Dominators for Multiple-objective Quasiconvex Maximization Problems

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Abstract. In this paper we address the problem of finding a dominator for a multiple-objective maximization problem with quasiconvex functions. The one-dimensional case is discussed in some detail, showing how a Branch-and-Bound procedure leads to a dominator with certain minimality properties. Then, the well-known result stating that the set of vertices of a polytope S contains an optimal solution for single-objective quasiconvex maximization problems is extended to multiple-objective problems, showing that, under upper-semicontinuity assumptions, the set of (k-1)-dimensional faces is a dominator for k-objective problems. In particular, for biobjective quasiconvex problems on a polytope S, the edges of S constitute a dominator, from which a dominator with minimality properties can be extracted by Branch-and Bound methods.

Key words: Multiple-objective problems; Quasiconvex maximization; Dominators

1. Introduction

Given a nonempty closed subset S of \mathbb{R}^n and a function $F: S \subset \mathbb{R}^n \to \mathbb{R}^k$, define the *multiple-objective* problem (P[F; S]),

$$\max_{x \in S} F(x), \qquad (P[F; S])$$

which seeks those alternatives maximizing simultaneously the components F_1, F_2, \ldots, F_k of F, [7, 28, 31].

Although the term simultaneous maximization is not uniquely defined, it customarily means finding the set $\mathscr{E}[F;S]$ of *efficient* or *Pareto-optimal* solutions to (P[F;S]),

$$\mathscr{E}[F;S] = \{x \in S : \text{no } y \in S \text{ verifies } F_i(y) \ge F_i(x) \ \forall i = 1, 2, \dots, k$$
 with at least one inequality strict\}

In general $\mathscr{E}[F;S]$ lacks many desirable properties such as being connected or closed, and this seems to be quite often the case and not only in pathological

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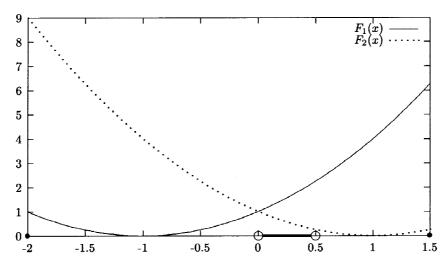


Figure 1. Biobjective convex maximization.

examples: take, for instance, the biobjective convex maximization problem in one variable (n = 1, k = 2) with $F(x) = ((x + 1)^2, (x - 1)^2)$ and S = [-2, 1.5], plotted in Figure 1.

Since $F(-2) \ge F(x) \ \forall x \in]-2,0]$, with at least one inequality strict, and $F(1.5) \ge F(x) \ \forall x \in [0.5, 1.5[$, with at least one ineuality strict too, it follows that the set of Pareto-optimal points must be contained in $\{-2\} \cup [0, 0.5[\cup \{1.5\}]$. In fact, it is readily seen from the plot that

$$\mathscr{E}[F; S] = \{-2\} \cup [0, 0.5] \cup \{1.5\},$$

which is a disconnected non-closed set. See following sections and also e.g. [3] for other instances.

Moreover, although there exist procedures to check whether a given point is efficient or not, e.g. [7, 31], an algorithm to construct $\mathscr{E}[F;S]$ is only available for a few classes of problems, such as multiple-objective linear problems, [28].

This drawback has been overcome in the literature by means of two strategies: either $\mathscr{E}[F;S]$ is sought, but, due to the unability for obtaining it, an approximation (sometimes with unknown degree of precision) is provided, e.g. [8, 18], or else the concept of efficiency is relaxed and replaced by a manageable surrogate of it.

In this paper we follow the second approach by using the concept of *dominator*, [5, 16, 21, 30], also called *weak kernel*, e.g. in [31] which is defined as any subset $S_0 \subset S$ such that, for any feasible $x \notin S_0$, S_0 contains a feasible alternative at least as good as x with respect to all objectives. See Section 2 for a formal definition.

It should be remarked that this concept is not only useful as a surrogate of the idea of Pareto-efficiency, but also as a tool in the resolution of some single-objective problems. Indeed, some of the most popular optimization methods for single-objective problems of the form

$$\max_{x \in S} \Psi(x) \tag{1.1}$$

require the feasible region S to be bounded. Such is the case, among others, of the Branch and Bound methods for global optimization, e.g. [15], which, in their simplest version, require, as pre-processing, the construction of a bounded polyhedron P (usually a hyper-rectangle, or a simplex) including either the whole feasible region, or, at least a bounded subset $S_0 \subset S$ known to contain an optimal solution. Moreover, the speed of convergence of the procedure is known to deteriorate with the volume of P, so P should be as small as possible in order to obtain reasonable computation times.

How to construct P will depend, of course, on the specific properties of the problem at hand. In particular, if (1.1) has the form

$$\max_{x \in S} \Phi(F(x)), \tag{1.2}$$

for some $\Phi: F(S) \to \mathbb{R}$ componentwise non-decreasing, then it is well known that, if (1.1) has optimal solutions, then any dominator for the multiple-objective problem $\max_{x \in S} F(x)$ also contains optimal solutions for (1.1), [21]. In other words, we can take as S_0 any bounded dominator for the multiple-objective problem, and as P any superset of S_0 with the required geometry.

This property has been successfully exploited, among others, in [5, 21, 22, 30] for problems of Linear Regression and Continuous Location, in which the globalizing function Φ is an arbitrary non-decreasing function and the function F is componentwise concave. Our aim here is to address the (harder) problem in which the function F is componentwise (quasi)-convex, showing as main result (Proposition 19) that, under upper-semicontinuity assumptions, the search of a dominator can be restricted to the (k-1)-dimensional faces of S.

The rest of this paper is structured as follows. In Section 2 we formally introduce the concept of dominators and discuss some general properties. These properties are used in Section 3 to address the one-dimensional case, for which dominators with certain minimality properties can be obtained.

Section 4 is devoted to show that, for multiple-objective multi-dimensional problems, one can construct dominators contained in low dimensional faces of the polytope S.

The paper ends with an application of these results to the construction of a dominator for a biobjective problem in Continuous Location. The reader is referred also to [25] for another successful application of the technique developed in this paper.

2. Dominators

Defining for each $x \in S$ the upper level set at x of F on S, $\mathcal{S}^{\geq}(x)$ as

$$\mathcal{S}^{\geqslant}(x) = \{ y \in S : F_i(y) \geqslant F_i(x) \text{ for all } i = 1, 2, \dots, k \},$$

the set $\mathscr{E}[F;S]$ of efficient solutions may be defined by

$$\mathscr{E}[F; S] = \{ x \in S : \text{If } y \in \mathscr{S}^{\geqslant}(x) \text{ then } x \in \mathscr{S}^{\geqslant}(y) \}$$
$$= \{ x \in S : \text{If } y \in \mathscr{S}^{\geqslant}(x) \text{ then } F(x) = F(y) \}$$

DEFINITION 1. A set $S^* \subset S$ is said to be a dominator for (P[F; S]) iff for each $x \in S$ there exists some $x^* \in S^*$ which has, componentwise, a value not smaller than x. In other words, S^* is a dominator iff

$$(\forall x \in S) \exists x^* \in \mathcal{S}^{\geq}(x) \cap S^*$$

Hereafter, the class of dominators for (P[F; S]) will be denoted by $\mathcal{D}[F; S]$. A direct consequence of the definition is the following:

PROPOSITION 2. One has

- 1. $S \in \mathcal{D}[F; S]$. In particular, $\mathcal{D}[F; S]$ is nonempty.
- 2. If $D \in \mathcal{D}[F; S]$ and D^* satisfies $D \subset D^* \subset S$, then $D^* \in \mathcal{D}[F; S]$.
- 3. For any class $\{S_i : j \in J\}$ of nonempty sets in \mathbb{R}^n ,

If
$$S_j^* \in \mathcal{D}[F; S_j](\forall j \in J)$$
 then $\bigcup_{j \in J} S_j^* \in \mathcal{D}\left[F; \bigcup_{j \in J} S_j\right]$

4. For any class $\{S_i : j \in J\}$ of nonempty sets in \mathbb{R}^n ,

$$\bigcap_{j \in J} \mathscr{D}[F; S_j] \subset \mathscr{D}\left[F; \bigcup_{j \in J} S_j\right]$$

5. If $D \in \mathcal{D}[F; S]$, then

$$\mathscr{D}[F;D] \subset \mathscr{D}[F;S]$$
.

