# ON THE $n$-COLOR RADO NUMBER FOR THE EQUATION 

$$
x_{1}+x_{2}+\cdots+x_{k}+c=x_{k+1}
$$

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Abstract. For integers $k, n, c$ with $k, n \geq 1$, the $n$-color Rado number $R_{k}(n, c)$ is defined to be the least integer $N$, if it exists or $\infty$ otherwise, such that for every $n$-coloring of the set $\{1,2, \ldots, N\}$, there exists a monochromatic solution in that set to the equation

$$
x_{1}+x_{2}+\cdots+x_{k}+c=x_{k+1}
$$

In this paper, we mostly restrict to the case $c \geq 0$, and consider two main issues regarding $R_{k}(n, c)$ : is it finite or infinite, and when finite, what is its value? Very few results are known so far on either one.

On the first issue, we formulate a general conjecture, namely that $R_{k}(n, c)$ should be finite if and only if every divisor $d \leq n$ of $k-1$ also divides $c$. The "only if" part of the conjecture is shown to hold, as well as the "if" part in the cases where either $k-1$ divides $c$, or $n \geq k-1$, or $k \leq 7$, except for two instances to be published separately.

On the second issue, we obtain new bounds on $R_{k}(n, c)$ and determine exact formulae in several new cases, including $R_{3}(3, c)$ and $R_{4}(3, c)$. As for the case $R_{2}(3, c)$, first settled by Schaal in 1995, we provide a new shorter proof.

Finally, the problem is reformulated as a Boolean satisfiability problem, allowing the use of a SAT solver to treat some instances.

## 1. Introduction

Throughout the paper, we shall denote by $\mathbb{Z}, \mathbb{N}$ and $\mathbb{N}_{+}$the set of integers, nonnegative integers and positive integers, respectively. Let $n \in \mathbb{N}_{+}$. An $n$-coloring of a set $A$ is a function

$$
\Delta: A \longrightarrow C
$$

where $C$ is some finite set of cardinality $|C|=n$. Here, we shall mostly deal with $n$-colorings of integer intervals $[1, N]$, where

$$
[a, b]=\{a, a+1, \ldots, b\}
$$

for integers $a \leq b$.
Given an $n$-coloring $\Delta$ of $[1, N]$ and a linear equation $L$ in $k+1$ variables with integer coefficients, a solution $\left(x_{1}, \ldots, x_{k}, x_{k+1}\right)$ to $L$ is said to be monochromatic if $\Delta\left(x_{1}\right)=\Delta\left(x_{2}\right)=\cdots=\Delta\left(x_{k+1}\right)$.

[^0]1.1. Some earlier results. In 1916, Schur [22] proved that for every $n \geq 1$, there exists a least integer $S_{2}(n)=N$, such that for every $n$-coloring of $[1, N]$, there exists a monochromatic solution to the equation $x_{1}+x_{2}=x_{3}$.

The integers $S_{2}(n)$ are called the Schur numbers and are currently known only for $n \leq 4$, namely: $S_{2}(1)=2, S_{2}(2)=5, S_{2}(3)=14$ and $S_{2}(4)=45$. While these stated values for $n \leq 3$ can easily be settled by hand, the one for $n=4$ relies on an exhaustive computer search [2]. For $n=5$, the currently available bounds for $S_{2}(5)$ are $161 \leq S_{2}(5) \leq 306$, and it is conjectured in [8] that the lower bound 161, established in [7], is perhaps sharp. For the upper bound 306, see [18]. For general $n \geq 1$, Schur [22] obtained the following bounds:

$$
\left(3^{n}+1\right) / 2+1 \leq S_{2}(n) \leq\lfloor n!e\rfloor+1 .
$$

Slightly improved upper bounds were subsequently provided by Whitehead [23] and Honghui Wan [10, whereas for lower bounds, the inequalities

$$
S_{2}(m+n) \geq 2 S_{2}(m) S_{2}(n)-S(m)-S(n)+1
$$

of Abbott and Hanson [1], and $S_{2}(5) \geq 161$ of Fredricksen and Sweet [8], together yield the sharpening $S_{2}(n) \geq c 322^{n / 5} \geq c 3.17^{n}$ for $n \geq 6$, where $c$ is some absolute positive constant.

In 1933, Rado [13] generalized the work of Schur to arbitrary systems of linear equations. Given $n \geq 1$ and a system of linear equations $L$, the least integer $N$ (if it exists) such that for every $n$-coloring of the set $[1, N]$, there is a monochromatic solution to the system $L$, is called the $n$-color Rado number for $L$. If no such integer $N$ exists, then this Rado number is defined to be infinite.

Given a linear system $L$ as above, and $n \geq 1$, there are two main issues regarding its corresponding $n$-color Rado number: is it finite or infinite? When it is finite, what is its exact value?

In particular, for the equation $x_{1}+x_{2}+\cdots+x_{k}=x_{k+1}$ where $k \geq 2$, Rado's results imply that for all $n \geq 1$, the corresponding $n$-color Rado number is actually finite, i.e., there exists a least integer $S_{k}(n)=N$ such that for every $n$-coloring of $[1, N]$, there is a monochromatic solution of that equation [13].

In 1982, Beutelspacher and Brestovansky [5] showed that $S_{k}(2)=k^{2}+k-1$ for $k \geq 2$. More than twenty years later, Sanz [18] established the value $S_{3}(3)=43$ with an exhaustive computer search.

Burr and Loo [3] were able to determine the 2-color Rado numbers for the equations $x_{1}+x_{2}+c=x_{3}$ and $x_{1}+x_{2}=k x_{3}$ for every integer $c$ and for every positive integer $k$. There are several results due to Schaal and other authors, about 2-color and 3 -color Rado numbers for particular equations; see [11, 12, 15, 21].
1.2. The main conjecture. Let $n, k, c$ be integers with $n, k \geq 1$ and $c \geq 0$. In this paper, we shall be concerned with the above-mentioned equation

$$
\begin{equation*}
x_{1}+x_{2}+\cdots+x_{k}+c=x_{k+1} . \tag{1}
\end{equation*}
$$

Notation 1.1. We shall denote by $R_{k}(n, c)$ the $n$-color Rado number corresponding to equation (1), i.e., the smallest positive integer $N$, if it exists, such that every $n$-coloring of $[1, N]$ admits a monochromatic solution to it.

For what values of the parameters $n, k, c$ is $R_{k}(n, c)$ finite? By Rado's result, recalled above, it holds that $R_{k}(n, 0)=S_{k}(n)$ is always finite. Now, for $n=2$,

Schaal [19] showed that $R_{k}(2, c)$ is finite if and only if $k$ or $c$ is even, in which case

$$
R_{k}(2, c)=(k+1)^{2}+(c-1)(k+2) .
$$

Later on, he further showed that the 3-color Rado number $R_{2}(3, c)$ is always finite [20], and obtained the exact value $R_{2}(3, c)=13 c+14$ for all $c \geq 0$.

In this paper, we provide further instances of $n, k, c$ for which $R_{k}(n, c)$ is shown to be finite or infinite, respectively. These results, as well as those of Schaal on $R_{k}(2, c)$ and $R_{2}(3, c)$, fit into one single conjecture.

