ s.t. } \text{st}(R) = z_1, \\ \text{st}(C) = z_2, \text{st}(S) = z_3.$$

TYPE AND STRUCTURE OF A PLR.

PROPOSITION

$$|\mathcal{PLR}_{r \times s}^n| = \sum_{(l,l') \in [r] \times [s]} \sum_{w \in [l \cdot s]} \sum_{(z,z') \in \mathcal{Z}_{l,w} \times \mathcal{Z}_{l',w}} \frac{r!}{(r-l)! \cdot \prod_{i \in [w]} \lambda_i!} \cdot \Delta_{r \times s}^n(z, z')$$

Where:

$$\Delta_{l \times s}^n(z, z') = \sum_{l'' \in [n]} \sum_{z'' \in \mathcal{Z}_{l'',w}} \Delta_{r \times s}^n(z, z', z'').$$

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PROBLEM

How to obtain $\Delta_{r \times s}^n(z, z')$ and $\Delta_{r \times s}^n(z, z', z'')$?

TYPE AND STRUCTURE OF A PLR.

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$$\Delta_{r \times s}^n(z, z').$$

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$$|\mathcal{PLR}_{(R,C)}^n| \rightarrow \mathfrak{A}(R, C)$$

- ▶ $\mathfrak{A}(R, C)$: $(0, 1)$ -matrices having R and C as row and column sum vectors.
- ▶ \mathcal{PLR}_M^n : Set of PLR of n symbols having $M \in \mathfrak{A}(R, C)$ as shape.

$$|\mathcal{PLR}_{(R,C)}^n| = \sum_{M \in \mathfrak{A}(R,C)} |\mathcal{PLR}_M^n|.$$

Equivalent problems:

- ▶ n -edge-colouring a bipartite graph of incidence matrix $\text{Sh}(P)$ (Existence problem is NP-complete even for $n = 3$ [Holyer, 1981]).
- ▶ *1-color tomography problem* [Kuba, 1999]: Reconstructing a binary matrix starting from its row and column sums.

$$\Delta_{r \times s}^n(z, z').$$

Gale-Ryser theorem [Gale, Ryser, 1957]: $\mathfrak{A}(R, C) \neq \emptyset \Leftrightarrow C \preceq R^*$.
 $R = (3, 5, 2, 2) \rightarrow R^* = (4, 4, 2, 1, 1) \succeq (3, 3, 3, 2, 1)$. (*Dominance order*).

Formulas and algorithms:

- ▶ Monomial symmetric functions [Sukhatme, 1938; David, 1951 (≤ 12)].
- ▶ Character of the symmetric group [Snapper, 1971].
- ▶ Lower bound [Wei, 1982].
- ▶ Recurrence formulas [Wang, 1988; Wang and Zhang, 1998; Pérez Salvador, 2002].
- ▶ General formulas [Dias, 2002].
- ▶ Asymptotic methods [Barvinok, 2010].
- ▶ Combinatorial methods [Brualdi, 1980; Brualdi, 2006; Fonseca, 2009].
- ▶ Simulation methods (social networks, ecology) [Snijders, 1991; Rao, 1996; Chen, 2005; Bezakova, 2007; Blanchet, 2009].

$$\Delta_{r \times s}^n(z, z').$$

Algebraic approach: Gröbner bases of boolean ideals for counting problems [Bayer, 1982; Alon, 1995; Bernasconi, 1997].

THEOREM

$R = (r_1, \dots, r_r) \in \mathcal{T}_{r,w}$ and $C = (c_1, \dots, c_s) \in \mathcal{T}_{s,w}$ s.t. $C \preceq R^*$.

$\mathfrak{A}(R, C) = V(I)$, where:

$$\begin{aligned} I = & \langle \left(\sum_{j \in [s]} x_{ij} - r_i \right) : i \in [r] \rangle + \langle \left(\sum_{i \in [r]} x_{ij} - c_j \right) : j \in [s] \rangle \\ & + \langle x_{ij} \cdot (1 - x_{ij}) : i \in [r], j \in [s] \rangle \subseteq \mathbb{Q}[x_{11}, \dots, x_{rs}]. \end{aligned}$$

Moreover, $|\mathfrak{A}(R, C)| = \dim_{\mathbb{Q}}(\mathbb{Q}[x_{11}, \dots, x_{rs}]/I)$.

TYPE AND STRUCTURE OF A PLR.

