Algebraic combinatorics in bounded induction

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ABSTRACT

Keywords:

Bounded induction Stone-Čech compactification Ramsey theorem Peano Arithmetic In this paper, new methods for analyzing models of weak subsystems of Peano Arithmetic are proposed. The focus will be on the study of algebro-combinatoric properties of certain definable cuts. Their relationship with segments that satisfy more induction, with those limited by the standard powers/roots of an element, and also with definable sets in Bounded Induction is studied. As a consequence, some considerations on the Π_1 -interpretability of $I\Delta_0$ in weak theories, as well as some alternative axiomatizations, are reviewed. Some of the results of the paper are obtained by immersing Bounded Induction models in its Stone-Cech Compactification, once it is endowed with a topology.

Contents

1.	Introduction	2
	1.1. Motivation	4
	1.2. Aim of the paper	7
	1.3. Structure of the paper	7
2.	Preliminaries and some starting results	7
	2.1. Subsystems of $I\Delta_0$	7
	2.2. Open induction and subsystems	8
	2.3. Exponentiation and segments determined by (standard) powers or integer roots	9
	2.4. Algebra on Stone-Čech compactification of a (discrete) semigroup	10
3.	Part I: maximal segments, kernels and gaps	11
	3.1. Definability of kernels	13
	3.2. Collapse of kernels and interpretability	16
	3.3. Arithmetic of gaps and segments	18
4.	Part II: algebraic combinatorics in IE_k	21
	4.1. Combinatorics in gaps	21

	4.2.	FP and FS-sets within gaps	2
	4.3.	Combinatorics in presence of weak overspill on \mathbb{N}	ŀ
5.	Closing	g remarks	3
Ackno	wledge	ments	3
Refere	ences .		3

1. Introduction

A significant amount of the research corpus on (models of) classic subsystems of Peano Arithmetic (PA) is primarily focused on achieving two types of goals. A first type refers to studying how powerful a theory (e.g. fragment of PA) is for supporting the results of elementary number theory. A second one is devoted to obtaining independence results for (open) central problems to Complexity Theory [16]. The first objective is usually focused on theories weaker than $I\Sigma_1$, because that theory is already enough to develop a large part of Number Theory. A *frontier* for such an analysis is Parikh's theory $I\Delta_0$ (the theory of bounded induction), where several problems are related to its power for developing number theory remains open, some of these being related to open problems in Complexity Theory. Its "frontier" nature is because a large number of these questions can be solved if $I\Delta_0$ is extended with an axiom that ensures that a certain total function grows faster than any polynomial function. For example, the exponential function. Please let us recall that according to a celebrated result of R. Parikh this kind of functions are not definable in $I\Delta_0$ (cf. [22], Th. V.1.4):

Result 1.1. Every provably recursive function in $I\Delta_0$ is bounded by a polynomial.

Concerning the second research line (the searching for independence results), the situation is similar. For example, in strong fragments, a classical approach to solve this type of problems is based on showing the existence, in any model of the subsystem, of an initial segment (or final extension), that is a model of a stronger theory. This technique is useful to separate theories [40], as well as for connecting model-theoretic and complexity issues (see e.g. the survey [38]). In the same way as in the first research topic, when one works with fragments of PA, almost all positive results are also limited to proper extensions of $I\Delta_0$ (such as $I\Sigma_{n+1}$ or $B\Sigma_{n+1}$).

The scenario substantially changes when the theories $\{IE_n\}_{n\in\omega}$ (fragments of $I\Delta_0$ defined like $I\Sigma_n$ but using bounded quantifiers instead) are considered. It is not known if this hierarchy is strict. In fact there are relatively very few results about the relationship between IE_k and IE_{k+1} ($k \ge 1$). In stronger subsystems, global coding tools are used to prove results, but this type of functions are not definable in Bounded Induction by Result 1.1, so they can not be used (in particular, arithmetization).

We could affirm that one of the reasons for this situation is that the relationship of these theories with classes of computational complexity is not so evident to be able to translate problems between both fields. Let us consider, for example, the case of the class of provably recursive theories on IE_n . In general, a useful approach to studying fragments of PA is to consider, given a T theory, the class R(T) of provably recursive functions in T and to study, for example, the existence of initial segments that are closed under that class of functions in (some) Arithmetic models (see [9] for a refinement of up to IE_1). However, the approach could be not useful here in our framework, because $\{R(IE_n)\}_{n < \omega}$ collapses.

Proposition 1.2. $R(I\Delta_0) = R(PA^-).$

Proof. (taken from [4]) Assume that f is provably recursive in IE_n by the formula $\exists z\theta(x, y, z)$, with $\theta \in E_k$. Then there is a Π_1 -sentence $\forall u\psi(u)$ (suppose $\psi \in E_{n+1}$), such that $PA^- + \forall u\psi(u) \vdash \forall x \exists ! y \exists z\theta(x, y, z)$. Then

$$\mathrm{PA}^- \vdash \forall u \psi(u) \to \forall x \exists v \exists y, z < v \varphi(x, y, z, v)$$

where $\varphi \in \mathbf{E}_{k+1}$ is

$$\varphi(x, y, z, v) := \theta(x, y, z) \land \forall y_1, y_2, z_1, z_2 < v(\theta(x, y_1, z_1) \land \theta(x, y_2, z_2) \to y_1 = y_2)$$

Therefore,

$$\mathrm{PA}^- \vdash \forall x \exists v \exists u (\neg \psi(u) \lor \exists y, z < v \varphi(x, y, z, v))$$

By 1.1, there is $p(x) \in \mathbb{N}[x]$ that bounds v, u. Hence the E $\max_{\{n+3,k+1\}}$ -formula

$$\gamma(x, y, z) \equiv \exists v < p(x) \begin{cases} \forall u \le p(x)\psi(u) \land \varphi(x, y, z, v) \land y < v \land z < v \\ \lor \\ \exists u \le p(x) \neg \psi(u) \land y = 0 \end{cases}$$

represents the standard graph of f and $PA^- \vdash \forall x \exists ! y \exists z \gamma(x, y, z)$. \Box

A plausible way to attack the problem of the possible collapse of $\{IE_n\}_{n\in\omega}$ (or, for instance, whether $IE_1 \vdash IE_2$) is to find intermediate alternative axiom schemes to facilitate the analysis. Due to the essentially different behavior of the recursive function classes in Bounded Induction, one option might be to design schemes that describe algebraic properties or properties about, for example, the order < over \mathbb{N} . This idea has been used frequently in the strong fragments. In [36] Kreuzer and Yokoyama show interesting intermediate schemes between $I\Sigma_1$ and $I\Sigma_2$ of different nature, which are equivalent if $I\Sigma_1$ is considered as the base theory, and which are equivalent in turn to the well-foundedness of ω^{ω} .

Among the published works on IE_n , some of them deserve to be highlighted. There are refinements of proven facts for $I\Sigma_n$ (for instance the existence in nonstandard models of IE_1 of a nonstandard initial segment model of PA [29,42]), studies of algebraic nature (see e.g. [5,43] for IE_0 and [44] for IE_1) and other studies about alternative axiomatizations [3,15]. Whether $I\Delta_0$ is finitely axiomatizable -or whether $\{IE_n\}_{n<\omega}$ collapses- is the most important open problem in the field, connected with the P = NP? problem (cf. [7,22]).