By Proposition 2, the class $\mathcal{D}[F;S]$ is nonempty since the whole feasible set S is one of its elements. However S does not seem to be the most appropriate dominator since it possibly contains (too) many dominated alternatives, being too far from the ideal aim of a smallest possible dominator.

PROPOSITION 3. Suppose each F_j is upper-semicontinuous on S, then any class of compact nested dominators is closed under intersections. In other words: if (I, \leq) is a totally ordered set, and $\{D_i\}_{i\in I}$ is a class of compact dominators with $D_i \subset D_j$, $j \in I$, $i \leq j$, then

$$\bigcap_{i\in I} D_i \in \mathscr{D}[F;S].$$

Proof. Take any $x \in S$. By the upper-semicontinuity of the functions F_j , all upper level sets $\{y \in S : F_j(y) \ge F_j(x)\}$ are closed, so their intersection $\mathcal{S} \ge (x)$ is also closed. By the definition of dominators and their compactness, it follows for each $i \in I$ that $\mathcal{S}^\ge(x) \cap D_i$ is a nonempty compact set, thus $\{\mathcal{S}^\ge(x) \cap D_i\}_{i \in I}$ constitutes a

class of nested compact sets. By compactness their intersection (i.e., $\mathscr{S}^{\geqslant}(x) \cap \bigcap_{i \in I} D_i$) is nonempty.

However, it is evident that the whole class $\mathscr{D}[F;S]$ is not closed under intersections (take constant functions F_1,\ldots,F_k , then any singletons $\{x\},\{y\}\subset S$ are dominators, with empty intersection). Hence, a unique smallest dominator is unlikely to exist. We then relax the idea of smallest dominator by introducing the concept of (weak) minimal dominators. First define for each $x\in S$ the strict upper level set of F on S, $\mathscr{S}^>(x)$ as

$$\mathcal{S}^{>}(x) = \{ y \in S : F_i(y) > F_i(x) \text{ for all } i = 1, 2, ..., k \}.$$

DEFINITION 4. A dominator S^* is said to be minimal for (P[F; S]) iff no proper subset of S^* belongs to $\mathcal{D}[F; S]$. In other words, $S^* \subset S$ is minimal iff

$$(x, y \in S^*, x \neq y) \Rightarrow x \notin \mathcal{S}^{\geqslant}(y)$$

A dominator $S^* \subset S$ is said to be weak minimal for (P[F; S]) iff

$$(x, y \in S^*) \Rightarrow x \not\in \mathcal{S}^>(y)$$

The class of minimal (respectively weak minimal) dominators for problem (P[F;S]) will be denoted by $\mathcal{D}_M[F;S]$ (respectively $\mathcal{D}_{WM}[F;S]$).

As a simple illustration of the concepts, consider the 2-dimensional 2-objective optimization problem $\max_{x \in S} F(x)$, depicted in Figure 2, where the feasible region S is the polyhedron in \mathbb{R}^2 with vertices a = (0, -3), b = (4, -1), c = (4, 0), d = (0, 3), and F is given by

$$F_1(x_1, x_2) = x_1$$

$$F_2(x_1, x_2) = |x_2|$$

Then, the Pareto optimal set is given by

$$\mathscr{E}[F;S] = \{d\} \cup [a,b],$$

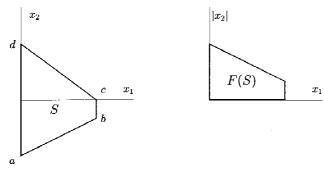


Figure 2. S and F(S).

only two minimal dominators exist, namely

$$S_1 = [a, b]$$

$$S_2 = [a, b] \cup \{d\},$$

whereas the polygonal S_3 ,

$$S_3 = \{d\} \cup [a, b] \cup [b, c]$$

is also weak minimal.

We observe in this example that the two minimal dominators are proper subsets of $\mathscr{E}[F;S]$. This result is more general, as stated in the following:

PROPOSITION 5. Suppose that S is compact and each F_i is upper semicontinuous on S. Then

- 1. $\mathscr{E}[F;S]$ is a weak minimal dominator.
- 2. Minimal dominators exist.
- 3. $\mathscr{E}[F;S] = \bigcup_{S^* \in \mathscr{D}_M[F;S]} S^*$.

Proof. By the upper-semicontinuity assumption, for each $x \in S$ the set $\mathscr{S}^{\geqslant}(x)$ is compact. Hence, by Theorem 6 of Chapter 2 of [31] $\mathscr{E}[F;S]$ is a dominator, which, by construction, is also weak minimal. Hence 1 holds.

To show 2, define on $\mathscr{E}[F,S]$ the equivalence relation

$$\rho = \{(x, y) \in \mathcal{E}[F; S] \times \mathcal{E}[F; S] : F(x) = F(y)\}.$$

Taking exactly one element in every equivalence class, we obtain a set S^* which is, by construction, a minimal dominator. Indeed, it is a dominator because $\mathscr{E}[F;S]$ is a dominator, as shown in Part 1. Moreover it is minimal: if there exists some dominator $M \subset S^*$, $M \neq S^*$, for any $x \in S^* | M$ there would exist some $y \in M$ with $F(y) \ge F(x)$. But by construction of S^* we would have $F(y) \ne F(x)$ contradicting the fact that x is efficient. Hence, minimal dominators exist.

For Part 3, we first show that every efficient point is in some minimal dominator: let $x^* \in \mathcal{E}[F;S]$, and construct a subset S^* of $\mathcal{E}[F;S]$ taking exactly one element of every equivalence class (with respect to the equivalence relation ρ above), x^* being the element chosen from its equivalence class. Using the reasoning above, it is seen that S^* is a minimal dominator, and $x^* \in S^*$.

Finally to show that any minimal dominator is included in the efficient set, take $x^* \in S^*$, for some $S^* \in \mathcal{D}_M[F;S]$, and assume $x^* \notin \mathcal{E}[F;S]$. Then, there exists some $y \in S$ with $F(y) \ge F(x)$, and at least one inequality strict. Since $S^* \in \mathcal{D}_M[F;S]$, there must exist some $y^* \in S^*$ with $F(y^*) \ge F(y) \ge F(x^*)$, thus the set $S^* \setminus \{x^*\}$ will also be a dominator, contradicting the minimality of S^* . Hence, $x^* \in \mathcal{E}[F;S]$. \square

REMARK 6. The upper-semicontinuity assumption is needed in order to guarantee the nonvoidness of $\mathcal{D}_{WM}[F;S]$, as the following counterexample shows: Let $S \subset \mathbb{R}^2$ be the triangle whose endpoints are (-1,0), (1,0), (0,1), and let $F_1:S \to \mathbb{R}$ be

defined as $1/(1-x_2)$ on the relative interior of the two top-edges, and zero elsewhere. Since

$$\lim_{\substack{(x_1,x_2)\to(0,\,1),\\(x_1,x_2)\in bd(S)}} F_1(x_1,x_2) = +\infty\,,$$

the maximum of F_1 on S is not attained, thus any $D \in \mathcal{D}[F_1; S]$ must contain a sequence of boundary points converging to (0, 1), implying that D contains points x, y with $F_1(x) > F_1(y)$. Hence, no weak minimal dominator exists.