Conjecture 1.2. For integers $k, n, c$ with $k \geq 2, n \geq 1$ and $c \geq 0$, the $n$-color Rado number $R_{k}(n, c)$ is finite if and only if every divisor $d \leq n$ of $k-1$ also divides $c$.

In the sequel, we settle this conjecture if either $k-1$ divides $c$, or $n \geq k-1$, or $k \leq 7$, except for two special cases which need a completely different approach and will be presented elsewhere.
1.3. Contents. In Section 2 we settle the "only if" part of Conjecture 1.2 by giving a sufficient condition on $n, k, c$ ensuring that $R_{k}(n, c)$ is infinite. In Section 3, we study how $R_{k}(n, c)$ behaves under changes on $n$ and $c$, and then exploit the results to settle Conjecture 1.2 in cases $n \geq k-1$ and $k \leq 7$. In Section 4 we introduce new numbers $S_{k}^{*}(n)$, smaller and easier to study than $S_{k}(n)$, and show how they help bounding $R_{k}(n, c)$ from below. This is then exploited in Section 5 in conjunction with SAT solvers, to get new formulae for some instances of $R_{k}(n, c)$.

## 2. An obstacle to finiteness

We start by treating the cases $n=1$ or $k=1$. For $n=1$, the Rado number $R_{k}(n, c)$ is given by the formula

$$
\begin{equation*}
R_{k}(1, c)=k+c \tag{2}
\end{equation*}
$$

for all $k \geq 1$ and $c \geq 0$, as readily verified. The case $k=1$ is also easy to determine. Indeed, for $c=0$ we have

$$
R_{1}(n, 0)=1
$$

for all $n \geq 2$, whereas the following holds for $c \geq 1$.
Proposition 2.1. For all $n \geq 2$ and $c \geq 1$, we have $R_{1}(n, c)=+\infty$.
Proof. The statement follows from the following 2-coloring of $\mathbb{N}_{+}$:

$$
\begin{aligned}
\Delta: \mathbb{N}_{+} & \longrightarrow\{0,1\} \\
x & \longmapsto \text { the class of }\lceil x / c\rceil \bmod 2 .
\end{aligned}
$$

Since $\lceil(x+c) / c\rceil=\lceil x / c\rceil+1$, implying $\Delta(x+c) \equiv \Delta(x)+1 \bmod 2$, there are no monochromatic solution to the equation $x_{1}+c=x_{2}$, thereby implying $R_{1}(n, c)=$ $+\infty$ as stated.

Therefore, from now on, we shall assume $k, n \geq 2$. Here is an obstacle to the finiteness of $R_{k}(n, c)$, which settles the "only if" part of Conjecture 1.2

Proposition 2.2. If there exists a divisor $d \leq n$ of $k-1$ which does not divide $c$, then $R_{k}(n, c)=+\infty$.

Proof. Color each integer by its class mod $d$, taken in the set $\{0,1, \ldots, d-1\}$. This yields a $d$-coloring of $\mathbb{N}_{+}$. Let $x_{1}, \ldots, x_{k+1} \in \mathbb{N}_{+}$be monochromatic for this coloring, say all of the same color class $r \bmod d$. Then

$$
x_{1}+\cdots+x_{k}-x_{k+1} \equiv(k-1) r \equiv 0 \bmod d,
$$

where the second congruence follows from the hypothesis that $d$ divides $k-1$. Now since $c \not \equiv 0 \bmod d$ by hypothesis, it follows that $\left(x_{1}, \ldots, x_{k+1}\right)$ cannot satisfy the equation

$$
x_{1}+\cdots+x_{k}-x_{k+1}=-c .
$$

Therefore this equation does not admit any monochromatic solution. It follows that $R_{k}(d, c)=+\infty$, whence also $R_{k}(n, c)=+\infty$, as claimed.

In fact, the above condition is the only general one we are aware of which implies $R_{k}(n, c)=\infty$. This is what led us to formulate Conjecture 1.2,

## 3. VARYing $n$ and $c$

In this section, we shall vary the parameters $n$ and $c$ and show how this affects the value of $R_{k}(n, c)$. This study will ultimately allow us to settle Conjecture 1.2 in the cases where either $k-1$ divides $c$, or $n \geq k-1$, or $k \leq 7$, except for two key instances which need a different approach and will be published separately.
3.1. Reducing $n$. Our first result is a relation between $n$-color and $(n-1)$-color Rado numbers. Trivially, one has $R_{k}(n, c) \geq R_{k}(n-1, c)$, but a sharper inequality holds.

Lemma 3.1. Let $k, n, c \in \mathbb{N}_{+}$. Then $R_{k}(n, c) \geq(k+1) R_{k}(n-1, c)+c-1$.
Proof. To ease notation, set $M=R_{k}(n-1, c)$ and $N=k M+c$. Thus

$$
(k+1) R_{k}(n-1, c)+c-1=N+M-1,
$$

and our aim is to show that $R_{k}(n, c) \geq N+M-1$. In order to do that, it suffices to construct an $n$-coloring of the integer interval [ $1, N+M-2$ ] for which that interval contains no monochromatic $x_{i}$ 's satisfying the equation

$$
\begin{equation*}
x_{1}+\cdots+x_{k}+c=x_{k+1} . \tag{3}
\end{equation*}
$$

The minimality property of $R_{k}(n, c)$ will then imply the desired inequality.
By definition of $M$, there exists an $(n-1)$-coloring

$$
\begin{equation*}
\Delta:[1, M-1] \longrightarrow[1, n-1] \tag{4}
\end{equation*}
$$

such that $[1, M-1]$ contains no $\Delta$-monochromatic $x_{i}$ 's satisfying equation (3). We now extend (4) to an $n$-coloring

$$
\Delta^{\prime}:[1, N+M-2] \longrightarrow[1, n]
$$

as follows:

$$
\Delta^{\prime}(x)=\left\{\begin{array}{cl}
\Delta(x) & \text { if } x \in[1, M-1], \\
n & \text { if } x \in[M, N-1] \\
\Delta(x-(N-1)) & \text { if } x \in[N, N+M-2] .
\end{array}\right.
$$

It remains to show that $[1, N+M-2]$ is free of a $\Delta^{\prime}$-monochromatic solution to (33). Assuming the contrary, let $x_{1}, \ldots, x_{k+1} \in[1, N+M-2]$ satisfy (3) and be of the same $\Delta^{\prime}$-color in $[1, n]$.