A noticeable result about the theories IE_n shall be used in this paper (due to Paris and Dimitracopoulos [14], that is proved by using a formalization of truth definition for E_n -formulas):

Result 1.3. For all *n* there is a Π_1 -sentence σ_n such that

$$I\Delta_0 \vdash \sigma_n \vdash IE_n$$

Please, let us note that the Result 1.3 implies that, if $\{IE_n\}_n$ collapses then $I\Delta_0$ is finitely axiomatizable. Result 1.3 has been used by R. Kaye for studying models of IE₁ [29].

The above result is useful for our purposes because the reciprocal one -i.e., looking for a property that, added to IE_n a theory equivalent to $I\Delta_0$ can be obtained- is more difficult to study in general, because it is possible that adding to IE_1 a numerical property (or for example Ramsey type scheme) stronger theories than the $I\Delta_0$ itself are obtained. Two examples would be [10] (numerical property) and [8] (Ramsey type scheme). In the first paper, P. D'Aquino described the difference between IE_1 and $I\Delta_0 + exp$ employing a property of number theory:

$$IE_1 + P \vdash I\Delta_0 + exp$$

where P is an axiom that states that every Pell equation has a solution. In the second one, C. Cornaros studied the strength of weak forms of the Regularity Principle in the presence of IE₁, proving that a Bounded Weak Regularity Principle on E_1 formulas is equivalent -over IE₁- to I Δ_0 + exp.

In general, to investigate the relationship between the theorems in PA (or $I\Delta_0$) and those of PA⁻ it is common to focus on the inductive formula(s) necessary to demonstrate the corresponding theorem.

Definition 1.4. A formula $\varphi(x)$ is called *inductive* if $PA^- \vdash \varphi(0)$ and $PA^- \vdash \forall x(\varphi(x) \rightarrow \varphi(x+1))$.

The analysis for the inductive formulas is interesting by itself, as well as the study of their behavior in a model [23]. An interesting result for this paper is the following (in Section 3.2 variants of the result using a fixed inductive formula will be studied).

Theorem 1.5. (Wilkie-Paris, see [50]) Let $\varphi(x)$ be a bounded formula. The following conditions are equivalent:

- 1. $I\Delta_0 + \exp \vdash \forall x \varphi(x)$
- 2. There exists an inductive formula ψ such that $PA^- \vdash \forall x(\psi(x) \rightarrow \varphi(x))$

1.1. Motivation

The underlying general hypothesis that guides this work is that to advance in the study of (very weak) subsystems of bounded induction it could be interesting to explore the design and use of new tools that exploit properties of natural numbers that are essentially different in nature from those of classical fragments. For example, techniques from other research fields that are not traditionally related to the study of arithmetic models (such as the topology of semigroups), thus changing the point of view with which the models are analyzed.

The working hypothesis is relevant when the interest is on theories that extend very weak theories: IE_0^1 , the extension of PA⁻ allowing Euclidean division; and RIE_0^1 , that extends it assuring the existence of integer roots. Therefore, we will work with the absence of definable functions of rapid growth (that is, asymptotically faster than any polynomial) which obstructs us from using global arithmetization techniques (that are useful to reproduce classical demonstrations).

Concerning the selection of algebraic properties that may be interesting and how to use them, let us consider the paradigmatic case of the set of prime numbers. It is unknown whether $I\Delta_0$ demonstrates its infinitude. Thus, whilst we cannot assume that this set is unbounded in a $I\Delta_0$ model, some combinatorial properties associated with its infinity in the standard model can be studied. For example, a very interesting combinatorial property of such a set is that it contains arbitrarily long arithmetic progressions (a celebrated result of Green and Tao [47], see also [19]). We intend to approach the study of properties of this type using some results of the Ramsey type.

In general terms, Ramsey type theorems state that if the universe to be studied is divided into a series of parts, then at least one of them enjoys a particular property (commonly associated with a certain kind of regularity). In the finite case, Ramsey's theory refers to how large the set of parts must be to ensure that the property is satisfied, while in the infinite case the properties focus on detecting a certain regularity in one of those parts. Unlike the seminal result -Ramsey's theorem- these properties can now be studied with non-standard methods (see e.g. [26]). A well-known example on that line is Schur's theorem: given n > 0, for any prime number p large enough there is a non-trivial solution of $x^n + y^n = z^n \pmod{p}$. A nice proof is based on the fact that there are three monochromatic elements in a certain \mathbb{N} partition (a Ramsey property). Although there are studies of the peculiarities about Ramsey Theorem on theories weaker than $I\Sigma_1$ (see for example [52]), in general, these are focused on theories stronger than $I\Delta_0$.

We are concerned with two general questions. On the one hand, are these types of combinatorial properties on infinite sets (actually, a suitable version) provable in Bounded Induction, independently of the ignorance about the provability of their infinitude? On the other hand, if it is possible to demonstrate a

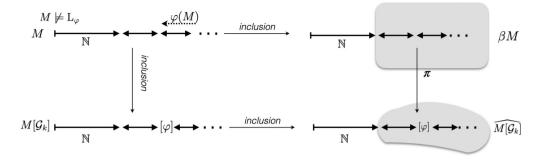


Fig. 1. Compactification of the model augmented with Δ_0 -definable sets without first element.

result of (algebro-)combinatorial regularity of that type, what is the relationship between its corresponding axiomatization as an axiom scheme and $I\Delta_0$ itself? It should be noted that the results of a combinatorial nature had already been obtained. To quote a recent one, the work of C. Cornaros [8] on the principle of regularity for bounded formulas. The principle of regularity states that if the elements of an infinite set are colored with a finite palette then there is a monochromatic infinite subset. Regularity might be seen as a result of Ramsey type.

To simplify the handling of these types of sets and their properties we will approach the question in a dual (although equivalent) way. Instead of working with combinatorial properties of bounded inductive sets in a model of weak arithmetic, we will work with nonempty sets without the first element instead (i.e., the corresponding minimization principle fails).

Combinatorial properties of that nature (those showing remarkable algebraic regularities in certain *infi*nite sets) can be attacked by working on an extension of the model that enables us to work within a richer structure. A widely used choice in the case of \mathbb{N} is its Stone-Čech compactification, $\beta\mathbb{N}$, which has proved to be an extremely useful structure for rewriting properties of Combinatorial Number Theory in another algebraic (topological) ones about $\beta\mathbb{N}$. Typically, it can be rewritten in terms of the existence of a specific ultrafilter (or a point in another type of compactification). Next, the corresponding property on $\beta\mathbb{N}$ can be addressed using very powerful techniques (see the surveys [2,13,24], or the recent example [45] where the author analyzes of the mentioned set of prime numbers from that perspective). Our interest in $\beta\mathbb{N}$ (or βM being M a model of weak arithmetic) lies in the fact that there are strong relationships between the infinite sets in \mathbb{N} and elements (ultrafilters) in $\beta\mathbb{N}$ (as expected for the construction of this one). Infinite sets of \mathbb{N} induce open neighbors in $\beta\mathbb{N}$ that contains non-principal ultrafilters with interesting properties. The selection of a suitable ultrafilter will allow to establish the Ramsey type discussed findings.