3. Multiple-objective one-dimensional problems

In this section we address the multiple-objective problem (P[F; S]) when S is given as a finite union of compact intervals in \mathbb{R} , and each F_i is quasiconvex on each interval. We first discuss some properties of one-dimensional single-objective quasiconvex minimization problems, which are then used to tackle (P[F; S]), first when S reduces to a single compact interval and then in the general case. For the basic properties of quasiconvex functions we refer the reader to [1].

3.1. SINGLE-OBJECTIVE QUASICONVEX MINIMIZATION PROBLEMS ON AN INTERVAL

Let $I \subset \mathbb{R}$ be a nonempty compact interval, and let $g: I \to \mathbb{R}$ be quasiconvex. We will denote by $\operatorname{cl}_{I} g$ the closure of g relative to I, namely

$$\operatorname{cl}_{I} g(x) = \inf \{ t : \exists \{x_r\}_r \subset I, \text{ such that } x_r \to x, \ g(x_r) \to t \}$$

$$= \lim_{x_r \to x} \inf g(x_r)$$
(3.3)

LEMMA 7. One has:

- 1. $g(x) \ge \operatorname{cl}_1 g(x)$ for all $x \in I$.
- 2. $\inf_{x \in I} g(x) = \inf_{x \in I} \operatorname{cl}_{I} g(x)$.
- 3. cl₁ g is quasiconvex and lower-semicontinuous.
- 4. The set $\arg\min_{x\in I}\operatorname{cl}_{\mathrm{I}}g(x)$ of optimal solutions to $\min_{x\in I}\operatorname{cl}_{\mathrm{I}}g(x)$ is a nonempty compact subinterval of I.

Proof. 1 to 3 immediately follow from the definition of quasiconvexity and (3.3). By the lower semicontinuity of $\operatorname{cl}_{\operatorname{I}} g$, the set $\operatorname{arg\,min}_{x\in I} g(x)$ is compact and nonempty; since $\operatorname{cl}_{\operatorname{I}} g$ is also quasiconvex, it follows that $\operatorname{arg\,min}_{x\in I} \operatorname{cl}_{\operatorname{I}} g(x)$ is also convex, thus it is a compact interval, and Part 4 follows.

We recall that a function g is said to be *semistrictly quasiconvex*, [1], if it satisfies the following:

$$g(a) < g(b) c \in]a, b[$$
 \Rightarrow $g(c) < g(b)$

The next lemma shows that, due to the quasiconvexity of g, the behavior of g and

 $\operatorname{cl}_{\mathrm{I}} g$ are closely related, the relationship being stronger for semistrictly quasiconvex g:

LEMMA 8. Let $x^* \in \arg\min_{x \in I} \operatorname{cl}_{\mathrm{I}} g(x)$, and let $z_1, z_2 \in I$ such that $z_1 \in]x^*, z_2[$. One has:

- 1. $g(z_1) \leq g(z_2)$.
- 2. If g is also semistrictly quasiconvex and $g(z_1) = g(z_2)$, then $]x^*, z_2[\subset \arg\min_{x \in I} g(x).$

Proof. By definition of $\operatorname{cl}_{\mathrm{I}} g$ and Part 2 of Lemma 7, one can take a sequence $\{x_r\}$ in I converging to x^* such that $\inf_{r} g(x_r) = \inf_{x \in I} g(x) = \operatorname{cl}_{\mathrm{I}} g(x^*)$.

Since $z_1 > x^*$, there exists r_0 such that $x_r < z_1$ for all $r \ge r_0$, thus

$$z_1 \in]x_r, z_2[$$
 for all $r \ge r_0$

Given $r \ge r_0$, if it were the case that $g(z_2) < g(z_1)$, then

$$g(z_2) < g(z_1)$$

$$\leq \max\{g(z_2), g(x_r)\}$$

Hence, $g(z_1) \le g(x_r)$ for each $r \ge r_0$ thus one would have

$$g(z_2) < g(z_1)$$

$$\leq \inf_r g(x_r)$$

$$= \inf_{r \in I} g(x),$$

which is a contradiction. Hence, $g(z_2) \ge g(z_1)$, which shows 1.

To show 2, by the quasiconvexity of g it is enough to show that, if $g(z_1) = g(z_2)$, then $\{z_1, z_2\} \subset \arg\min_{x \in I} g(x)$. Suppose that, on the contrary, $g(z_1) = g(z_2) > \inf_{x \in I} g(x)$. Then, by Lemma 7,

$$g(z_1) = g(z_2)$$

$$> \operatorname{cl}_1 g(x^*),$$

and we could take a sequence $\{x_r\}$ converging to x^* with $g(x_r)$ converging to $\operatorname{cl}_1 g(x^*)$ and $g(x_r) < g(z_2)$ for each r. Since $z_1 \in]x^*, z_2[$, it would follow that $z_1 \in]x_r, z_2[$ for some r, thus, by the strict quasiconvexity of g, $g(z_1) < g(z_2)$, which would be a contradiction. Hence $g(z_1) = g(z_2) = \operatorname{cl}_1 g(x^*)$, showing that

$$[z_1, z_2] \subset \arg\min_{x \in I} g(x)$$
.

By the quasiconvexity of both g and $\operatorname{cl}_1 g$, and the optimality of x^* and $[z_1, z_2]$ for $\min_{x \in I} \operatorname{cl}_1 g(x)$, it then follows that

$$[x^*, z_1] \subset \arg\min_{x \in I} g(x)$$
,

and the result holds.

Another interesting property, which will be exploited in the sequel, states that, once problem $\min_{x \in I} \operatorname{cl}_I g(x)$ has been solved, any problem $\inf_{x \in J} g(x)$ with nested feasible interval $J \subset I$ is immediately solved. Indeed, denoting by $\operatorname{i}(J)$ the interior of J, one has:

PROPOSITION 9. Let $J := [a, b] \subset I$ be two compact intervals in \mathbb{R} . One has:

1.
$$\operatorname{cl}_{\operatorname{I}} g \leq \operatorname{cl}_{\operatorname{J}} g$$
 on J , and

$$\operatorname{cl}_{\mathbf{I}} g(x) = \operatorname{cl}_{\mathbf{I}} g(x) \quad \text{for all } x \in \mathrm{i}(J)$$
(3.4)

2. If $(\arg\min_{x\in I}\operatorname{cl}_{I}g(x))\cap\operatorname{i}(J)\neq\emptyset$, then

$$\inf_{x \in J} g(x) = \min_{x \in I} \operatorname{cl}_{I} g(x) \tag{3.5}$$

3. If $(\arg\min_{x\in I}\operatorname{cl}_{I}g(x))\cap\operatorname{i}(J)=\emptyset$, then

$$\inf_{x \in J} g(x) = \min\{g(a), g(b)\}$$
 (3.6)

Proof. Part 1 is a direct consequence of the definition of the closure of g and Lemma 7.