First, that common color cannot be $n$, for otherwise all $x_{i}$ 's would belong to [ $M, N-1$ ], thereby yielding

$$
x_{k+1}=x_{1}+\cdots+x_{k}+c \geq k M+c=N
$$

a contradiction.
Therefore, that common color of $x_{1}, \ldots, x_{k+1}$ belongs to [ $\left.1, n-1\right]$. Hence, some $x_{i}$ 's belong to $[1, M-1]$ and the rest to $[N, N+M-2]$. How do they distribute among these two intervals? First, we may assume that

$$
x_{1} \leq \cdots \leq x_{k}
$$

Note further that $x_{k}<x_{k+1}$, since

$$
x_{k+1}=x_{1}+\cdots+x_{k}+c \geq k+c \geq 2
$$

by hypothesis on $c$.
Clearly, the $x_{i}$ 's cannot all belong to $[1, M-1]$ by our hypothesis on (4). It follows that the largest one, namely $x_{k+1}$, belongs to $[N, N+M-2]$. We claim that

$$
x_{1}, \ldots, x_{k-1} \in[1, M-1] \text { and } x_{k} \in[N, N+M-2] .
$$

Indeed, by (3) and the fact that $x_{k+1} \in[N, N+M-2]$, at most one among $\left\{x_{1}, \ldots, x_{k}\right\}$ may belong to $[N, N+M-2]$, since

$$
2 N>N+M-2
$$

as readily verified. Similarly, at least one among $\left\{x_{1}, \ldots, x_{k}\right\}$ must belong to [ $N, N+M-2]$, for otherwise $x_{k} \leq M-1$, and by (3) we would have

$$
x_{k+1} \leq k(M-1)+c<N,
$$

a contradiction. Subtracting $N-1$ from $x_{k}$ and $x_{k+1}$, it follows that

$$
x_{1}, \ldots, x_{k-1}, x_{k}-(N-1), x_{k+1}-(N-1)
$$

are $\Delta$-monochromatic, belong to $[1, M-1]$, and satisfy (3), a contradiction. Therefore, the interval $[1, N+M-2]$ contains no $\Delta^{\prime}$-monochromatic solution to (3), and the proof is finished.

Applying the above result inductively, we obtain the following absolute lower bound.

Theorem 3.2. Let $k, n, c \in \mathbb{N}_{+}$. Then $R_{k}(n, c) \geq \frac{(k+1)^{n}-1}{k}(k+c-1)+1$.
Proof. The inequality holds for $n=1$, since $R_{k}(1, c)=k+c$ by (2). For general $n \geq 2$, we apply induction and Lemma 3.1.
3.2. Reducing $c$. We now vary the parameter $c$ and show how $R_{k}(n, c)$ is affected. Several consequences will then be presented in subsequent sections.

Lemma 3.3. Let $\alpha, \beta \in \mathbb{Z}$ such that $\alpha \geq 1$ and $\beta \geq 1-\alpha$. Then for all integers $k, n, c$ with $k, n \geq 2$ and $c \geq 0$, we have

$$
R_{k}(n, \alpha c-\beta(k-1)) \leq \alpha R_{k}(n, c)+\beta
$$

Proof. Let $I$ be the integer interval

$$
I=\left[1, \alpha R_{k}(n, c)+\beta\right]
$$

and let $\Delta: I \rightarrow[1, n]$ be any $n$-coloring of $I$. We must show that there exist $x_{1}, \ldots, x_{k+1} \in I$ which are monochromatic under $\Delta$ and which satisfy the equation

$$
x_{1}+\cdots+x_{k}+(\alpha c-\beta(k-1))=x_{k+1} .
$$

By the defining minimality property of $R_{k}(n,(\alpha c-\beta(k-1)))$, this will suffice to establish the stated inequality.

Let $J=\left[1, R_{k}(n, c)\right]$. We affinely embed $J$ in $I$ as follows:

$$
\begin{aligned}
h: J & \longrightarrow I \\
\quad z & \longmapsto \alpha z+\beta .
\end{aligned}
$$

Note that our hypotheses on $\alpha, \beta$ ensure that if $z \geq 1$, then $\alpha z+\beta \geq 1$ and, more generally, that $h(J) \subseteq I$. By composing $\Delta$ with $h$, we obtain an $n$-coloring $\Delta^{\prime}=\Delta \circ h$ on $J$, namely the map $\Delta^{\prime}: J \rightarrow[1, n]$ defined by

$$
\Delta^{\prime}(z)=\Delta(\alpha z+\beta)
$$

for all $z \in J$. Now, by definition of the upper bound $R_{k}(n, c)$ of $J$, there exists a $\Delta^{\prime}$-monochromatic solution

$$
z_{1}+\cdots+z_{k}+c=z_{k+1}
$$

with $z_{i} \in J$ for all $i$. Multiplying by $\alpha$ and adding $\beta$ 's, we get

$$
\left(\alpha z_{1}+\beta\right)+\cdots+\left(\alpha z_{k}+\beta\right)+(\alpha c-\beta(k-1))=\left(\alpha z_{k+1}+\beta\right) .
$$

Now $\left(\alpha z_{i}+\beta\right) \in I$ for all $i$, and since the $z_{i}$ are $\Delta^{\prime}$-monochromatic, it follows that the $\left(\alpha z_{i}+\beta\right)$ are $\Delta$-monochromatic. This concludes the proof of the lemma.
3.3. The case $n \geq k-1$. A first consequence of the above lemma is the verification of Conjecture 1.2 in the cases where either $c \equiv 0 \bmod k-1$ (Proposition 3.4) or $n \geq k-1$ (Theorem 3.5).

Proposition 3.4. If $c$ is a multiple of $k-1$, then $R_{k}(n, c)$ is finite. More precisely, if $c=q(k-1)$ for some integer $q \geq 1$, then

$$
R_{k}(n, c) \leq(q+1) S_{k}(n)-q .
$$

Proof. It follows from Lemma 3.3, with values $\beta=-q, \alpha=q+1$ (so that $\alpha+\beta \geq 1$ as required), and $c=0$ (the $c$ in that lemma, not the present one), that

$$
\begin{aligned}
R_{k}(n, q(k-1)) & \leq(q+1) R_{k}(n, 0)-q \\
& =(q+1) S_{k}(n)-q
\end{aligned}
$$

We may now settle Conjecture 1.2 in the case $n \geq k-1$.
Theorem 3.5. Assume $n \geq k-1$ and $c \geq 0$. Then $R_{k}(n, c)$ is finite if and only if $c$ is a multiple of $k-1$.

Proof. If $c$ is a multiple of $k-1$, the statement follows from the above proposition. On the other hand, if $c$ is not a multiple of $k-1$, then Proposition 2.2, with divisor $d=k-1$, implies $R_{k}(n, c)=+\infty$.
3.4. Further consequences. Our next consequence of Lemma 3.3 is a bound on $R_{k}(n, c)$ in terms of $R_{k}(n, 1)$.
Proposition 3.6. For all $c \geq 1$, we have $R_{k}(n, c) \leq c R_{k}(n, 1)$. In particular, if $R_{k}(n, 1)$ is finite, then $R_{k}(n, c)$ is finite for all $c \geq 0$.

Proof. Applying Lemma 3.3 with $c=1$ (the $c$ in that lemma again, not the current one) yields

$$
R_{k}(n, \alpha-\beta(k-1)) \leq \alpha R_{k}(n, 1)+\beta
$$

for all integer $\alpha \geq 1$ and $\beta \geq 1-\alpha$. Now, setting $\alpha=c$ (the current one) and $\beta=0$ in the above relation yields the stated inequality.

Finally, we find that if $c \geq k-1$, then $R_{k}(n, c)$ may be bounded below by $R_{k}(n, \bar{c})-q$, where $\bar{c}$ is the class of $c \bmod k-1$ and $q$ is the floor of $c /(k-1)$.
Proposition 3.7. Assume $c \geq k-1$, and let $c=q(k-1)+\bar{c}$ be the Euclidean division of $c$ by $k-1$, with $q \geq 0$ and $0 \leq \bar{c} \leq k-2$. Then

$$
R_{k}(n, \bar{c}) \leq R_{k}(n, c)+q
$$

Proof. It follows from Lemma 3.3, with values $\alpha=1$ and $\beta=q$, that

$$
R_{k}(n, \bar{c})=R_{k}(n, c-(k-1) q) \leq R_{k}(n, c)+q .
$$

3.5. The case $k \leq 7$. We now verify Conjecture 1.2 in case $k \leq 7$, except for the instances $R_{5}(3,2)$ and $R_{6}(4,1)$ which require completely different methods and will be presented elsewhere.