Continuing with the idea, a model $M \models PA^-$ can be viewed as an ordered semigroup, thus it is susceptible to being studied algebraically. If the model is endowed with a topology (for example the discrete one), results and techniques developed for topological semigroups can be used. The potential application would not be interesting if, additionally, if one does not work on an adequate compactification of the structure, in order to obtain a richer topology than the discrete one, with interesting properties as βM . We think that we have in that way a certain useful analogy between this kind of compactness and saturation in models of Arithmetic. Both properties allow assuring the existence of a certain element if an adequate approximation to it (a sequence in the compact case or a finitely satisfiable type in the case of saturation) can be obtained. We are particularly interested in that relationship for the special case of Stone-Čech compactification βM .

More precisely, the classic construction of $\beta \mathbb{N}$ has inspired us the following framework (shown in Fig. 1). In more detail, a model of Arithmetic M is extended to a structure $M[\mathcal{G}_k]$ where the Δ_0 -definable sets without minimum are also elements; they will be interpreted in a similar way to Henkin's cuts. The notion of cut induces a definable equivalence relationship ~ between that sets (two sets are equivalent if they induce the same cut), switching from that mode to working with equivalency classes [φ]. The set $M[\mathcal{G}_k]$ is endowed with the Hausdorff topology generated by the basis:

$$\{\{a\}\}\$$
 if $a \in M$ and $\{(a,b) \subseteq M[\mathcal{G}_k] : M \models a < \varphi < b\}$ if $[\varphi] \notin M$

(Please, let us note that $M[\mathcal{G}_k]$ under the equivalence relation ~ is the so-called Kolmogorov quotient of this topology.)

With that topology the set M is dense in $M[\mathcal{G}_k]$. Since βM is the biggest compactification of M, there is a projection

$$\pi:\beta M\to \widehat{M[\mathcal{G}_k]}$$

between βM and any right-topological semigroup compactification of $M[\mathcal{G}_k]$ we select (cf. th. 21.4 of [24]). Since $\pi \upharpoonright_M = Id_M$ is an homomorphism between dense sets (βM is endowed with the semigroup structure that it will be described in Sect. 2.4), then π is also homomorphism (cf. th. 4.22 of [24]).

The construction is interesting basically because the sets without the first element become elements of the structure, thus we are going to be able to draw information from those sets by studying (subsets of) the semigroup $\mathbb{S}_{\varphi} := \pi^{-1}(\{[\varphi]\})$, when $[\varphi]$ is an idempotent element in $M[\mathcal{G}_k]$. Then some well-known results about βM as compact right topological semigroup can be applied. Although the reference to $\widehat{M[\mathcal{G}_k]}$ can be omitted throughout the paper by defining the subsemigroups only in terms of βM (as we do in fact below), we believe that the semigroup \mathbb{S}_{φ} has interest itself, as it will be commented in the closing remarks of this paper.

Specifically, we are interested in translating a basic property of \mathbb{N} related to the existence of relative primes in a set (i.e., elements of the set that can not be decomposed as a product of numbers from the set itself). Given $A \neq \emptyset$, let S(A) be the set of its lower bounds,

$$S(A) = \{x : \forall b \in A(x < b)\}$$

Given two sets A, B, it will be said that they define the same cut, $A \sim_S B$, if S(A) = S(B). Let A^* be the set of its relative *composite elements*,

$$A^* := (A \cdot A) \cap A = \{x \in A : \exists b, c \in A (x = b \cdot c)\}$$

A basic property of \mathbb{N} is that if $0, 1 \notin A$, then $A^* \approx_S A$. It is not hard to prove that $I\Delta_0$ is enough to prove that fact for any Δ_0 -definable set. However, the reciprocal (namely if that property characterizes $I\Delta_0$), is somewhat more complicated. We must justify that we can limit ourselves to validate the principle of minimization for sets with adequate internal arithmetic structure to be able to apply that property. In general terms, in this paper alternative axiomatizations of bounded induction theories are obtained through schemes that describe the following property, for any non-empty Δ_0 -definable set:

$$A \sim_S A^* \Rightarrow \{0, 1\} \cap A \neq \emptyset \tag{(\diamond)}$$

The necessary combinatorial properties will be demonstrated by analyzing some appropriate ultrafilters of the set \mathbb{S}_{φ} described above.

If we focus on the idea of exploiting the algebraic properties of the compactification βM , it is important to note that the behavior of the operations (which are extensions that of M) is very different from the original ones. For example, in [25] Hindman, Maleki and Strauss shown that if a and b be distinct positive integers, then the equation $u + a \cdot p = v + b \cdot p$ has no trivial solutions with $u, v \in \beta \mathbb{N}$ and $p \in \beta \mathbb{N} \setminus \mathbb{N}$ (in that paper the abelian groups where the corresponding equation holds in their corresponding Stone-Čech compactifications are characterized). However, other interesting relationships have been proven (see e.g. the survey [24]). More details of this issue -once formalized- will be shown in the subsection 2.4. On the other hand, other results do establish a certain transfer of properties (for example that of [39]), although in general, it is not obtained straightforwardly. For example, for any sufficiently strong theory of arithmetic, the set of Diophantine equations provably unsolvable in the theory is algorithmically undecidable, as a consequence of the MRDP theorem. In contrast, in [27] the author proves the decidability of Diophantine equations provably unsolvable in Robinson's arithmetic Q. In that paper, the author axiomatizes the universal fragment of \mathbb{Q} itself.

1.2. Aim of the paper

Driven by the above motivations, some methods for analyzing fragments of Bounded Induction, which are not refinements of those used for classic fragments (those that were applied to $I\Sigma_n$ or $B\Sigma_n$) are presented. The proposal is developed in two stages. In the first part, several results about Δ_0 -definable sets will be established. The model will be submerged in a more complex structure in the second part (the aforementioned Stone-Čech compactification), where techniques related to algebraic combinatorics will be applied, taking into account the Δ_0 -definability of the sets involved in the study.

1.3. Structure of the paper

As we discussed, the paper has a differentiated structure in two related parts. The first part is dedicated to the analysis of segments and cuts in weak induction models, to provide these sets with an algebraic structure (semi-group). In the second part, the construction will be used to submerge the model in a topological structure (compact topological semigroup) where results of a very powerful combinatorial nature are available. When applied, it will provide information on the structure of the model as well as alternative axiomatizations of weak induction.

Specifically, the content of the paper is as follows. Section 2 is devoted to listing the main elements and results used in the paper. Next, it is studied the so-called *T*-kernel, the biggest initial segment of a model of a weak theory which is a model of T, when $T = IE_n$, studying its approximation by means algebraic initial segments as well as its definability (Sect. 3). Section 4 presents how to show the maximal character of $I\Delta_0$ as possible Π_1 -interpretable theory in RIE $_0^1$ (3.18).