PROPOSITION

$$|\mathcal{PLR}_{r \times s}^n| = \sum_{(l,l') \in [r] \times [s]} \sum_{w \in [l \cdot s]} \sum_{(z,z') \in \mathcal{Z}_{l,w} \times \mathcal{Z}_{l',w}} \frac{r!}{(r-l)! \cdot \prod_{i \in [w]} \lambda_i!} \cdot \Delta_{r \times s}^n(z, z')$$

Where:

$$\Delta_{l \times s}^n(z, z') = \sum_{l'' \in [n]} \sum_{z'' \in \mathcal{Z}_{l'',w}} \Delta_{r \times s}^n(z, z', z'').$$

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$R = (r_1, \dots, r_r) \in \mathcal{T}_{r,w}$, $C = (c_1, \dots, c_s) \in \mathcal{T}_{s,w}$ and $S = (s_1, \dots, s_n) \in \mathcal{T}_{n,w}$.

$\mathcal{PLR}_{(R,C,S)} = V(I)$, where:

$$\begin{aligned} I = & \langle (\sum_{i \in [r]} x_{ijk}) \cdot (1 - \sum_{i \in [r]} x_{ijk}) : j \in [s], k \in [n] \rangle + \langle r_i - \sum_{j \in [s], k \in [n]} x_{ijk} : i \in [r] \rangle + \\ & \langle (\sum_{j \in [s]} x_{ijk}) \cdot (1 - \sum_{j \in [s]} x_{ijk}) : i \in [r], k \in [n] \rangle + \langle c_j - \sum_{i \in [r], k \in [n]} x_{ijk} : j \in [s] \rangle + \\ & \langle (\sum_{k \in [n]} x_{ijk}) \cdot (1 - \sum_{k \in [n]} x_{ijk}) : i \in [r], j \in [s] \rangle + \langle s_k - \sum_{i \in [r], j \in [s]} x_{ijk} : k \in [n] \rangle + \\ & \langle x_{ijk} \cdot (1 - x_{ijk}) : i \in [r], j \in [s], k \in [n] \rangle \subseteq \mathbb{Q}[x_{111}, \dots, x_{rsn}]. \end{aligned}$$

Moreover, $|\mathcal{PLR}_{(R,C,S)}| = \dim_{\mathbb{Q}}(\mathbb{Q}[x_{111}, \dots, x_{rsn}]/I)$.

$$\Delta_{r \times s}^n(z, z', z'').$$

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Two PLR $P = (p_{rc})$, $Q = (q_{rc}) \in \mathcal{PLR}_{(R,C,S)}$ are isotopic if the following system has solution:

$$\left\{ \begin{array}{l} \sum_{j \in [r]} x_{ij} = 1, \forall i \in [r], \\ \sum_{j \in [s]} y_{ij} = 1, \forall i \in [s], \\ \sum_{j \in [n]} z_{ij} = 1, \forall i \in [n], \\ \sum_{i \in [r]} x_{ij} = 1, \forall j \in [r], \\ \sum_{i \in [s]} y_{ij} = 1, \forall j \in [s], \\ \sum_{i \in [n]} z_{ij} = 1, \forall j \in [n], \\ x_{ik} \cdot y_{jl} \cdot (z_{p_{ij}q_{kl}} - 1) = 0, \forall i, j \in [r] \text{ and } k, l \in [s] \text{ s.t. } p_{ij}, q_{kl} \in [n], \\ x_{ik} \cdot y_{jl} = 0, \forall i, j \in [r] \text{ and } k, l \in [s] \text{ s.t. } q_{kl} = \emptyset, \end{array} \right.$$

$$\Delta_{r \times s}^n(z, z', z'').$$

$Z = (z_1, z_2, z_3)$			$\Delta(Z)$	$\Delta_I(Z)$	$\Delta_P(Z)$
z_1	z_2	z_3			
1	1	1	1	1	1
2	1 ²	1 ²	2	1	1
1 ²	1 ²	1 ²	4	1	1
3	1 ³	1 ³	6	1	1
21	21	21	1	1	1
		1 ³	6	1	1
	1 ³	1 ³	18	1	1
1 ³	1 ³	1 ³	36	1	1
4	1 ⁴	1 ⁴	24	1	1
31	21 ²	21 ²	4	1	1
		1 ⁴	24	1	1
	1 ⁴	1 ⁴	96	1	1
2 ²	2 ²	2 ²	2	1	1
		21 ²	4	1	1
		1 ⁴	24	1	1
	21 ²	21 ²	12	2	2
		1 ⁴	48	1	1
	1 ⁴	1 ⁴	144	1	1
21 ²	21 ²	21 ²	40	5	3
		1 ⁴	120	2	2
	1 ⁴	1 ⁴	288	1	1
1 ⁴	1 ⁴	1 ⁴	576	1	1
5	1 ⁵	1 ⁵	120	1	1
41	21 ³	21 ³	18	1	1
		1 ⁵	120	1	1
	1 ⁵	1 ⁵	600	1	1