For Part 2, let $x^* \in \arg\min_{x \in I} \operatorname{cl}_{I} g(x) \cap \operatorname{i}(J)$; then, by Parts 1, 2 of Lemma 7 and Part 1 of this proposition,

$$\min_{x \in I} \operatorname{cl}_{I} g(x) = \operatorname{cl}_{I} g(x^{*})$$

$$= \operatorname{cl}_{J} g(x^{*})$$

$$= \min_{x \in J} \operatorname{cl}_{J} g(x)$$

$$= \inf_{x \in I} g(x)$$

$$\geq \inf_{x \in I} g(x)$$

$$= \min_{x \in I} \operatorname{cl}_{I} g(x)$$

Part 3 immediately follows from Lemma 8 if $\arg\min_{x\in I}\operatorname{cl}_{1}g(x)$ contains points in $I\setminus J$. In the remaining case, $\arg\min_{x\in I}\operatorname{cl}_{1}g(x)$ consists of just one endpoint of J, say a. If a sequence $\{x_{i}\}\subset J$ exists converging to a with $g(x_{i})$ converging to $\min_{x\in I}\operatorname{cl}_{1}g(x)=\operatorname{cl}_{1}g(a)$, then the result follows from the definition of $\operatorname{cl}_{1}g$. Otherwise there exists $x^{*} < a$ with $g(x^{*}) < g(a)$ and then the quasiconvexity of g implies that, for any $x\in J$,

$$g(x^*) < g(a)$$

$$\leq \max\{g(a), g(x)\},\$$

thus $g(x) \ge g(a)$, showing (3.6).

3.2. MULTIPLE-OBJECTIVE QUASICONVEX MAXIMIZATION PROBLEMS ON AN INTERVAL

In this subsection we show how to find a (weak) minimal dominator for the problem (P[F; I]) when I = [a, b] is a compact interval of \mathbb{R} .

By Lemma 7, for each i = 1, 2, ..., k, the set arg $\min_{x \in I} \operatorname{cl}_{I} F_{i}(x)$ is a nonempty closed subinterval of I, thus it has the form $[\alpha^{i}, \beta^{i}]$.

LEMMA 10. Let $x \in]a, \min_{1 \le i \le k} \beta^i[$ (respectively $x \in]\max_{1 \le i \le k} \alpha^i, b[$). Then, $a \in \mathcal{S}^{\geqslant}(x)$ (respectively $b \in \mathcal{S}^{\geqslant}(x)$).

Proof. Given $x \in]a, \min_{1 \le i \le k} \beta^i[$ and $j \in \{1, 2, ..., k\}$, it follows that $x < \beta^j$, thus, by the definition of β^j there exists $y^j \in \arg\min_{y \in I} \operatorname{cl}_I F_j(y)$ such that $x \in]a, y^j[$. Hence, by Lemma 8, $F_j(x) \le F_j(a)$ for all j, showing that $a \in \mathcal{S}^{\triangleright}(x)$. The other case is similar.

PROPOSITION 11. Define D_I^0 as

$$D_I^0 = \{a, b\} \cup \left[\min_{1 \le i \le k} \beta^i, \max_{1 \le i \le k} \alpha^i\right],$$

where it is understood that $[\min_{1 \le i \le k} \beta^i, \max_{1 \le i \le k} \alpha^i] = \emptyset$ if $\min_{1 \le i \le k} \beta^i > \max_{1 \le i \le k} \alpha^i$. Define also

$$D_{I} = \begin{cases} \{\alpha\}, & \text{if } F(a) \geq F(b) \\ \{b\}, & \text{if } F(b) \geq F(a), F(b) \neq F(a) \\ D_{I}^{0} \setminus \{x \neq a : a \in \mathcal{S}^{\geq}(x)\} \cup \{x \neq b : b \in \mathcal{S}^{\geq}(x)\} \}, & \text{otherwise} \end{cases}$$

One then has

- 1. $D_I^0 \in \mathcal{D}[F;I]$.
- 2. $D_I \in \mathcal{D}_{WM}[F;I]$.
- 3. If $a \in \mathcal{S}^{\geqslant}(b)$, $b \in \mathcal{S}^{\geqslant}(a)$, or each F_i is semistrictly quasiconvex on [a, b], then $D_I \in \mathcal{D}_M[F; I]$.

Proof. Part 1 follows from Lemma 10. To show 2, if $a \in \mathcal{S}^{\geq}(b)$ one would have for each $x \in I$ and $i \in \{1, 2, ..., k\}$ that

$$F_i(x) \le \max\{F_i(a), F_i(b)\}\$$

= $F_i(a)$,

thus $a\in \mathscr{S}^{\geqslant}(x)$; hence $D_I=\{a\}\in \mathscr{D}[F;I]$, which is (weak) minimal being a singleton. A similar result is obtained when $b\in \mathscr{S}^{\geqslant}(a)$, thus to finish the proof of 2, we assume that $a\not\in \mathscr{S}^{\geqslant}(b)$ and $b\not\in \mathscr{S}^{\geqslant}(a)$. In particular, $\{a,b\}\subset D_I$. Given $x\in [a,b]$, it follows from Part 1 that there exists some $y\in \mathscr{S}^{\geqslant}(x)\cap D_I^0$; if $y\not\in D_I$, then $y\not\in \{a,b\}$ and either $a\in \mathscr{S}^{\geqslant}(y)\subset \mathscr{S}^{\geqslant}(x)$ or $b\in \mathscr{S}^{\geqslant}(y)\subset \mathscr{S}^{\geqslant}(x)$, hence $\emptyset\neq \{a,b\}\cap \mathscr{S}^{\geqslant}(x)\subset D_I\cap \mathscr{S}^{\geqslant}(x)$; if $y\in D_I$ then $D_I\cap \mathscr{S}^{\geqslant}(x)\neq \emptyset$. Thus $D_I\in \mathscr{D}[F;S]$.

To show that $D_I \in \mathcal{D}_{WM}[F; I]$, suppose that, by contradiction, $x, y \in D_I$ exist such

that $x \in \mathcal{G}^{>}(y)$. Since either $x \in [a, y[\text{ or } x \in]y, b]$, we can assume w.l.o.g. that $x \in [a, y[$. Then, for each i

$$F_i(y) < F_i(x)$$

$$\leq \max\{F_i(a), F_i(y)\}$$

thus $F_i(y) < F_i(a)$ for each *i*. Hence $a \in \mathcal{S}^{>}(y)$, thus $y \not\in D_I$, which is a contradiction. Hence $D_I \in \mathcal{D}_{WM}[F; I]$, and this shows 2.

The minimality property of Part 3 was shown above for $a \in \mathcal{S}^{\geqslant}(b)$ or $b \in \mathcal{S}^{\geqslant}(a)$, so we show now the case of semistrictly quasiconvex functions F_i . Suppose that, on the contrary, $x, y \in D_I$ exist such that $x \in \mathcal{S}^{\geqslant}(y) \setminus \{y\}$. Since $y \in D_I$, one gets that $x \notin \{a, b\}$; then, $x \in]a, y[\cup]y, b[$, thus w.l.o.g. we assume $x \in]a, y[$. Since $x \in D_I \setminus \{a\}$, $a \notin \mathcal{S}(x)$, thus there exists some i with $F_i(a) < F_i(x)$, thus

$$F_i(a) < F_i(x)$$

$$\leq \max\{F_i(a), F_i(y)\},$$

thus $F_i(x) \le F_i(y)$, and, since $x \in \mathcal{S}^{\ge}(y)$, $F_i(x) = F_i(y)$, which contradicts the semistrict quasiconvexity of F_i . Hence, $D_i \in \mathcal{D}_M[F;I]$.