The second part of the paper starts presenting an algebraic (sec. 5) and topological (sec. 4) treatment, including Ramsey theoretic arguments, of the *upper neighborhood* of a bounded definable set without the first element. This study allows to give alternative axiomatizations of $I\Delta_0$ representing the property (\diamond) described above, as well as to state combinatorial regularities near to the kernel.

2. Preliminaries and some starting results

In this section, a number of notions and results to be used throughout the paper will be introduced. Specifically, some results on $I\Delta_0$, IE_k ($k \ge 1$), and $IE_0 = IOpen$. Also, two types of initial segments are defined, and finally, a brief introduction on βM (being M a semigroup) will be shown, focusing on its algebraic and topological features.

2.1. Subsystems of $I\Delta_0$

The language of the Arithmetic is $L = \{+, \cdot, <, 0, 1\}$. Quantifications like $\exists x (x < t \land \varphi)$ or $\forall x (x < t \rightarrow \varphi)$ (denoted by $\exists x < t\varphi$ and $\forall x < t\varphi$) are called *bounded*.

The set Δ_0 is the set of bounded formulas (which have only bounded quantifiers, respectively). Similar to the definitions of the classical fragments in Peano Arithmetic PA, the classes E_n , U_n are defined according to the alternation of (bounded) quantifiers:

$$\begin{split} & \mathcal{E}_0(=\mathcal{U}_0) \text{ is the class of open formulas,} \\ & \mathcal{E}_{n+1} := \{ \exists x_1 < t_1 \cdots \exists x_m < t_m \varphi \mid \varphi \in \mathcal{U}_n \}, \text{ and} \\ & \mathcal{U}_{n+1} := \{ \forall x_1 < t_1 \cdots \forall x_p < t_p \varphi \mid \varphi \in \mathcal{E}_n \} \end{split}$$

(with t_i *L*-terms). Therefore $\Delta_0 = \mathbf{E}_{\omega} = \bigcup_k \mathbf{E}_k$.

The (finite) theory PA⁻ has as models the nonnegative part of discretely ordered rings.

The paper will work on the well–known Induction (I), Least element (L) and Collection (B) schemas, and the theories

$$\mathbf{E}\Gamma = \mathbf{P}\mathbf{A}^- + \{\mathbf{E}_{\varphi} \mid \varphi \in \Gamma\}$$

where E = I, L, B and $\Gamma = E_n$, U_n (see [30] and [51]). It is unknown whether $\{IE_n\}_{n < \omega}$ collapses. Only partial results are known. Among them, the following ones:

- $IE_0 \equiv LE_0$, $IE_0 \nvDash IE_1$ ([43]).
- IE_0 plus the full collection scheme does not prove IE_1 [3].
- $\operatorname{IE}_{n+1} \iff \operatorname{LU}_n \Longrightarrow \operatorname{IE}_n \iff \operatorname{IU}_n \iff \operatorname{LE}_n (n \ge 1; \operatorname{see} [28, 51]).$

Another useful scheme for the paper is the following one:

$$\mathbf{G}_{\psi}(x) := \forall \vec{v} < x [\exists u < x \psi(u, \vec{v}) \rightarrow \exists u < x(\psi(u, \vec{v}) \land \forall w < u \neg \psi(w, \vec{v}))].$$

It is easy to see that

$$\mathrm{PA}^- + \{ \forall u \mathrm{G}_{\psi}(u) : \psi \in \mathrm{E}_k \}$$

is an alternative axiomatization of IE_k .

2.2. Open induction and subsystems

If models of IE₀ are considered, then a large number of results about their algebraic nature can be proved. A seminal paper on this topic was [43], where J. C. Shepherdson shows the following (recursive) nonstandard model $\mathbb{N}_X \models \mathrm{IE}_0$. Let K be the field of real algebraic numbers and X an indeterminate. The universe of \mathbb{N}_X is the set of polynomials

$$\{a_p X^{p/q} + \dots + a_1 X^{1/q} + a_0 : a_p, \dots, a_1 \in K, a_p > 0, a_0 \in \mathbb{Z}, p \in \mathbb{N}, q \in \mathbb{N}^*\},\$$

the addition and multiplication are naturally defined, and the order "<" is interpreted as: P(X) < Q(X) if the leader coefficient of Q(X) - P(X) is > 0.

Shepherdson's construction will be generalized (in the first part of the paper) to some kind of extensions and models. The most direct one is to replicate it to obtain, given any $M \models IE_0$, an end-extension $M \subset_e M_X \models IE_0$ (The model thus constructed will be called **Shepherdson's extension of** M).

An algebraic characterization of IE_0 (also proved in [43]) is the following:

Result 2.1. The models of IE_0 are those models of PA^- that are the integer part of their real algebraic closure.

See [21] for other analogous results that provide algebraic theories equivalent to various fragments of IE₀.

Please note also that Shepherdson's algebraic construction produces a recursive model if a recursive one is extended. Thus it is not useful to build (extensions of) of models of IE_1 , due to *Tennenbaum phenomena* in IE_1 :

Result 2.2. There are no recursive nonstandard models of IE_1 [51].

Shepherdson's work illustrates how the construction of real closed fields can be related to obtaining adequate extensions of the fraction field of an arithmetic model [12]. See also the more recent paper [46], where Tanaka and Tsuboi present a new construction of real closed fields using an elementary extension of an ordered field that has an integer part that satisfies PA. It is also interesting to review the relationship obtained by D'Aquino, Knight and Strachencko in [11], where it is shown that the real closure of a model I Σ_4 is recursively saturated, thus establishing a relationship between a model-theoretic property of (ordered) fields and classical induction schemes on Arithmetic. Alternatively, it is possible to study the inverse path, that is, to analyze the real closed fields that admit an integer part whose non-negative part is a model of PA (see [6]).

Throughout the paper, two remarkable subsystems of IE_0 are considered, which will be used as base theories. The theory IE_0^1 augments PA^- with the induction scheme for open formulas whose atomic subformulas are linear inequalities. The theory is equivalent to

$$PA^{-} + \{ \exists u (vu \le x < v(u+1)) \}$$

(i.e. PA^- plus Euclidean division), and it is strictly weaker than IE₀. The theory of integer roots,

$$RIE_0^1 := IE_0^1 + \{ \exists u (u^k \le x < u^{k+1}) : k < \omega \}$$

is also weaker than IE₀. Moreover, the extension $IE_0^1 \subset RIE_0^1$ is proper (all these results can be found in [5]).

2.3. Exponentiation and segments determined by (standard) powers or integer roots

Let us suppose $PA^- \subseteq T$ and $\mathbb{N} \models T$. A function $f : \omega^n \to \omega$ is provably recursive in T if there is $\theta(\vec{x}, y, z) \in \Delta_0$ such that $\exists z \theta(\vec{x}, y, z)$ defines the graph of f and $T \vdash \forall \vec{x} \exists y \exists z \theta(\vec{x}, y, z)$. Such functions have been characterized for strong fragments (cf. [22]).

According to Result 1.1 the function $(x, y) \mapsto x^y$ is not provably recursive in $I\Delta_0$, although there is a bounded formula $\eta(x, y, z) \in \Delta_0$ that represents this function in $I\Delta_0$ and such that the basic properties are provable in $I\Delta_0$ (see e.g. [20]). To simplify the following notation is introduced: $x^y = z$ will mean $\eta(x, y, z)$, and exp is the sentence $\forall x \forall y \exists z (x^y = z)$.