Only the "if" part of the conjecture remains open. It states that if every divisor $d \leq n$ of $k-1$ also divides $c$, then $R_{k}(n, c)$ should be finite. This is known to be true in the following cases:

- if $c=0$, in which case $R_{k}(n, 0)$ is finite by Rado's results [13].
- if $n=2$, by Schaal's results recalled in [19, Section 1.2).
- if $c \equiv 0 \bmod k-1$, in which case $R_{k}(n, c)$ is finite by Proposition 3.4.
- if $n \geq k-1$, by Theorem 3.5.

Consequently, for $k \leq 7$, it remains to verify the "if" part of the conjecture only for $n \in[3, k-2]$. In particular, Conjecture 1.2 holds for $k \leq 4$.

The case $k=5$. It remains to discuss the case $n=3$ and $c \not \equiv 0 \bmod 4$.

- For $c$ odd, there is a divisor $d \leq n$ of $k-1=4$ not dividing $c$, namely $d=2$. Therefore $R_{5}(3, c)=\infty$ in that case, by the "only if" part of the conjecture, i.e., by Proposition 2.2,
- For $c \equiv 2 \bmod 4:$ according to the conjecture, we must show that $R_{5}(3, c)$ is finite in this case. Now Lemma 3.3 reduces that statement to the sole finiteness of $R_{5}(3,2)$. Indeed, it implies $R_{5}(3,2 \alpha) \leq \alpha R_{5}(3,2)$ for all $\alpha \geq 1$, by setting $\beta=0$ and $c=2$ there. Finally, it turns out that $R_{5}(3,2)$ is indeed finite, as will be proved elsewhere.
We conclude that the conjecture holds for $k=5$.
The case $k=6$. We may assume $n \in[3,4]$ and $c \not \equiv 0 \bmod 5$. The only divisor $d \leq n$ of $k-1=5$ is 1 , which divides any $c$. Therefore, according to the conjecture, we must show that $R_{6}(n, c)$ is finite for $n \in[3,4]$ and all $c \geq 1$. Proposition 3.6 and the obvious bound $R_{k}(3, c) \leq R_{k}(4, c)$ reduce that statement to the sole finiteness of $R_{6}(4,1)$. Again, $R_{6}(4,1)$ turns out to be finite, with a proof similar to that for
$R_{5}(3,2)$ and to appear in the same paper. We conclude that the conjecture also holds for $k=6$.

The case $k=7$. We may assume $n \in[3,5]$ and $c \not \equiv 0 \bmod 6$. The only divisors $d \leq n$ of $k-1=6$ are 1,2 and 3 . If either 2 or 3 does not divide $c$, then $R_{7}(n, c)=\infty$ in these cases, by the "only if" part of the conjecture, i.e., by Proposition 2.2, Now, if both 2 and 3 divide $c$, then $c \equiv 0 \bmod 6$, an already settled case. Therefore, the conjecture holds for $k=7$, with a proof entirely contained in the present paper in contrast to the cases $k=5$ and 6 .

The conjecture remains open for $k \geq 8$. However, in order to settle the smallest open case $k=8$, it would suffice, by using the same reduction tools as above, to show that $R_{8}(6,1)$ is finite.
3.6. Yet another case of the conjecture. Here is yet another case where Conjecture 1.2 is shown to hold. Its interest lies in the fact that $n$ is smaller than $k-1$, in contrast with the situation in Section 3.3

Proposition 3.8. Let $n$ be an integer such that $\operatorname{gcd}(n, 6)=1$ and $n \geq 5$. Let $k=6 n+1$. The following statements are equivalent.
(1) $R_{k}(n, c)$ is finite.
(2) Every divisor $d \leq n$ of $k-1$ also divides $c$.
(3) $k-1$ divides $c$.

Proof. That (1) implies (2) directly follows from Proposition 2.2. Now assume (2). Applying this to $d=2, d=3$ and $d=n$, it follows that $2 \cdot 3 \cdot n$ divides $c$, since these numbers are pairwise coprime. That is, $k-1$ divides $c$, as stated in (3). Finally, that (3) implies (1) directly follows from Proposition 3.4

Needless to say, the same proof tools yield the same equivalences under this set of hypotheses: $\operatorname{gcd}(n, 30)=1, n \geq 5$ and $k=30 n+1$; or, for that matter, under the set of hypotheses $\operatorname{gcd}(n, 210)=1, n \geq 7$ and $k=210 n+1$; and so on.

## 4. Bounds and exact values

In this section, we define new numbers $S_{k}^{*}(n)$ which are smaller and easier to determine than $S_{k}(n)$. We then show how they provide a lower bound to $R_{k}(n, c)$, and we give an estimate for them. Finally, we use them to obtain a new shorter proof for Schaal's formula on $R_{2}(3, c)$, exact formulae for $R_{3}(3, c)$ and $R_{4}(3, c)$, and sharper bounds in some other cases.
4.1. The numbers $S_{k}^{*}(n)$. Before defining the $S_{k}^{*}(n)$ proper, we introduce and generalize some terminology from additive number theory.

Given nonempty sets $A_{1}, A_{2}$ of integers, their sumset $A_{1}+A_{2}$ is defined as

$$
A_{1}+A_{2}=\left\{x_{1}+x_{2} \mid x_{1} \in A_{1}, x_{2} \in A_{2}\right\} .
$$

In particular, if $A_{2}=\{c\}$ is a singleton, then $A_{1}+A_{2}=\left\{x+c \mid x \in A_{1}\right\}$; it is the translate of $A_{1}$ by $c$, and will be denoted by $A_{1}+c$ instead of $A_{1}+\{c\}$. If $A_{1}=A_{2}=A$, then we denote $2 A=A+A$, and more generally, for any integer $k \geq 1$, we denote by $k A$ the $k$-fold sumset of $A$ with itself, i.e.,

$$
k A=\underbrace{A+\cdots+A}_{k}=\left\{x_{1}+\cdots+x_{k} \mid x_{i} \in A \text { for all } 1 \leq i \leq k\right\} .
$$

A set $A$ of integers is said to be sum-free if $A \cap 2 A=\emptyset$. As a generalization well-suited to our purposes here, and with $k, c$ integers with $k \geq 1$, we shall say that $A$ is $(k X+c)$-free if

$$
(k A+c) \cap A=\emptyset .
$$

More generally, given a set $C$ of integers, we shall say that $A$ is $(k X+C)$-free if it is ( $k X+c$ )-free for all $c \in C$, or equivalently, if

$$
(k A+C) \cap A=\emptyset .
$$

Clearly, for $c \in \mathbb{Z}$, the set $A$ is $(k X+c)$-free if and only if it contains no solution to the equation

$$
\begin{equation*}
x_{1}+\cdots+x_{k}+c=x_{k+1} \tag{5}
\end{equation*}
$$

with $x_{1}, \ldots, x_{k+1} \in A$.
Thus, the $n$-color Rado number $R_{k}(n, c)$ may equivalently be described as the smallest integer $N$, if it exists, or $\infty$ if not, such that for every $n$-coloring of $[1, N]$, at least one of the color classes $A$ in $[1, N]$ fails to be $(k X+c)$-free.