REMARK 12. In Part 1 of Proposition 11, a dominator has been constructed, consisting of at most three intervals, two of which are reduced to a point. Moreover, such a dominator is easily derived once all the single-objective one-dimensional problems $\min_{x \in I} \operatorname{cl}_{I} F_{i}(x)$, i = 1, 2, ..., k have been solved.

On the other hand, it follows from the quasiconvexity of the functions F_i that the set $\{x \in I : a \in \mathscr{S}^{\geqslant}(x)\}$ (respectively $\{x \in I : b \in \mathscr{S}^{\geqslant}(x)\}$) is an interval having a (respectively b) as one of its endpoints. This implies that the set D_I , shown in Part 2 of Proposition 11 to be weak minimal, also consists of at most three intervals, two of which are reduced to the endpoints of I.

In the case of continuous F_i , finding the set $\{x : a \in \mathcal{S}^{\geq}(x)\}$ is reduced to finding, for each $i = 1, \ldots, k$, the highest root of the nonlinear equation $F_i(x) = F_i(a)$, which, due to the quasiconvexity of F_i can be solved with any prespecified accuracy by e.g. binary search.

REMARK 13. For the biobjective case (k = 2), the interval $[\min_k \beta^k, \max_k \alpha^k]$ is, by construction, such that, within it, both F_1 and F_2 are monotonic: one non-decreasing and the other nonincreasing. Hence, for the biobjective case, there is no loss of generality in assuming that functions F_i are not only quasiconvex but also quasiconcave on the intervals $[\min_k \beta^k, \max_k \alpha^k]$.

EXAMPLE 1. Let I = [0, 4], and consider the three quasiconvex functions F_1, F_2, F_3 defined as

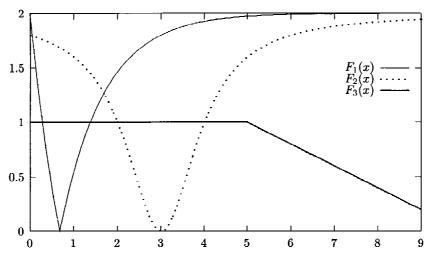


Figure 3. Functions of Example 1.

$$F_1(x) = |4e^{-x} - 2|$$

$$F_2(x) = \frac{2(x-3)^2}{1 + (x-3)^2}$$

$$F_3(x) = \min\left\{1, 2 - \frac{x}{5}\right\},$$

depicted in Figure 3.

In order to construct the dominator(s) described in Proposition 11, we must determine first the set $[\alpha_i, \beta_i]$ of minima on I for each F_i . These are respectively $\{\ln 2\} = \{0.6931\}, \{3\}$ and [0, 4]. This yields

i	$lpha_i$	$oldsymbol{eta}_i$
1	0.6931	0.6931
2	3	3
3	0	4

For this we obtain the dominator

$$D_I^0 = \{0, 4\} \cup [0.6931, 3]$$

Moreover, by comparing the endpoints, we get

$$F(0) = (2, 1.8, 1)$$

$$F(4) = (1.9267, 1, 1),$$

thus $F(0) \ge F(4)$. Hence, by Proposition 11, the set $D_I = \{0\}$ is not only a weak minimal dominator but also a minimal dominator.

Suppose now that the feasible region is the interval I = [5, 9]. In this case we obtain

i	$lpha_{_i}$	$oldsymbol{eta}_i$
1	5	5
2	5 5 9	5
3	9	9

From this it is easily seen that

$$D_I^0 = D_I = [5, 9]$$
.

Since all the functions are semistrictly quasiconvex in I, it follows that I is a minimal dominator for $\max_{x \in I} F(x)$.

Finally, for I = [4, 9] we similarly obtain $D_I^0 = D_I = [4, 9]$, but in this case D_I^0 is not a minimal dominator, since [5, 9] is a strictly included dominator (which may be seen to be minimal).

3.3. MULTI-OBJECTIVE QUASICONVEX MAXIMIZATION PROBLEMS ON A SET OF INTERVALS

As a natural extension of the model presented in Section 3.2, we address here the problem

$$\max_{x \in X} F(x) ,$$

where

- $X = \bigcup_{1 \le i \le t} I_i$, with $\{I_i\}_{1 \le i \le t}$ being a family of compact (possibly degenerate) intervals of the real line, not necessarily disjoint,
- F_1, \ldots, F_k are quasiconvex on each I_i , $i = 1, \ldots, t$. (Note that this is a weaker assumption than each component of F to be quasiconvex in the convex hull of $\bigcup_{1 \le i \le t} I_i$).

By Proposition 2, if one finds, for each $i=1,2,\ldots,t$ some dominator $D_i\in \mathscr{D}[F;I_i]$, then any $D\in \mathscr{D}[F;\bigcup_{1\leq i\leq t}D_i]$ would serve as a dominator for $(P[F;\bigcup_{1\leq i\leq t}I_i])$. Moreover, if a (weak) minimal dominator is sought, redundant alternatives should be purged, either in the construction of the sets D_i (by imposing e.g. $D_i\in \mathscr{D}_{\mathscr{WM}}[F;I_i]$) or when they are merged to produce a (small) final dominator.

To approximate this goal one can use a Branch-and-Bound scheme, similar to the one described in [14]: we start with a list \mathscr{L} of compact intervals, the union of which is known to be a dominator for $(P[F; \bigcup_{1 \le i \le l} I_i])$, and then refine iteratively

the elements in \mathcal{L} , by making pairwise comparisons, in such a way that, at any stage, one has

$$\bigcup_{I\in\mathscr{L}}I\in\mathscr{D}[F;\bigcup_{1\leqslant i\leqslant t}I_i]$$

To perform comparisons among elements in \mathcal{L} we introduce, for each interval I := [a, b] contained in some I_i , the vectors M(I), $UB(I) \in \mathbb{R}^k$ of evaluations at the midpoint of I and a componentwise upper bound of F, respectively:

$$M(I)_{j} = F_{j} \left(\frac{a+b}{2}\right)$$

$$UB(I)_{j} \ge \max_{x \in I} F_{j}(x)$$

REMARK 14. By the quasiconvexity of F_j on $I \subset I_i$, it follows that one may choose

$$UB(I)_{i} = \max\{F_{i}(a), F_{i}(b)\} \quad j = 1, 2, ..., k$$

Note also that for $I = \{a\}$ we have M(I) = UB(I).

From the definitions of the vectors M and UB one immediately obtains the following way to check whether some interval J can be discarded from further consideration in the Branch-and-Bound scheme.

PROPOSITION 15. Given nonempty compact intervals I, J, suppose F is continuous on I and on J. Then the following statements are equivalent:

- 1. $I \in \mathcal{D}[F; J]$, i.e. for any $y \in J$ there exists $x \in I$ with $F(x) \ge F(y)$
- 2. $0 \le \min_{y \in J} \max_{x \in I} \min_{1 \le j \le k} (F_j(x) F_j(y))$.

This is implied by both

$$\bigcap_{1 \le j \le k} \left\{ x \in I : F_j(x) \ge UB(J)_j \right\} \ne \emptyset$$
(3.7)

and

$$0 \le \min_{1 \le i \le k} \left(M(I)_j - UB(J)_j \right), \tag{3.8}$$

while (3.8) always implies (3.7).