Remark 2.3. Throughout the paper, k_e will be an integer such that IE_{k_e} proves the basic properties of $x^y = z$ as well as the weak overspill principle for U_{m+1} -formulas, for some m such that $x^y = z \in E_m$.

Two kinds of segments are used in this paper. The segment determined by standard powers of $a \in M \models$ PA⁻ is

$$a^{\omega} := \{ b \in M : \exists n < \omega(b < a^n) \}$$

It is straightforward to prove that

$$a^{\omega} = \bigcap \{ N \subseteq_e M : N \models PA^- \text{ and } a \in N \}$$

When $a \notin \omega$, it is defined the segment bounded by the integer standard roots of a as

$$a^{\frac{1}{\omega}} := \{ b \in M : \forall n < \omega(b^n < a) \}$$

for which it can be shown that

$$a^{\frac{1}{\omega}} = \bigcup \{ N \subset_e M : N \models PA^- \text{ and } a \notin N \}.$$

Notation: As usual, it will be written $\varphi(M)$ instead of $\{a \in M : M \models \varphi(a)\}$, and L(M) denotes the language L augmented by constants denoting the elements of M. To simplify notation, it will be written " $M \models T_1 + \neg T_2$ " to mean " $M \models T_1$ and $M \not\models T_2$ ".

2.4. Algebra on Stone-Čech compactification of a (discrete) semigroup

Let us recall that Stone-Čech compactification βS of a discrete space S is the compact Hausdorff space whose elements are the ultrafilters on S (being the elements of S identified with the principal ones). A basis of the topology is given by the (clopen) sets

$$\overline{A} = \{ p \in \beta S : A \in p \}$$
 for each $A \subseteq S$

The set S is dense in βS . It will be denoted by S^* to the set $\beta S \setminus S$.

The construction is functorial in nature, i.e. for each function $f: S \to S$ induces a (unique) extension $\beta f: \beta S \to \beta S$ such that the following diagram is commutative:

$$\begin{array}{cccc} S & \stackrel{f}{\to} & S \\ i & \downarrow & \downarrow & i \\ \beta S & \stackrel{\beta f}{\to} & \beta S \end{array}$$

The extension βf is defined by

$$\beta f(p) := \{ A \subseteq S : f^{-1}(A) \in p \}$$

The operation * of a discrete semigroup $(M, *_M)$ is extended to βM as follows (cf. [24], chap. 4, sect. 1): If $p, q \in \beta M$, the ultrafilter p * q is defined by:

$$A \in p * q$$
 if and only if $\{x : \{y : x * y \in A\} \in q\} \in p$.

The structure $(\beta M, *)$ is a (compact) right-topological semigroup. An excellent reference for this topic is the exposition [24].

The idempotent elements of such a semigroup are especially interesting for our purposes. A result that assures its existence is the following:

Result 2.4. [17] Every compact right topological semigroup has an idempotent element.

The interest in these elements lies in their combinatorial reading, since they determine a certain Ramseylike regularity. Let us formalize such a property. A *FP*-set in a semigroup (S, \cdot) is a set which contains all the finite products of a set, i.e. a subset like

$$FP(\{x_n\}_{n<\omega}) := \{\prod_{n\in F} x_n : F \subset \omega \text{ and } |F| < \aleph_0\},\$$

for some $\{x_n\}_{n<\omega}$ sequence (in additive notation, $FS(\{x_n\}_{n<\omega})$). A version of Galvin-Glazer theorem of the Finite Sums Theorem, in terms of idempotents elements, is the following:

Result 2.5. (see e.g. Th. 5.8 of [24]) If $A \in p$ and p is a nonprincipal idempotent ultrafilter in $(\beta M, \cdot)$, then A is a FP-set.

In the second part of the paper the following property on the fixed points will be used.

Result 2.6. (see e.g. Th. 3.35 in [24]) Let $\beta f : \beta M \to \beta M$. If $p \in \beta M$ is a fixed point for βf then

$$\{x \in M : f(x) = x\} \in p$$

For example, consider $\beta f : \beta M \to \beta M$ (being $M \models PA^-$) be the extension of $f(x) := x^2$. By above result it has not any fixed point in $\beta M \setminus \{0, 1\}$.

It is also possible to generalize the composition of βf as follows: given a ultrafilter $r \in \omega^*$ (recall that $\omega^* = \beta \omega \setminus \omega$), it is defined the *r*-iteration of βf as

$$(\beta f)^r(p) = q \quad \Longleftrightarrow \quad \{n : (\beta f)^n(p) = q\} \in r$$

this type of iteration is very useful because of the following property:

Theorem 2.7. (folklore, see e.g. [2]) Let (X, f) be a topological dynamic system (thus X is compact). The following conditions are equivalent:

1. $x \in X$ is recurrent (that is, for any U neighborhood in X the set $\{n \in \omega : f^n(x) \in U\}$ is infinite) 2. There exists $r \in \omega^*$ such that $f^r(x) = x$

Therefore, since recurrent points exist, there also exists $r \in \mathbb{N}^*$ such that $f^r(x) = x$ for some $x \in X$. Particularize us this fact for our case. If the dynamic system $(\beta M, \beta f)$ is considered, it can conclude that

Result 2.8. For any $f: M \to M$ there exist $r \in \mathbb{N}^*$ and $p \in \beta N$ such that $(\beta f)^r(p) = p$.

3. Part I: maximal segments, kernels and gaps

The initial segments that will be used in the rest of the article are studied in this section.

Definition 3.1. [3] Let $M \models PA^-$ and $T \subseteq Th(\mathbb{N})$. The *T*-kernel of *M* is

$$M(T) := \bigcup \{ I \subseteq_e M \mid I \models T \}$$

Since M(T) is the union of an increasing chain of initial segments, if T is Π_1 -axiomatizable (for example IE_n) then $M(T) \models T$.

Given a formula $\varphi(x, \vec{a})$, the notation

$$\varphi(x, \vec{a}) < u$$

(or simply $\varphi < u$) will be used to mean that $\exists v \leq u\varphi(v, \vec{a})$, and $u < \varphi$ means $\neg(\varphi < u)$.

12

Definition 3.2. Let $M \models PA^-$. An L(M)-formula $\varphi(x)$ is said a gap in M if $M \not\models L_{\varphi}$.

The following lemma is a very useful characterization of the elements that do not belong to a IE_k -kernel:

Lemma 3.3. Let $M \models PA^-$. If $a \in M \setminus M(IE_k)$ $(0 \le k \le \omega)$ then there exists $n \in \omega$ and ψ an E_k -gap in M with parameters in a^{ω} such that $M \models \psi < a^n$. Moreover, if $M \models IE_0^1$ then it can be selected ψ with parameters in [0, a].

Proof. If $a \notin M(\mathrm{IE}_k)$ then $a^{\omega} \not\models \mathrm{IE}_k$. Thus is $\psi(x, \vec{b}) \in \mathrm{E}_k$ with $\vec{b} \in a^{\omega}$, such that $a^{\omega} \not\models \mathrm{L}_{\psi}$.