We are now in a position to define the $S_{k}^{*}(n)$.
Definition 4.1. Let $n, k$ be positive integers with $k \geq 2$. Denote by $S_{k}^{*}(n)$ the least integer $N$ such that, for every $n$-coloring of $[1, N]$, there exists a monochromatic solution $\left(x_{1}, \ldots, x_{k+1}\right)$ to the equation

$$
\begin{equation*}
x_{1}+\cdots+x_{k}+s=x_{k+1} \tag{6}
\end{equation*}
$$

for some $s \in[-k+2,1]$.
Equivalently, let $I=[-k+2,1]$. Then $S_{k}^{*}(n)$ is the least integer $N$ such that, for every $n$-coloring of $[1, N]$, at least one of its color classes fails to be $(k X+I)$-free.

Indeed, a set $A$ is $(k X+I)$-free if and only if it is $(k X+s)$-free for all $s \in I$, or equivalently if, for all $s \in I$, it contains no solution to equation (6).

It directly follows from the definitions that

$$
S_{k}^{*}(n) \leq \min _{c \in[-k+2,1]} R_{k}(n, c) .
$$

In particular, at $c=0$, we have $S_{k}^{*}(n) \leq S_{k}(n)$.
4.2. Bounding $R_{k}(n, c)$ with $S_{k}^{*}(n)$. Our first result constructs a $(k X+c)$-free set $A$ from a $(k X+I)$-free set $B$, where $I=[-k+2,1]$. This will then be used to bound $R_{k}(n, c)$ in terms of $S_{k}^{*}(n)$.

Notation 4.2. For a subset $B \subseteq \mathbb{Z}$ and a positive integer $\lambda$, we denote

$$
\lambda \cdot B=\{\lambda y \mid y \in B\} .
$$

The dot here is important since it helps distinguish $\lambda \cdot B$ from the $\lambda$-fold sumset $\lambda B=B+\cdots+B$.

Lemma 4.3. Let $k, c$ be integers such that $k \geq 2$ and $c \geq 0$. Let $B \subseteq \mathbb{Z}$ be $a$ $(k X+I)$-free subset, where $I=[-k+2,1]$. Let

$$
A=(c+k-1) \cdot B+[-(c+k-2), 0] .
$$

Then $A$ is $(k X+c)$-free.

Proof. Assume that $A$ is not $(k X+c)$-free. Then there exist $x_{1}, \ldots, x_{k+1} \in A$ such that

$$
\begin{equation*}
x_{1}+\cdots+x_{k}+c=x_{k+1} . \tag{7}
\end{equation*}
$$

By construction, each $x_{i}$ decomposes as

$$
x_{i}=(c+k-1) y_{i}+r_{i}
$$

for some $y_{i} \in B$ and $r_{i} \in[-(c+k-2), 0]$. Equation (7) then yields

$$
(c+k-1)\left(y_{1}+\cdots+y_{k}-y_{k+1}\right)+\left(r_{1}+\cdots+r_{k}-r_{k+1}+c\right)=0 .
$$

It follows that $(c+k-1)$ divides $\left(r_{1}+\cdots+r_{k}-r_{k+1}+c\right)$, and that

$$
\begin{equation*}
\frac{\left(r_{1}+\cdots+r_{k}-r_{k+1}+c\right)}{(c+k-1)}=-\left(y_{1}+\cdots+y_{k}-y_{k+1}\right) \tag{8}
\end{equation*}
$$

Now, since $r_{i} \in[-(c+k-2), 0]$ for all $i \in[1, k+1]$, we have

$$
\begin{array}{rll}
-(c+k-2) k & \leq r_{1}+\cdots+r_{k} & \leq 0 \\
0 & \leq & -r_{k+1}
\end{array} \leq(c+k-2), ~ l
$$

from which it follows that

$$
-(c+k-2) k+c \leq\left(r_{1}+\cdots+r_{k}-r_{k+1}+c\right) \leq(c+k-2)+c .
$$

Dividing by $(c+k-1)$ and using (8), we get

$$
\begin{equation*}
\frac{-(c+k-2) k+c}{(c+k-1)} \leq-\left(y_{1}+\cdots+y_{k}-y_{k+1}\right) \leq \frac{(c+k-2)+c}{(c+k-1)} \tag{9}
\end{equation*}
$$

Since the middle term is an integer, inequalities (9) remain valid if we replace the leftmost term by its ceiling $\rceil$ and the rightmost one by its floor $\rfloor$.

The numerator $-(c+k-2) k+c$ in the leftmost term may be written as $-(c+k-1) k+(c+k-1)+1$, so that

$$
\frac{-(c+k-2) k+c}{(c+k-1)}=-k+1+\frac{1}{(c+k-1)},
$$

whose ceiling equals $-k+2$.
In turn, the numerator of the rightmost term of (9) may be written as $(c+k-$ 1) $+(c-1)$, so that

$$
\frac{(c+k-2)+c}{(c+k-1)}=1+\frac{(c-1)}{(c+k-1)},
$$

whose floor equals 0 if $c=0$, or 1 if $c \geq 1$. In either case, (9) yields

$$
-k+2 \leq-\left(y_{1}+\cdots+y_{k}-y_{k+1}\right) \leq 1 .
$$

Thus, setting $s=-\left(y_{1}+\cdots+y_{k}-y_{k+1}\right)$, we have $s \in[-k+2,1]$, and the equality

$$
y_{1}+\cdots+y_{k}+s=y_{k+1}
$$

implies that $B$ is not $(k X+s)$-free, contrary to the assumption.
This implies the following lower bound on $R_{k}(n, c)$ in terms of $S_{k}^{*}(n)$.
Proposition 4.4. Let $k, n \geq 2$. Then $R_{k}(n, c) \geq(c+k-1)\left(S_{k}^{*}(n)-1\right)+1$.

Proof. Set $M=S_{k}^{*}(n)$ and $N=(c+k-1)(M-1)+1$. By the minimality property of $S_{k}^{*}(n)$, there exists a special $n$-coloring $\Delta$ of $[1, M-1]$ all of whose color classes $B$ are $(k X+I)$-free, where $I=[-k+2,1]$. Let

$$
\pi:[1, N-1] \longrightarrow[1, M-1]
$$

be the map defined by

$$
\pi(x)=\left\lceil\frac{x}{(c+k-1)}\right\rceil
$$

for all $x \in[1, N-1]$, and let $\Delta^{\prime}$ be the $n$-coloring of $[1, N-1]$ defined by

$$
\Delta^{\prime}(x)=\Delta(\pi(x))
$$

for all $x \in[1, N-1]$. Clearly, for all $y \in[1, M-1]$ and all $r \in[-(c+k-2), 0]$, one has

$$
\pi((c+k-1) y+r)=y
$$

i.e., the map $\pi$ is constant on the subinterval $(c+k-1) y+[-(c+k-2), 0]$. It follows that each color class $A$ in $[1, N-1]$ under $\Delta^{\prime}$ is of the form

$$
\begin{equation*}
A=(c+k-1) \cdot B+[-(c+k-2), 0] \tag{10}
\end{equation*}
$$

for some color class $B$ in $[1, M-1]$ under $\Delta$. Now, since each such $B$ is $(k X+I)$ free, it follows from (10) and Lemma 4.3 that $A$ is $(k X+c)$-free. By the minimality property of $R_{k}(n, c)$, it follows that

$$
N-1 \leq R_{k}(n, c)-1
$$

as claimed.
Remark 4.5. Proposition 4.4 provides a lower bound on $R_{k}(n, c)$ involving $S_{k}^{*}(n)$, whereas Proposition 3.4 provides, if $c$ is a nonnegative multiple of $k-1$, an upper bound on $R_{k}(n, c)$ involving $S_{k}(n)$. Thus, combining both bounds for $c=q(k-1)$ with $q \in \mathbb{Z}_{+}$, we get

$$
(q+1)(k-1)\left(S_{k}^{*}(n)-1\right)+1 \leq R_{k}(n, c) \leq(q+1)\left(S_{k}(n)-1\right)+1
$$