$Z = (z_1, z_2, z_3)$			$\Delta(Z)$	$\Delta_I(Z)$	$\Delta_P(Z)$
z_1	z_2	z_3			
32	2 ² 1	2 ² 1	6	2	2
		21 ³	24	2	2
		1 ⁵	120	1	1
	21 ³	21 ³	90	3	3
		1 ⁵	360	1	1
	1 ⁵	1 ⁵	1,200	1	1
31 ²	31 ²	2 ² 1	4	1	1
		21 ³	24	1	1
		1 ⁵	120	1	1
	2 ² 1	2 ² 1	12	2	2
		21 ³	60	3	3
		1 ⁵	240	1	1
	21 ³	21 ³	252	5	4
		1 ⁵	840	2	2
	1 ⁵	1 ⁵	2,400	1	1
2 ² 1	2 ² 1	2 ² 1	58	8	4
		21 ³	180	8	6
		1 ⁵	600	2	2
	21 ³	21 ³	504	8	6
		1 ⁵	1440	2	2
	1 ⁵	1 ⁵	3,600	1	1
21 ³	21 ³	21 ³	1,296	8	4
		1 ⁵	3,240	2	2
	1 ⁵	1 ⁵	7,200	1	1
1 ⁵	1 ⁵	1 ⁵	14,400	1	1

3-SEMINETS.

THEOREM

The number of isomorphism classes of 3-seminets with one, two, three, four and five points are 1, 4, 11, 52 and 220, respectively. That of paratopism classes are 1, 2, 5, 18 and 59, respectively. □

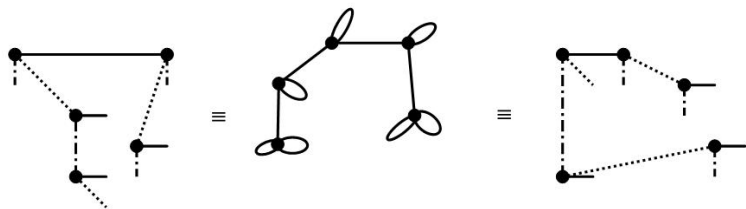
34 of the 85 3-seminets are uniquely determined by the PLR's structures.

$(1, 1, 1)$	$(2, 1^2, 1^2)$	$(1^2, 1^2, 1^2)$	$(3, 1^3, 1^3)$	$(21, 21, 21)$	$(21, 21, 1^3)$	$(21, 1^3, 1^3)$	$(1^3, 1^3, 1^3)$	
$(4, 1^4, 1^4)$	$(31, 21^2, 21^2)$	$(31, 21^2, 1^4)$	$(31, 1^4, 1^4)$	$(2^2, 2^2, 2^2)$	$(2^2, 2^2, 21^2)$	$(2^2, 2^2, 1^4)$		

The rest are not uniquely determined:

$(21^2, 21^2, 21^2)$	$(21^2, 21^2, 21^2)$	$(21^2, 21^2, 21^2)$

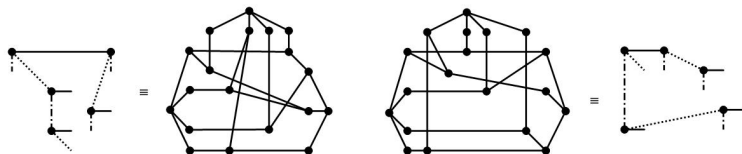
SEMINETS-GRAPHS.



Seminet-graph (G_1):

- ▶ Vertices \equiv points.
- ▶ Connected vertices \equiv collinear points.
- ▶ Lines containing only one point are identified with loops.
- ▶ It can be related to more than one paratopism class.

SEMINETS-GRAPHS.



Seminet-graph (G_2):

- ▶ Vertices \equiv Points, lines and parallel classes.

$$\{u_1, \dots, u_n\} \cup \{v_{1,1}, \dots, v_{1,l_1}, v_{2,1}, \dots, v_{3,l_3}\} \cup \{w_1, w_2, w_3\}.$$

- ▶ Each vertex w_i is connected to all the vertices $v_{i,j}$, for all $i \in [3]$ and $j \in [l_i]$.
- ▶ each vertex $v_{i,j}$ is connected to those vertices u_k such that the line related to the former contains the point associated to the latter.
- ▶ Uniquely related to a paratopism class.
- ▶ If G_2 is not acyclic, then its girth is 6.

SEMINETS-GRAPHS.

We have considered:

- ▶ $l \equiv$ Number of vertices of $G_1(S)$ contained in at least one loop.
- ▶ $a \equiv$ Number of articulation points of $G_1(S)$.
- ▶ $t \equiv$ Number of transversal of $G_1(S)$.
- ▶ $c \equiv$ Clustering coefficient of $G_1(S)$.
- ▶ $st_1 \equiv$ Number of spanning trees in $G_1(S)$.
- ▶ $st_2 \equiv$ Number of spanning trees in $G_2(S)$.

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CLASSIFICATION OF 3-SEMINETS WITH AT MOST 5 POINTS.

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THANK YOU FOR YOUR ATTENTION!!