The additional property for IE_0^1 -models is easy to justify, by using Euclidean division on the parameters. \Box

Remark 3.4. It is easy to prove, using the above lemma, that M(PA) is nonstandard in every Δ_0 -recursively saturated model of PA⁻. This fact is true because the following type can be realized:

$$\mathbf{p}(v) := \{ m < v \land \forall u < v^m \mathbf{G}_{\varphi_m}(u) : m \in \omega \}$$

(where $\{\varphi_m\}_{m<\omega}$ is a recursive enumeration of E₁). The element that realizes $\mathbf{p}(v)$ belongs to $N = M(\text{IE}_1)$, so this segment is nonstandard. According to Paris' result given in [42] (see also [29]), N(PA) (which coincides with M(PA)) is nonstandard.

In fact, it is sufficient to have U₃-recursive saturation.

From now on it will be studied some relationships between the *T*-kernels and the initial segments defined in Sect. 2.3. In the case of Shepherdson's model, \mathbb{N}_X is satisfied that

$$2^{\omega} = X^{\frac{1}{\omega}} = \mathbb{N} = \mathbb{N}_X(\mathrm{I}\Delta_0)$$

In general, if M is a model of a weak theory and $a \in M \setminus M(\mathrm{IE}_k)$, then $M(\mathrm{IE}_k) \subseteq_e a^{\frac{1}{\omega}}$. Likewise, if $a \in M(\mathrm{IE}_k)$, then $a^{\omega} \subseteq M(\mathrm{IE}_k)$. Therefore, it can be "approximated" the IE_k -kernels with that kind of initial segments. Such an approach can either achieve equality or at least to be arbitrarily approximate, by properly choosing elements. This fact suggests the following definition, where a name is given to the possible relationships between both types of segments.

Definition 3.5. Let $M \models PA^-$, and $0 \le k \le \omega$. It is said that:

- 1. *M* is *k*-short if there is $a \in M$ such that $a^{\omega} = M(IE_k)$. Otherwise it is said that *M* is *k*-long.
- 2. *M* is *k*-high if there is $a \in M$ such that $a^{\frac{1}{\omega}} = M(IE_k)$. Otherwise it is said that *M* is *k*-low.

Throughout the following sections, several relations between those kinds of segments and other problems on Weak Arithmetics will be proven (they are described in Fig. 2). Several of them will be used or refined in the second part of the paper.

By generalizing Shepherdson's construction, several examples of the models defined above can be shown (see Table 1):

- Shepherdson's extension of a model of IE_k is k-high $(1 \le k \le \omega)$, and \mathbb{N}_X is ω -short.
- A ω -low model of IE₀ is $M_{\downarrow} := \bigcup_{i < \omega} M_i$, where

$$M_n = (\cdots (\mathbb{N}_{X_n})_{X_{n-1}} \cdots)_{X_0}$$

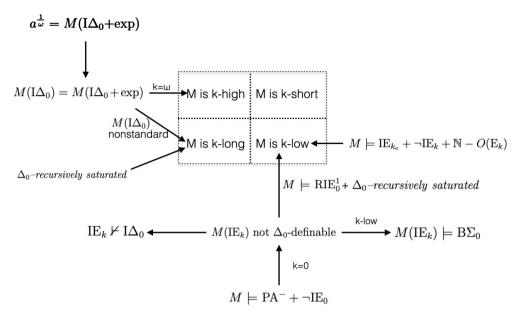


Fig. 2. Results about kernels (for models of PA⁻).

Table 1Examples of initial segments.

	k-high	k - low
k-short	\mathbb{N}_X	M_{\downarrow}
$k - \log$	$M_X, M \models I\Delta_0 + exp$	Corollary 3.9

 $(M_{\downarrow} \text{ is an integer part of its real closure because this closure satisfies } K^{M_{\downarrow}} \cong \bigcup_{i < \omega} K^{M_i}$, hence $M_{\downarrow} \models IE_0$, and $M_{\downarrow}(I\Delta_0) = \mathbb{N}$.

For stronger theories the following result holds:

Proposition 3.6. Let $M \models PA^-$ If

$$a^{\frac{1}{\omega}} = M(\mathrm{I}\Delta_0 + \exp)$$

for some a, then

$$M(I\Delta_0) = M(I\Delta_0 + \exp)$$

(thus M is ω -high and, if $M(I\Delta_0)$ is nonstandard, ω -long).

Proof. Otherwise, if $a \in M(I\Delta_0) \setminus M(I\Delta_0 + \exp)$, it is sufficient to apply Δ_0 -overspill in $M(I\Delta_0)$ on the Δ_0 -formula $\gamma(v) \equiv \exists u < a(v^v = u)$ to reach a contradiction. \Box

3.1. Definability of kernels

For stronger theories the following result holds: The problem of Δ_0 -definability of IE_k-kernels is related to the collection schema. It holds that if $M \models IE_{k+1} + \neg IE_{k+2}$, then $M(IE_{k+2}) \models BE_k$ [3].

Proposition 3.7. If $M(IE_k) \not\models B\Sigma_0$ then it is Δ_0 -definable.

Proof. Let $\varphi(x, y)$ be a L(M)-formula such that $M(IE_k) \not\models B_{\varphi}$, that is

$$M(\mathrm{IE}_k) \models \forall x \leq b \exists y \varphi(x, y) \land \forall z \exists x \leq b \forall y < z \neg \varphi(x, y)$$

then the formula $\exists x \leq b \forall y < z \neg \varphi(x, y)$ defines $M(IE_k)$ in M. \Box

In Δ_0 -recursively saturated models, properties of kernels and Δ_0 -definability issues are related. Please let us note that in bounded induction, Δ_0 -recursive saturation has a different behavior than in strong fragments (see e.g. [9], Sect. 2). It is known that if $a \in M(I\Delta_0)$ and $2^{a^{\omega}} \subsetneq M(I\Delta_0)$ then the relational structure [0, a] is recursively saturated [37]. Therefore, if $M(I\Delta_0 + exp) = M(I\Delta_0)$, then $M(I\Delta_0)$ is short recursively saturated. This kind of models presents nice properties (see e.g. [18] for an exposition) Let us first see some results on this type of model.

Theorem 3.8. Let $M \models PA^-$ be Δ_0 -recursively saturated and $0 \le k \le \omega$. It holds that:

- 1. M is k-long.
- 2. If $M \models \text{RIE}_0^1$ verifies that $M(\text{IE}_k)$ is Δ_0 -definable, then M is k-low. This fact is not true in general for models of IE_0^1 .

Proof. To prove (1), consider $\{\varphi_i : i \in \omega\}$ a recursive enumeration of E_k and $b \in M(IE_k)$. Then

$$\mathbf{p}(v;b) = \{v > b^n \land \bigwedge_{1 \le i \le n} \mathbf{G}_{\varphi_i}(v^m) : n, m < \omega\}$$

 $(v^m \text{ is the term } v \stackrel{(m)}{\cdots} v)$ is a finitely satisfiable type in M. Any element realizing $\mathbf{p}(v)$ belongs to $M(\mathrm{IE}_k) \setminus b^{\omega}$

To prove (2), assume that $a \in M \models \text{RIE}_0^1$ verifies $a^{\frac{1}{\omega}} = M(\text{IE}_k)$. If $\varphi \in \Delta_0$ defines $M(\text{IE}_k)$, then the type

$$\mathbf{p}(v;a) = \{\neg \varphi(v)\} \cup \{v^n < a : n \in \omega\}$$

is finitely satisfiable. Contradiction.