Proof. The equivalence between 1 and 2 is evident. Since (3.7) is equivalent to the existence of some $x \in I$ with

$$F(x) \ge F(y) \ \forall y \in J \,, \tag{3.9}$$

it clearly implies 1. On the other hand, (3.8) is equivalent to (3.9) for x fixed to the midpoint of I. Hence, (3.8) implies (3.7) and the result follows.

Although condition (3.8) is easier to implement, the stronger test (3.7) is also of practical interest since this intersection set, if nonempty, has a simple structure due to the quasiconvexity of F, as indicated by the following simple result:

PROPOSITION 16. One has for any values c_i

- 1. Each set $\{x \in I : F_j(x) \ge c_j\}$ consists of at most two intervals, each with an endpoint of I as one of its endpoints.
- 2. For $k_0 = 1, 2, ..., k$, the set $\bigcap_{1 \le j \le k_0} \{x \in I : F_j(x) \ge c_j\}$ is a collection of $n(k_0)$ intervals, with

$$n(1) \le 2$$

 $n(k_0) \le n(k_0 - 1) + 1$, $k_0 = 2, 3, ..., k$

The basic steps of the Branch-and-Bound procedure are described below:

```
Algorithm 1

Initialization:
Set \mathcal{L} := \{cl(D_{I_j}), \ j=1,\ldots,t\}
Set r := 1

Iteration r = 1, 2, \ldots,:
for all I \in \mathcal{L} do

If, for some J \in \mathcal{L}, \ J \neq I, (3.8) or (3.7) hold, then delete I from \mathcal{L};
Else, if I is non-degenerate do split I into I_1 and I_2 at the midpoint of I; replace I by I_1 and I_2 in \mathcal{L};
GoTo Iteration r + 1
```

Before discussing the output of the algorithm in the limit $(r = \infty)$, let us present an illustrative example.

EXAMPLE 1 (Cont.)

Let F be the three-objective one-dimensional function described in the first part of the Example, and assume now that the feasible region X consists of the two compact segments $I_1 = [0, 4]$, and $I_2 = [5, 9]$.

In the Initialization phase, we must construct the sets $cl(D_{I_j})$, j = 1, 2. This was already done in the first part of the Example, thus we start with the list

```
\mathcal{L} = \{\{0\}, [5, 9]\}.
```

Then, we go to Iteration 1. For each interval I (degenerate or not) in \mathcal{L} , the vectors M(I), UB(I) must be constructed. (Observe that this task becomes trivial using Remark 14 above.) Evaluations at the endpoints, 0, 5, 9 and the midpoint 7 yield

$$F(0) = (2, 1.8000, 1)$$

$$F(5) = (1.9730, 1.6000, 1)$$

$$F(7) = (1.9964, 1.8824, 0.6000)$$

$$F(9) = (1.9995, 1.9459, 0.2000)$$

We then obtain

I	M(I)	UB(I)
{0}	(2, 1.8000, 1)	(2, 1.8000, 1)
[5, 9]	(1.9964, 1.8824, 0.6000)	(1.9995, 1.9459, 1)

We will only use the simplest test, namely, (3.8) in the algorithm. Since no pair of intervals in \mathcal{L} satisfies condition (3.8), we go to Iteration 2 with the list of intervals

$$\mathcal{L} = \{\{0\}, [5, 7], [7, 9]\}.$$

Two new midpoints appear, namely, 6 and 8, with objective values

$$F(6) = (1.9901, 1.8000, 0.8000)$$
$$F(8) = (1.9987, 1.9231, 0.4000).$$

This enables us to update the table of vectors M, UB yielding

I	M(I)	UB(I)
{0}	(2, 1.8000, 1)	(2, 1.8000, 1)
[5, 7]	(1.9901, 1.8000, 0.8000)	(1.9964, 1.8824, 1)
[7, 9]	(1.9987, 1.9231, 0.4000)	(1.9995, 1.9459, 0.6000)

As in the previous iteration, no pair of intervals satisfies condition (3.8), and we go to Iteration 3 with the updated list of intervals

$$\mathcal{L} = \{\{0\}, [5, 6], [6, 7], [7, 8], [8, 9]\}$$

The new midpoints give objective values

$$F(5.5) = (1.9837, 1.7241, 0.9000)$$

$$F(6.5) = (1.9940, 1.8491, 0.7000)$$

$$F(7.5) = (1.9978, 1.9059, 0.5000)$$

$$F(8.5) = (1.9992, 1.9360, 0.3000)$$

I	M(I)	UB(I)
{0}	(2.18000.1)	(2. 1.8000, 1)

With this, our new table of vectors M, UB is given by

I	M(I)	UB(I)
{0}	(2, 1.8000, 1)	(2, 1.8000, 1)
[5, 6]	(1.9837, 1.7241, 0.9000)	(1.9901, 1.8000, 1)
[6, 7]	(1.9940, 1.8491, 0.7000)	(1.9964, 1.8824, 0.8000)
[7, 8]	(1.9978, 1.9059, 0.5000)	(1.9987, 1.9231, 0.6000)
[8, 9]	(1.9992, 1.9360, 0.3000)	(1.9995, 1.9459, 0.4000)

In this case, the sufficient condition for dominance is satisfied for the pair of intervals {0} and [5, 6], so the interval [5, 6] can be excluded for further considerations.

We would then obtain a reduced list

$$\mathcal{L} = \{\{0\}, [6, 7], [7, 8], [8, 9]\}$$

to start Iteration 4, if desired.

The following theorem shows that the successive steps of the algorithm above provide a sequence of nested compact dominators, converging to a dominator which, under mild further assumptions on the functions F_i , enjoys minimality properties:

PROPOSITION 17. Denote by D_r the union of all intervals of \mathcal{L} at the end of iteration r, and by D^* the compact set

$$D^* = \bigcap_{r=1}^{\infty} D_r$$

- 1. $D_1 = X = \bigcup_{1 \le i \le t} I_i$ and $D_{r+1} \subset D_r$ for all r.
- 2. If F is upper-semicontinuous, then

$$D^* \in \mathcal{D}[F;X] \tag{3.10}$$

3. Moreover, if F is continuous, then

$$D^* \in \mathcal{D}_{\mathcal{W}, \mu}[F; X] \,. \tag{3.11}$$

Proof. The first property is evident from the algorithm.

By construction, each D_r is compact, thus their intersection is also compact. Moreover, $D_r \in \mathcal{D}[F; X]$, thus, by Proposition 3, (3.10) follows.

To show (3.11), suppose, on the contrary, that there exist $x_1, x_2 \in D^*$ with $x_1 \in \mathcal{S}^{>}(x_2)$. If, for each i = 1, 2 and $r = 1, 2, \ldots$, we denote by \mathcal{I}_i^r the class of intervals I_i^r in the list at stage r with $x_i \in I_i^r$, it will follow from the splitting process that there exists some r_0 such that, for each $r \ge r_0$, and each $I_i^r \in \mathcal{I}_i^r$

$$x_1 \not\in I_2^r$$
, and $x_2 \not\in I_1^r$

Since the functions F_i are continuous, thus uniformly continuous on X, there would exist some r such that for each $I_i^r \in \mathcal{I}_i^r$

$$F_i(x) > F_i(y)$$
 for all $x \in I_1^r$ and $y \in I_2^r$, $j = 1, 2, ..., k$

Hence $UB(I_2^r) < M(I_1^r)$, implying that I_2^r (thus x_2) would have been deleted prior to stage r by (3.8), thus $x_2 \not\in D^*$, which is a contradiction.