Remarkably, it turns out that these two bounds sometimes coincide and hence yield the exact value of $R_{k}(n, c)$, as will be seen later on in some instances. At any rate, we get the following corollary.

Corollary 4.6. If $S_{k}(n)-1=(k-1)\left(S_{k}^{*}(n)-1\right)$, then for all $c=q(k-1)$ with $q \in \mathbb{N}$, we have $R_{k}(n, c)=(q+1)\left(S_{k}(n)-1\right)+1$.

Proof. This directly follows from the above remark.
4.3. A lower bound on $S_{k}^{*}(n)$. We first relate $S_{k}^{*}(n)$ to $S_{k}^{*}(n-1)$, and then derive an absolute lower bound on it.

Proposition 4.7. Let $k, n \geq 2$. Then $S_{k}^{*}(n) \geq(k+1) S_{k}^{*}(n-1)-k+1$.
Proof. Set $M=S_{k}^{*}(n-1)$ and $N=(k+1) M-k+1$. By the minimality property of $S_{k}^{*}(n-1)$, there exists a special $(n-1)$-coloring $\Delta$ of $[1, M-1]$ under which each color class $B \subseteq[1, M-1]$ is $(k X+s)$-free for all $s \in[-k+2,1]$.

Let us now extend $\Delta$ to the $n$-coloring

$$
\Delta^{\prime}:[1, N-1] \longrightarrow[1, n],
$$

defined as follows, for $x \in[1, N-1]$ :

$$
\Delta^{\prime}(x)=\left\{\begin{array}{cl}
\Delta(x) & \text { if } x \in[1, M-1],  \tag{11}\\
n & \text { if } x \in[M, k M-k+1], \\
\Delta(x-(N-1)) & \text { if } x \in[k M-k+2, N-1] .
\end{array}\right.
$$

Note, for later use, that the third interval above is just a translate of the first one, namely:

$$
\begin{equation*}
[k M-k+2, N-1]=[1, M-1]+(k M-k+1) \tag{12}
\end{equation*}
$$

It remains to show that each color class in $[1, N-1]$ under $\Delta^{\prime}$ is $(k X+s)$-free. The minimality property of $S_{k}^{*}(n)$ will then imply $N-1 \leq S_{k}^{*}(n)-1$, as desired.

Now, each color class $A$ under $\Delta^{\prime}$ is either equal to $[M, k M-k+1]$, or else it is of the form

$$
\begin{equation*}
A=B+\{0, k M-k+1\}=B+(k M-k+1) \cdot[0,1] \tag{13}
\end{equation*}
$$

for some color class $B \subseteq[1, M-1]$ under $\Delta$. This follows from (11) and (12).
So, let $A \subseteq[1, N-1]$ be a color class under $\Delta^{\prime}$, and let $s \in[-k+2,1]$. We now show that $A$ is $(k X+s)$-free.

Case 1. $A=[M, k M-k+1]$. We then have

$$
\min (k A+s)=k M+s \geq k M-k+2=\max (A)+1
$$

It follows that $(k A+s) \cap A=\emptyset$, as claimed.
Case 2. $A=B+(k M-k+1) \cdot[0,1]$ for some color class $B \subseteq[1, M-1]$, as stated in (13). Assume, for a contradiction, that $(k A+s) \cap A$ is not empty. Since

$$
k A+s=k B+s+(k M-k+1) \cdot[0, k]
$$

there exist $b_{1}, \ldots, b_{k+1} \in B$, and integers $u \in[0, k], v \in[0,1]$ such that

$$
b_{1}+\cdots+b_{k}+s+(k M-k+1) u=b_{k+1}+(k M-k+1) v
$$

It follows that $b_{1}+\cdots+b_{k}-b_{k+1}+s$ is a multiple of $(k M-k+1)$. Now, since $b_{i} \in[1, M-1]$ for all $i$ and since $s \in[-k+2,1]$, we have

$$
-M+3 \leq b_{1}+\cdots+b_{k}-b_{k+1}+s \leq k M-k .
$$

But the only multiple of $(k M-k+1)$ within this range is 0 . Therefore,

$$
b_{1}+\cdots+b_{k}-b_{k+1}+s=0
$$

i.e., $b_{1}+\cdots+b_{k}+s=b_{k+1}$ and hence belongs to $(k B+s) \cap B$. This contradicts the fact that the color class $B$ is $(k X+s)$-free.

Corollary 4.8. Let $k, n \geq 2$. Then $S_{k}^{*}(n) \geq \frac{(k+1)^{n}-1}{k}+1$.
Proof. Since $S_{k}^{*}(1)=2$, as easily seen, the inequality is satisfied for $n=1$. For general $n \geq 2$, we apply induction and Proposition 4.7.
4.4. Revisiting $R_{2}(3, c)$. An exact formula for $R_{2}(3, c)$ has been provided by Schaal [20]. We now provide a shorter proof for it, which exploits the above properties of the $S_{k}^{*}(n)$ and thereby avoids the case-by-case analysis of [20.

Proposition 4.9. We have $S_{2}^{*}(3)=14$.
Proof. Since $S_{2}(3)=14$ and $S_{2}^{*}(3) \leq S_{2}(3)$, we have $S_{2}^{*}(3) \leq 14$. The reverse inequality directly follows from Corollary 4.8. Alternatively, it suffices to exhibit a 3 -coloring of $[1,13]$ with all three color classes being $(2 X+I)$-free where $I=[0,1]$ as seen below:

$$
\begin{aligned}
& A_{1}=\{1,4,10,13\} \\
& A_{2}=\{2,3,11,12\} \\
& A_{3}=[5,9] .
\end{aligned}
$$

Each $A_{i}$ satisfies $\left(2 A_{i}+I\right) \cap A_{i}=\emptyset$, as readily checked and as required.
Corollary 4.10 (Schaal, [20]). We have $R_{2}(3, c)=13 c+14$ for all $c \geq 0$.
Proof. We have $S_{2}^{*}(3)=S_{2}(3)=14$. The first equality is an instance where the hypothesis

$$
S_{k}(n)=(k-1)\left(S_{k}^{*}(n)-1\right)+1
$$

of Corollary 4.6 is satisfied, here with $k=2$. That corollary then implies

$$
R_{2}(3, c)=(c+1)\left(S_{2}(3)-1\right)+1,
$$

i.e., $R_{2}(3, c)=13(c+1)+1=13 c+14$.