To conclude the proof, it remains to see that 3.8.(2) cannot be improved to IE_0^1 , by giving a countermodel: Let M_1 be the structure whose universe is

$$\{a_n X^n + \dots + a_1 X + a_0 : a_i \in \mathbb{Q}, a_n > 0, a_0 \in \mathbb{Z} \text{ if } n > 0 \text{ or } a_0 \in \mathbb{N} \text{ if } n = 0\}$$

The addition, product and order relation are defined as usual. It verifies that $M_1 \models \mathrm{IE}_0^1$ [5].

A first property of M_1 is that the set ω is defined by the open formula

$$\varphi(u) \equiv u^2 < X$$

Let M be a Δ_0 -recursively saturated elementary extension of M_1 . The following results hold:

- (i) For all $n \ge 2$ it is true that $M \models u^n < X \leftrightarrow u^2 < X$.
- (*ii*) For each $\psi(u, \vec{y}) \in \Delta_0$

$$M \models \forall \vec{y} \left[\exists u(u^2 < X \land \psi(u, \vec{y})) \to \exists v(\psi(v, \vec{y}) \land \forall w < v \neg \psi(w, \vec{y})) \right]$$

 $(iii) \ M \models \neg \varphi(u) \to \neg \mathbf{G}_{\neg \varphi}(u).$

From (i), (ii) and (iii) it is concluded that $X^{\frac{1}{\omega}} = M(I\Delta_0)$, but $M(I\Delta_0)$ is definable by $\varphi(u)$, an E_0 formula. \Box

Analogously countermodels for finite fragments of RIE_0^1 can be found, using similar models which are introduced in [5].

Corollary 3.9. Let $M \models IE_k$. Every Δ_0 -recursively saturated elementary extension M' of the Shepherdson extension M_X is k-low and k-long.

Proof. M' is k-long by Theorem 3.8(1). Moreover M is k-low as consequence of 3.8.(2): the formula

$$\varphi(u) \equiv u = 0 \lor (u|X \land (u+1)|X)$$

defines $M'(\mathrm{IE}_k)$ in M' (if $M \models \mathrm{IE}_k$). This is true because the following formulas are true in M' (they are true in M):

- $M' \models \neg \varphi(u) \rightarrow \neg \mathbf{G}_{\neg \varphi}(u)$
- $M' \models \varphi(u) \to \mathbf{G}_{\psi}(u)$ for each $\psi \in E_k$

The analysis made in Theorem 3.8 can be supplemented by outlining the hardness of deciding whether the reciprocal of 3.8.(2) holds: if such reciprocal fails, then $IE_k \nvDash I\Delta_0$.

Theorem 3.10. If there is a k-low model $M \models \text{RIE}_0^1$ with $M(\text{IE}_k)$ not Δ_0 -definable, then $\text{IE}_k \nvDash \text{I}\Delta_0$.

Proof. Suppose that there is such a model M and $\operatorname{IE}_k \vdash \operatorname{I}\Delta_0$. By Result 1.3 one would conclude that IE_k is finitely axiomatizable; so there are E_k -formulas $\varphi_1, \ldots, \varphi_n$ such that IE_k is equivalent to the theory

$$\mathrm{PA}^{-} + \{ \forall u \mathrm{G}_{\varphi_1}(u), \ldots, \forall u \mathrm{G}_{\varphi_n}(u) \}$$

Then the following formula defines $M(IE_k)$:

$$\gamma(u) := \bigwedge_{1 \le i \le n} \mathcal{G}_{\varphi_i}(u)$$

This is true because if $M \models \gamma(b)$, then $b^{\frac{1}{\omega}} \subsetneq M(\mathrm{IE}_k)$ (*M* is ω -low), so $b \in M(\mathrm{IE}_k)$. \Box

This theorem will be improved later (Corollary 3.24). Please let us note that the ω -low model M_{\downarrow} , described before would not be a candidate to apply the theorem above: $M_{\downarrow}(\text{IE}_k)$ is definable in M_{\downarrow} by

$$\psi(u) \equiv \forall v < u(v = 0 \lor \forall y \le v(v^2 \ne 2y^2))$$

which has only parameters in $M_{\downarrow}(IE_k) = \mathbb{N}$.

In Shepherdson's model $\mathbb{N}_X(\mathrm{IE}_1) = \mathbb{N}$ holds, being the segment definable by the formula $\psi(u)$ above. It is known that $\psi(u) \in \mathrm{E}_1$ in IE_0 (see [3]). It is interesting to note that this kind of characterization is not possible for $M(\mathrm{IE}_0)$ if it is only required to be a model of PA⁻:

Proposition 3.11. If $M \models PA^- + \neg IE_0$, then $M(IE_0)$ is not E_0 -definable in M with parameters in the own segment.

Proof. Let us assume otherwise, that there exists $\varphi(x, \vec{u}) \in E_0$ and $\vec{a} \in M(IE_0)$ such that $\varphi(x, \vec{a})$ defines $M(IE_0)$.

Let K^M be the real closure of M (by Result 2.1 M is not integer part of K^M , because $M \not\models IE_0$), and K' the real closure of $M(IE_0)$ within K^M . Since $M(IE_0)$ is a whole part of their actual closure, it verifies that

$$K' \models \forall x \exists y (x \ge 0 \to y \le x < y + 1 \land \varphi(y, \vec{a})) \tag{(*)}$$

By model completeness of the theory of real closed fields, $K' \prec K^M$, so K^M satisfies (*). Then $M(IE_0)$ would be an integer part in K^M , which is a contradiction. \Box

3.2. Collapse of kernels and interpretability

Once the definability of the kernels has been examined, the next concern to be studied is about two related questions. The first one asks whether there exists a model with internal structure like Shepherdson's model (that is, k-high and k-short for every k) but for *infinite and different* kernels, (that would imply that $I\Delta_0$ is not finitely axiomatizable). It will be seen that there is no such model under some assumptions (Corollary 3.15): in an ω -high and ω -short model the hierarchy of kernels $\{M(IE_k)\}_{k<\omega}$ collapses. Please note that this fact does not imply that $\{IE_n\}_n$ collapses, nor that $I\Delta_0$ is finitely axiomatizable modulo IE_0 , but it suggests the second question:

Open Problem 3.12. Is $I\Delta_0$ a theory Π_1 -interpretable in IE₀?

The aim will be to demonstrate on that issue that $I\Delta_0$ has certain maximal status as possible Π_1 -interpretable theory on RIE₀¹. In what follows it will be used the index k_e introduced in Remark 2.3.