4. Multiple-objective multi-dimensional problems

For the single-objective case (i.e., if k = 1 in $(P[F_1; S])$), it is a well-known result of Global Optimization that, if S is a polytope and F_1 is quasiconvex on S, then the set of vertices of S is a dominator for $(P[F_1; S])$, [15].

In other words, if, for j = 0, 1, ..., n, \mathcal{F}^j denotes the set of points of a polytope S contained in some j-dimensional face of S, then

$$\mathcal{F}^0 \in \mathcal{D}[F_1; S] \tag{4.12}$$

The next proposition extends assertion (4.12) to multiple-objective quasiconvex problems. To show it, we will use the following

LEMMA 18. Let P be a polyhedron in \mathbb{R}^n , and let H_1, H_2, \ldots, H_t be closed halfspaces in \mathbb{R}^n . If x^* is an extreme point of $P \cap \bigcap_{1 \le i \le t} H_i$, then x^* belongs to some face of P with dimension not greater than t.

Proof. Let P be represented as

$$P = \{x \in \mathbb{R}^n : a'_r x \leq b_r \text{ for all } r \in R\}$$

for some finite index set R, and let each H_i be given as

$$\{x \in \mathbb{R}^n : c'_i x \leq d_i\}$$

Define the sets of active indices $R(x^*)$ and $T(x^*)$ as

$$R(x^*) = \{r \in R : a'_r x^* = b_r\}$$

$$T(x^*) = \{i, 1 \le i \le t : c'_i x^* = d_i\}$$

Then x^* belongs to the face F of P,

$$F = P \cap \{x \in \mathbb{R}^n : a_r'x = b_r \ \forall r \in R(x^*)\}$$

We will show that F has dimension not greater than t. Indeed, since x^* is, by assumption, an extreme point of $P \cap \bigcap_{1 \le i \le t} H_i$, then the set of vectors $\{a_r\}_{r \in R(x^*)} \cup \{c_i\}_{i \in T(x^*)}$ has rank

$$rank({a_r}_{r \in R(x^*)} \cup {c_i}_{i \in T(x^*)}) = n$$

Hence, denoting by $|T(x^*)|$ the cardinality of $T(x^*)$, one obtains

$$\operatorname{rank}(\{a_r\}_{r \in R(x^*)}) \ge n - |T(x^*)|$$

$$\ge n - t.$$

thus the dimension of F cannot be greater than t.

PROPOSITION 19. Let S be a polytope in \mathbb{R}^n , let $k \leq n+1$, and let F_1, \ldots, F_k be quasiconvex functions on S, all but possibly one of which are upper-semicontinuous. Then

$$\mathcal{F}^{k-1} \in \mathcal{D}[F;S] \tag{4.13}$$

Proof. Without loss of generality we assume that F_1, F_2, \dots, F_{k-1} are uppersemicontinuous on S. We will show that, for any $x \in S$,

$$\mathcal{S}^{\geqslant}(x) \cap \mathcal{F}^{k-1} \neq \emptyset \tag{4.14}$$

Let $x \in S$, and denote by $\mathcal{A}(x)$ the index set

$$\mathcal{A}(x) = \{i, 1 \le i \le k-1, F_i(y) < F_i(x) \text{ for some } y \in S\}.$$

If $\mathcal{A}(x)$ is empty, we would have

$$F(y) \ge F(x) \forall y \in S$$
,

thus any vertex y^* of S satisfies $y^* \in \mathcal{S}^{\geq}(x)$. Hence

$$\emptyset \neq \mathcal{S}^{\geqslant}(x) \cap \mathcal{F}^0 \subset \mathcal{S}^{\geqslant}(x) \cap \mathcal{F}^{k-1}$$

showing (4.14).

We consider now the case $\mathcal{A}(x) \neq \emptyset$. For each $i \in \mathcal{A}(x)$, the convex set $\{y \in S : F_i(y) < F_i(x)\}$ is open in S (its complement is closed due to the uppersemicontinuity of F_i) and does not contain x. Hence, there exists some nonzero vector u^i such that

$$\langle u^i, y - x \rangle > 0$$
 for all $y \in S$ with $F_i(y) < F_i(x)$, (4.15)

where $\langle \cdot, \cdot \rangle$ stands for the usual scalar product in \mathbb{R}^n .

Consider the polyhedron S(x),

$$S(x) = S \cap \{ y \in \mathbb{R}^n : \langle u^i, y - x \rangle \le 0, \quad \forall i \in \mathcal{A}(x) \},$$

which is nonempty because $x \in S(x)$. Consider the optimization problem

$$\max_{y \in S(x)} F_k(y) \tag{4.16}$$

Since F_k is quasiconvex on the nonempty polyhedron S(x), (4.16) has an optimal solution at some vertex y^* of S(x). We will show that

$$y^* \in \mathcal{F}^{k-1} \cap \mathcal{S}^{\geqslant}(x) \tag{4.17}$$

Since y^* is a vertex of S(x), Lemma 18 implies that

$$y^* \in \mathcal{F}^{|\mathcal{A}(x)|} \subset \mathcal{F}^{k-1} \tag{4.18}$$

Since $x \in S(x)$ and y^* is optimal for (4.16),

$$F_k(y^*) \ge F_k(x) \tag{4.19}$$

By definition of $\mathcal{A}(x)$,

$$F_i(y^*) \ge F_i(x) \, \forall i \in \{1, 2, \dots, k-1\} | \mathcal{A}(x)$$
 (4.20)

and by (4.15) and the fact that $y^* \in S(x)$,

$$F_i(y^*) \ge F_i(x) \,\forall i \in \mathcal{A}(x) \tag{4.21}$$

Joining (4.18–4.21), (4.17) holds, thus
$$\mathscr{S}^{\geqslant}(x) \cap \mathscr{F}^{k-1} \neq \emptyset$$
, as asserted.

REMARK 20. The assumption of upper-semicontinuity of at least k-1 functions is not superfluous, as the following example shows: let k=2, n=2, $S=[0,1]\times[0,1]$, and the functions F_1, F_2 defined as

$$F_1(x) = \begin{cases} 0, & \text{if } x_2 > \frac{1}{2} \text{ or } x = (0, \frac{1}{2}) \\ 1, & \text{otherwise} \end{cases} \quad F_2(x) = \begin{cases} 0, & \text{if } x_2 > \frac{1}{2} \text{ or } x = (1, \frac{1}{2}) \\ 1, & \text{otherwise} \end{cases}$$

Both functions are quasiconvex but are not upper-semicontinuous; let $x^* = (\frac{1}{2}, \frac{1}{2})$. It is easily seen that

$$\mathcal{S}^{\geq}(x^*) = \{(\lambda, \frac{1}{2}) : 0 < \lambda < 1\}$$

thus $\mathscr{S}^{\geqslant}(x^*) \cap \mathscr{F}^1 = \emptyset$, showing that \mathscr{F}^1 is not a dominator.

As a consequence of Propositions 19 and 2, one obtains

COROLLARY 21. Let S be the union of t polytopes S_1, \ldots, S_t in \mathbb{R}^n . Let F_1, \ldots, F_k be $k \leq n+1$ real-valued functions on S. On each S_j , let all F_i be quasiconvex and all but possibly one F_i be lower-semicontinuous. Then the union of all k-1-faces of all S_j is a dominator for P[F; S].

Proposition 19 also enables us to derive localization results for single-objective problems.