Using the same method of proof, it is easy to establish the corresponding formula for $R_{2}(2, c)$, namely:

$$
R_{2}(2, c)=4 c+5
$$

for all $c \geq 0$.
4.5. A formula for $S_{k}^{*}(2)$. We end this section by deriving a formula for $S_{k}^{*}(2)$.

Proposition 4.11. Let $k \geq 2$. Then $S_{k}^{*}(2)=k+3$.
Proof. The bound $S_{k}^{*}(2) \geq k+3$ directly follows from Corollary 4.8. To prove the reverse inequality, it suffices to show that for every 2-coloring of $[1, k+3]$, one of the two color classes fails to be $(k X+I)$-free, where as usual $I=[-k+2,1]$.

Given a 2 -coloring of $[1, k+3]$, let $A_{1}, A_{2}$ be its two color classes. We may freely assume that $1 \in A_{1}$. Since

$$
k\{1\}+I=\{k\}+[-k+2,1]=[2, k+1],
$$

it follows that if $A_{1} \cap[2, k+1]$ failed to be empty, then $A_{1}$ would fail to be $(k X+I)$ free and we would be done.

Therefore, we may assume $A_{1} \cap[2, k+1]=\emptyset$, i.e., $[2, k+1] \subseteq A_{2}$. Since

$$
k[2, k+1]+I=\left[2 k, k^{2}+k\right]+[-k+2,1]=\left[k+2, k^{2}+k+1\right],
$$

which contains $\{k+2, k+3\}$, we may assume that none of $k+2, k+3$ belongs to $A_{2}$, for otherwise $A_{2}$ would fail to be ( $k X+I$ )-free and we would again be done.

Therefore, we may assume $\{1, k+2, k+3\} \subseteq A_{1}$. But then, $A_{1}$ fails to be $(k X+I)$-free, since setting

$$
x_{1}=\cdots=x_{k-1}=1, \quad x_{k}=k+2, \quad x_{k+1}=k+3, \quad s=-k+2,
$$

we have $x_{i} \in A_{1}$ for all $i, s \in I$, and

$$
x_{1}+\cdots+x_{k}+s=x_{k+1} .
$$

## 5. Computer-aided results

The problem of computing $S_{k}^{*}(n)$ or $R_{k}(n, c)$ can be translated as a Boolean satisfiability problem, as detailed in Section 5.3 for the computation of $R_{2}(4, c)$. The resulting Boolean translation, for given instances, may then be fed to a computer running a suitable SAT solver. All results in this section have been obtained by combining such computations with the theory developed above. The specific SAT solver we have used is March RW [9, the gold medal winner of the 2011 International SAT Competition.

Proposition 5.1. The following values of $S_{k}^{*}(n)$, for $n=3$ and $2 \leq k \leq 6$, and for $n=4$ and $2 \leq k \leq 3$, hold:
(1) $S_{2}^{*}(3)=14, S_{3}^{*}(3)=22, S_{4}^{*}(3)=32, S_{5}^{*}(3)=44, S_{6}^{*}(3)=58$.
(2) $S_{2}^{*}(4)=41, S_{3}^{*}(4)=86$.

Proof. The formula for $S_{2}^{*}(3)$ has been established in Proposition 4.9, All others have been obtained by running a SAT solver on the corresponding Boolean translations.

The numbers $S_{k}(n)$ are larger and more difficult to compute than the $S_{k}^{*}(n)$ 's. However, in a few instances where we have been able to compute them, the hypothesis of Corollary 4.6. namely

$$
S_{k}(n)-1=(k-1)\left(S_{k}^{*}(n)-1\right),
$$

turned out to be satisfied, thereby allowing an exact formula for the corresponding $R_{k}(n, c)$ 's.
5.1. Exact formulas for $R_{3}(3, c)$ and $R_{4}(3, c)$. We now establish previously unknown formulas for $R_{3}(3, c)$ and $R_{4}(3, c)$, as new applications of Corollary 4.6,
5.1.1. The case $R_{3}(3, c)$. We first need the value of $S_{3}(3)$, recently obtained in 18.

Theorem 5.2. We have $S_{3}(3)=43$.
Proof. The inequality $S_{3}(3) \leq 43$ is obtained by computer using the SAT solver March [9. Indeed, the solver established that the Boolean constraints derived from assuming the existence of a 3 -coloring of $[1,43]$ with $(3 X+0)$-free color classes cannot be satisfied.

The reverse inequality follows from a result of Znam [24] which, for $k=n=3$, yields the bound

$$
S_{3}(3) \geq \frac{2}{3}\left(4^{3}-1\right)+1=43
$$

Theorem 5.3. For every integer $c \geq 0$, we have

$$
R_{3}(3, c)= \begin{cases}\infty & \text { if c odd } \\ 21 c+43 & \text { if c even }\end{cases}
$$

Proof. For $c$ odd, we have $R_{3}(3, c)=+\infty$ by Proposition 2.2. Now assume $c=2 q$ with $q \in \mathbb{N}$. Since $S_{3}^{*}(3)=22$ by Proposition 5.1, and since $S_{3}(3)=43$ as stated above, we see that the hypothesis

$$
S_{k}(n)-1=(k-1)\left(S_{k}^{*}(n)-1\right)
$$

of Corollary 4.6 is again satisfied in this instance. Therefore, that corollary yields the formula

$$
R_{3}(3,2 q)=(q+1)\left(S_{3}(3)-1\right)+1
$$

that is, for $c=2 q: R_{3}(3, c)=42(q+1)+1=21 c+43$, as claimed.
5.1.2. The case $R_{4}(3, c)$. In view of applying Corollary 4.6 again, we now need the value of $S_{4}(3)$.
Theorem 5.4. We have $S_{4}(3)=94$.
Proof. Follows from a three-hour computation with the SAT solver March, using our Boolean translation of the problem along the lines of Section 5.3.

Theorem 5.5. For every integer $c \geq 0$, we have

$$
R_{4}(3, c)= \begin{cases}\infty & \text { if } c \notin 3 \mathbb{N} \\ 31 c+94 & \text { if } c \in 3 \mathbb{N}\end{cases}
$$

Proof. For $c$ not divisible by 3 , we have $R_{4}(3, c)=+\infty$ by Proposition 2.2, Assume now $c=3 q$ with $q \in \mathbb{N}$. Since $S_{4}^{*}(3)=32$ by Proposition 5.1 and since $S_{4}(3)=94$ as stated above, we see that the hypothesis

$$
S_{k}(n)-1=(k-1)\left(S_{k}^{*}(n)-1\right)
$$

of Corollary 4.6 is again satisfied in this instance. Therefore, that corollary yields the formula

$$
R_{4}(3,3 q)=(q+1)\left(S_{4}(3)-1\right)+1
$$

that is, for $c=3 q: R_{4}(3, c)=93(q+1)+1=31 c+94$, as claimed.
5.2. Some new bounds. Short of exact formulae, Remark 4.5 also enables us to obtain new bounds on suitable instances of $R_{k}(n, c)$.