COROLLARY 22. Let S be the union of t polytopes S_1, \ldots, S_t in \mathbb{R}^n . Let F_1, \ldots, F_k be $k \leq n+1$ real-valued functions on S, quasiconvex on each S_j . For any componentwise nondecreasing $\Phi: F(S) \to \mathbb{R}$ such that Problem

$$\max_{x \in S} \Phi(F_1(x), F_2(x), \dots, F_k(x))$$

has an optimal solution, the union of the set of k-1-faces of all S_j also contains an optimal solution.

In particular, for F_1, F_2, \ldots, F_k linear fractional functions with positive denominators on a polytope S, which are well-known to be quasiconvex (see e.g. [1], p. 165) and $\Phi(s_1, s_2, \ldots, s_k) = s_1 + s_2 + \cdots + s_k$, or $\Phi(s_1, s_2, \ldots, s_k) = \max\{s_1, s_2, \ldots, s_k\}$, we obtain

COROLLARY 23. The minimum of the sum (respect. the maximum) of k linear fractional functions with positive denominators over a polytope S in dimension $n \ge k-1$ is attained at some k-1-face of S.

This generalizes the results known for the case k = 2 (see the review of [27] and the references therein). For an application see [25].

REMARK 24. For biobjective problems (k = 2), since both F_j are quasiconvex on each edge, after embedding such edges as compact intervals of the real line, one can use the results in Section 3 to design an algorithm converging to a weak minimal dominator.

For the case of general k, Proposition 19 seems at the moment to be mainly of theoretical interest: In principle, Algorithm 1 can be generalized to the k-dimensional case, by replacing intervals by e.g. simplices, although the corresponding bounding scheme does not extend to the general case, and less efficient schemes, such as those proposed in [4, 14], should be used.

Nevertheless, this kind of localization results can be used to design new heuristic resolution methods of problems of the form $\min_{x \in S} \Phi(F(x))$, where k, the number of components of F, is very small, and, in particular, much smaller than the dimension n of the space.

We know then that the search for optimal solutions can be reduced to the k-1-dimensional faces of S, so that algorithms which alternate a global search in a given low-dimensional face with moves to adjacent low-dimensional faces, can be used.

5. Application: Location of a semi-obnoxious facility

Let $S = S_1 \cup S_2 \cup \cdots \cup S_t$, each S_i being a convex polygon in \mathbb{R}^2 . Two finite subsets \mathscr{A}^+ , \mathscr{A}^- of \mathbb{R}^2 are given. Associated with each $a \in \mathscr{A}^+$ we have a concave function $g_a : [0, +\infty) \to \mathbb{R}$ and a polyhedral gauge γ_a , [9, 10, 19], i.e., a Minkowski functional whose unit ball is a polytope.

Let $h:[0,+\infty)\to\mathbb{R}$ be a nonincreasing function, and consider the biobjective problem

$$\min_{x \in S} (F_1(x), F_2(x)), \tag{5.22}$$

where

$$F_1(x) = \sum_{a \in \mathcal{A}^+} g_a(\gamma_a(x - a))$$
$$F_2(x) = \max_{a \in \mathcal{A}^-} h(||x - a||),$$

 $\|\cdot\|$ being the euclidean norm.

This problem has its motivation in Continuous Location of semidesirable facilities, see [17, 23] for an introduction to Continuous Location in general and [6, 24] for semidesirable facility location models: A facility is to be located within region S, and will interact with individuals who want the facility close (those in \mathcal{A}^+) and others who want the facility far (those in \mathcal{A}^-). Interactions with \mathcal{A}^+ provide the first objective in (5.22): the minimization of the total transportation cost $F_1(x)$, where transportation cost from $a \in \mathcal{A}^+$ to x is given by a concave function g_a of the distance from a to x, the latter measured by the polyhedral gauge γ_a , [29].

On the other hand, interactions of the facility with \mathcal{A}^- provide the second objective F_2 , which measures the highest damage suffered by points in \mathcal{A}^- , where the damage suffered by $a \in \mathcal{A}^-$ is assumed to be given by a nonincreasing function h of the Euclidean distance from a to x, see [11, 24].

In practice, the two objectives of (5.22) are aggregated into a single criterion, yielding a problem of the form

$$\max \Phi\left(-\sum_{a \in \mathcal{A}^+} g_a(\gamma_a(x-a)), \max_{a \in \mathcal{A}^-} h(\|x-a\|)\right), \tag{5.23}$$

[6, 24] for some globalizing Φ , and the resulting problem (multimodal, as a rule), can be tackled e.g. by the 2-dimensional Branch and Bound method described in [13]. However, as shown below (Proposition 25), the search of an optimal solution for (5.23) can be restricted to a series of segments, thus (5.23) can be solved by simply using single-variable Global-Optimization techniques, [2, 12], which are usually much faster than their two-variable counterparts.

In order to obtain a dominator for (5.22) one should observe first that, since h is assumed to be nonincreasing, it suffices to obtain a dominator for problem

$$\max_{x \in S} (-F_1(x), \min_{a \in \mathcal{A}^-} ||x - a||) \tag{5.24}$$

(in fact, if h is decreasing, both problems are equivalent). Let us rewrite now (5.24) within our framework. For polyhedral gauges, using the concept of *elementary* convex set of [10], one can obtain a subdivision $\mathscr C$ of the plane into polyhedra in such a way that, within each $C \in \mathscr C$, each gauge γ_a is affine, see [9, 10] for further details. For instance, if each γ_a is the l_1 norm, then the polyhedral subdivision of the plane is obtained after constructing horizontal and vertical lines through each $a \in \mathscr{A}^+$, yielding a total of $O(|\mathscr{A}^+|^2)$ cells.

Moreover, defining, for each $a \in \mathcal{A}^-$, the *Voronoi cell V(a)* associated with a as

$$V(a) = \{x \in \mathbb{R}^2 : ||x - a|| \le ||x - b|| \text{ for all } b \in \mathcal{A}^-\},$$

the class $\mathcal{V} = \{V(a) : a \in \mathcal{A}^-\}$ also constitutes a polyhedral subdivision of \mathbb{R}^2 in $O(|\mathcal{A}^-|)$ polyhedra, which can be efficiently constructed in $O(|\mathcal{A}^-|\log |\mathcal{A}^-|)$, see e.g. [20, 26].

Consider now the class \mathcal{Z} of all Z of the form

$$S_i \cap C \cap V(a)$$

for some $i, 1 \le i \le t, C \in \mathcal{C}$ and $a \in \mathcal{A}^-$ which are nonempty. On each $Z \in \mathcal{Z}$, we have that $-F_1$ is convex (it is the composition of the convex function $-\Sigma_{a \in \mathcal{A}^+} g_a$ with the affine functions (within Z!) γ_a , and F_2 is also convex (recall that, for $Z \in \mathcal{Z}$ fixed, there exists some $a^* \in \mathcal{A}^-$ such that $\min_{a \in \mathcal{A}^-} \|x - a\| = \|x - a^*\|$). Hence, rewriting (5.24) as

$$\max_{x \in \bigcup_{T \in \mathscr{Q}^Z}} (-F_1(x), \min_{a \in \mathscr{A}^-} ||x - a||),$$

we can use Corollary 21 to obtain

PROPOSITION 25. The edges of the sets in \mathcal{Z} constitute a dominator for Problem (5.22).

After embedding the edges of polytopes in \mathcal{Z} as compact intervals of the real line, one can use the algorithm described in Section 3.3 to reduce the size of such dominator, converging (in case of decreasing h) to a weak minimal dominator.

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