Corollary 5.6. Let $c \in \mathbb{N}$. Then $40 c+41 \leq R_{2}(4, c) \leq 44 c+45$.
Proof. The lower bound on $R_{2}(4, c)$ follows from Proposition 4.4 and the value $S_{2}^{*}(4)=41$ stated in Proposition 5.1 whereas the upper bound follows from Proposition 3.4 and the known value $S_{2}(4)=45$.

Corollary 5.7. Let $c \in \mathbb{N}$. Then $121 c+122 \leq R_{2}(5, c) \leq 305 c+306$.
Proof. It is known that $S_{2}(5) \leq 316$; see [14, 17]. Moreover, Radziszowski showed in [14] that $r_{5}(3) \leq 307$, where $r_{5}(3)=r(3,3,3,3,3)$ denotes the 5 -color Ramsey number for unavoidable monochromatic triangles in any edge-colored complete graph of that order. Applying the relationship between Schur numbers and Ramsey numbers given by Roberston [16], one obtains $S_{2}(5) \leq r_{5}(3)-1 \leq 306$.

Propositions 3.4 and 4.4 then yield, for $k=2$ and $n=5$, the stated bounds on $R_{2}(5, c)$.
5.3. Seeking $R_{2}(4, c)$ by computer. Having fixed $c \geq 0$, we seek successive integers $M$ such that $[1, M]$ admits a 4 -coloring all of whose color classes are $(2 X+c)$-free; or equivalently, such that no triplet of the form $\{i, j, i+j+c\}$, with $1 \leq i \leq M-c$ and $i \leq j \leq M-i-c$, is monochromatic. When the largest possible such $M$ is found, we are done: $R_{2}(4, c)=M+1$.

We now reformulate this problem as a Boolean satisfiability problem [6]. We proceed as follows.

First, any 4-coloring of $[1, M]$ may be viewed as a function

$$
\Delta:[1, M] \longrightarrow\{0,1\}^{2}
$$

By setting $\Delta(i)=\left(x_{i}, x_{i+M}\right)$ for all $i \in[1, M]$, this 4 -coloring may be represented by $2 M$ binary variables $x_{1}, \ldots, x_{2 M}$ with values in $\{0,1\}$.

We now view 0 and 1 as representing the Boolean values False and True, respectively. This allows us to use the logical operators AND, OR and NOT, denoted respectively by $\wedge, \vee$ and $\neg$, on the set $\{0,1\}$, with the purpose of translating the nonequality of elements by the validity of some associated logical formula. Indeed, for any $x, y \in\{0,1\}$, we have

$$
\begin{equation*}
x \neq y \Longleftrightarrow(x \vee y) \wedge(\neg x \vee \neg y) \text { is True, } \tag{14}
\end{equation*}
$$

as readily checked. It is easy to similarly translate the non-equality of two given colors in $\{0,1\}^{2}$, since for $\left(x_{1}, y_{1}\right),\left(x_{2}, y_{2}\right) \in\{0,1\}^{2}$, we have

$$
\left(x_{1}, y_{1}\right) \neq\left(x_{2}, y_{2}\right) \Longleftrightarrow x_{1} \neq x_{2} \text { or } y_{1} \neq y_{2}
$$

Let us go back to our generic 4-coloring

$$
\Delta:[1, M] \longrightarrow\{0,1\}^{2}
$$

represented by the Boolean variables $x_{1}, \ldots, x_{2 M}$. Given an arbitrary subset $A \subseteq$ $[1, M]$, we may associate to $A$, in the way described above, a logical formula $\lambda(A)$ in the variables $x_{1}, \ldots, x_{2 M}$, in such a way that $A$ fails to be monochromatic if and only if $\lambda(A)$ is True. Therefore, given a family of subsets $A_{1}, \ldots, A_{r} \subseteq[1, M]$, prohibiting all of them to be monochromatic admits the following Boolean translation:

$$
\begin{equation*}
\text { no } A_{i} \text { is monochromatic } \Longleftrightarrow \bigwedge_{i=1}^{r} \lambda\left(A_{i}\right) \text { is True. } \tag{15}
\end{equation*}
$$

Now, applying this translation to the above-mentioned set of triplets, namely the subsets $\{i, j, i+j+c\} \subseteq[1, M]$ with $1 \leq i \leq M-c$ and $i \leq j \leq M-i-c$, we obtain by (15) a system of logical formulas that can be simultaneously satisfied if and only if $[1, M]$ admits a 4 -coloring all of whose color classes are $(2 X+c)$-free. Fixing successive values of $M$, we feed the associated system to a SAT solver, which will then attempt to say whether that system is satisfiable or not. As long as it is, we increase $M$. When, for some $M_{0}$, we reach nonsatisfiability as an output, we know we are done: $R_{2}(4, c)=M_{0}$.
5.4. Exact values of $R_{2}(4, c)$ for $c \leq 6$. Applying the above method, we have obtained, by computer, the following exact value of the Rado numbers $R_{2}(4, c)$ for $0 \leq c \leq 6$.
Theorem 5.8. We have

$$
\begin{aligned}
& R_{2}(4,0)=40 \cdot 0+45=45 \\
& R_{2}(4,1)=40 \cdot 1+43=83, \text { and } \\
& R_{2}(4, c)=40 \cdot c+41 \quad \text { for } 2 \leq c \leq 6
\end{aligned}
$$

Proof. The value $R_{2}(4,0)=45$ is due to [2]. A 4-coloring of $[1,82]$ with $(2 X+1)$-free color classes, implying $R_{2}(4,1) \geq 83$, is given by the following 4 color classes:

$$
\begin{aligned}
& \{1,2,15,16,21,22,28,29,35,36,47,54,55,61,62,67,68,81,82\}, \\
& \{3,4,5,6,17,18,19,20,33,34,49,50,63,64,65,66,77,78,79,80\} \\
& \{23,24,25,26,27,37,38,39,40,41,42,43,44,45,46,56,57,58,59,60\}, \\
& \{7,8,9,10,11,12,13,14,30,31,32,48,51,52,53,69,70,71,72,73,74,75,76\}
\end{aligned}
$$

thereby improving the lower bound $R_{2}(4,1) \geq 81$ given by Proposition 4.4 and the value $S_{2}^{*}(4)=41$ of Proposition 5.1 For $2 \leq c \leq 6$, the lower bound is achieved by Proposition 4.4, and is revealed to be sharp by computations using the SAT solver March [9. The running times on a standard desktop computer were as follows:

| Values of c | $R_{2}(4, c)$ | Time in seconds |
| :---: | ---: | ---: |
| $c=1$ | 83 | 13 |
| $c=2$ | 121 | 50 |
| $c=3$ | 161 | 1260 |
| $c=4$ | 201 | 2810 |
| $c=5$ | 241 | 9270 |
| $c=6$ | 281 | 593000 |

5.5. Conclusions and open problems. Combined with a separate forthcoming paper establishing the finiteness of $R_{5}(3,2)$ and $R_{6}(4,1)$, Conjecture 1.2 turns out to be true for $k \leq 7$. It remains to settle it in general. The smallest open case is $R_{8}(6,1)$, conjectured to be finite. In addition, in view of the above determination of $R_{2}(4, c)$ for $c \leq 6$, it is natural to conjecture that the formula $R_{2}(4, c)=40 c+41$ also holds for $c \geq 7$. Is it true or not? Along the same line, we are also interested in determining exact values or sharper bounds for $R_{2}(5, c)$.

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