Firstly it is going to be related to the collapse of $\{IE_n\}_{n<\omega}$ with a problem on Δ_0 -elementary extensions:

Theorem 3.13. Suppose that every model of IE_{k_e} has a Δ_0 -elementary extension model of RIE_0^1 , which is ω -high and ω -short. Then $\{IE_k\}_{k < \omega}$ collapses.

Proof. Let us see that $IE_{k_e} \vdash I\Delta_0$, being k_e any index satisfying the features required in Remark 2.3. All one needs to do is prove that if $N \models IE_{k_e}$ then $N \models I\Delta_0$.

Let N be such a model, and let $M \models \text{RIE}_0^1$ be a Δ_0 -elementary extension of N which is ω -high and ω -short.

Claim: $M(IE_{k_e}) = M(I\Delta_0).$

Proof of the claim: Assume that there is $a \in M(IE_{k_e}) \setminus M(I\Delta_0)$. Let $b \in M(I\Delta_0)$ and

$$\psi(u) \equiv \forall x < u \exists z < a(b^x = z \land z^x < a).$$

It holds that $M(IE_{k_e}) \models \psi(n)$ for each $n \in \omega$. By weak U_{m+1} -overspill, there is $\delta \in M(I\Delta_0) \setminus \omega$ such that $M \models \psi(\delta)$. Consider now the model of PA⁻:

$$I = \{ d \in M : \exists \alpha \in \delta^{\frac{1}{\omega}} (d < b^{\alpha}) \}$$

It is straightforward to prove that $I \models PA^-$ and $b^{\omega} \subsetneq I \subseteq a^{\frac{1}{\omega}}$ $(M(I\Delta_0)$ is nonstandard, so $\delta^{\frac{1}{\omega}} \neq \omega$ because the model is ω -high). Therefore, $b^{\omega} \neq \mathcal{M}(I\Delta_0)$, so M is ω -long. Contradiction. \Box

The claim, together with the fact that $N \subseteq M(\mathrm{IE}_{k_e}) = M(\mathrm{I}\Delta_0)$, implies that for each $a \in N$ and $\varphi(x) \in \Delta_0$ it has $N \models \mathrm{G}_{\varphi}(a)$ (because $M \models \mathrm{G}_{\varphi}(a)$). Therefore, $N \models \mathrm{I}\Delta_0$. \Box

Notice that in the demonstration of Theorem 3.13, the following result has been demonstrated, which we believe to be of interest in itself.

Corollary 3.14. Let $M \models \text{RIE}_0^1$ be such that $M(\text{I}\Delta_0) = a^{\frac{1}{\omega}} = b^{\omega} \neq \mathbb{N}$ for some $a, b \in M$. Then $M(\text{IE}_k) = M(\text{I}\Delta_0)$ for some k.

An interesting question about the special status of IE₁ in the hierarchy $\{IE_n\}_n$, is whether Corollary 3.14 is true for k = 1, since for k = 0 does not hold. Let us note that, by Theorem 3.13, there is not $M \models IE_0$ with a stratification like Shepherdson's model for infinite kernels.

Corollary 3.15. There is $k \ge 1$ such that for each ω -high $M \models \text{RIE}_0^1$ with $M(\text{I}\Delta_0)$ nonstandard, one of the following conditions is true:

- 1. $M(\mathrm{IE}_k) = M(\mathrm{I}\Delta_0).$
- 2. M is ω -long.

Proof. Similar to that of 3.13, taking $k = k_e$. \Box

Please note that Corollary 3.15 is already independent of the question of whether $\{IE_n\}_{n<\omega}$ collapses. The plausible collapse of IE_k -kernels suggests that $I\Delta_0$ may be Π_1 -interpretable in a class of models of RIE_0^1 (recall that if $IE_k \vdash I\Delta_0$ then $I\Delta_0$ is finitely axiomatizable). This issue will be examined below.

Definition 3.16. A Π_1 -formula $\Phi(u)$ interprets $I\Delta_0$ in a theory T if Φ defines an initial segment which is a model of $I\Delta_0$, that is:

- $T \vdash \Phi(0) \land \forall x, y(x \le y \land \Phi(y) \to \Phi(x)).$
- $T \vdash \forall x, y(\Phi(x) \land \Phi(y) \to \Phi(x \cdot y) \land \Phi(x+1) \land \Phi(x+y)).$
- For each $\psi \in \Delta_0$, $T \vdash \forall x (\Phi(x) \to G_{\psi}(x))$ (hence $T + \forall u \Phi(u) \vdash I \Delta_0$).

See [49] for some results about interpretability of theories in models of $I\Delta_0$.

The following result can be useful to analyze the question of whether $I\Delta_0$ is Π_1 -interpretable in RIE₀¹. (Please let us compare the next theorem with Th. 5.26 in chap. V of [22], as well as their proofs):

Theorem 3.17. Let T be a theory such that $\operatorname{RIE}_0^1 \subseteq \operatorname{T} \subseteq \operatorname{I}\Delta_0$ and $M \models T$. Assume that $\Phi(u)$ is a Π_1 -formula defining a proper initial segment closed by multiplication in M. Then $M(\operatorname{I}\Delta_0) \models \forall x \Phi(x)$.

Proof. Assume $\Phi(u) \equiv \forall x \psi(x, u)$, with $\psi \in \Delta_0$, and let $N = M(\forall u \Phi(u))$. If $M(I\Delta_0) \subseteq N$, the result holds. Suppose that $N \subseteq M(I\Delta_0)$. Let M' be a Σ_1 -recursively saturated elementary extension of M, and $N' = M'(\forall u \Phi(u))$. Then $N' \subseteq M'(I\Delta_0)$.

Let $\gamma(u) \equiv \forall x, y < u\psi(x, y)$. We claim that

$$M' \models \gamma(0) \land \forall u(\gamma(u) \to \gamma(u+1))$$

If $a \in M'$ is such that $M' \models \gamma(a)$, then $a^{\frac{1}{\omega}} \subseteq N'$ (because $a^{\frac{1}{\omega}} \models \forall u \Phi(u)$). By Σ_1 -recursive saturation, it can to conclude that $a^{\frac{1}{\omega}} \neq N'$. Thus, $a \in N'$, hence $a + 1 \in N'$. Therefore, $M' \models \gamma(a + 1)$.

By induction on $\gamma(u)$, $M'(\mathrm{I}\Delta_0) \models \forall u\gamma(u)$. This implies that $N = M(\mathrm{I}\Delta_0)$ (because $M(\mathrm{I}\Delta_0) \setminus N \subseteq N(M', \mathrm{I}\Delta_0) \setminus N'$). \Box

The section will be completed by demonstrating a necessary and sufficient condition on the Π_1 -interpretability of $I\Delta_0$ in RIE_0^1 which is consequence of above theorem.

Corollary 3.18. The following conditions are equivalent:

- 1. $I\Delta_0$ is finitely axiomatizable modulo RIE_0^1 .
- 2. There exists an extension of $I\Delta_0$ which is Π_1 -interpretable in RIE¹₀.

Proof.

(1) \Longrightarrow (2): If $I\Delta_0 \equiv RIE_0^1 + \forall u\Phi(u)$, let

$$\Phi'(x) \equiv \Phi(x) \land \forall u(u < x \to \Phi(u)) \land \forall u < x \forall v < x \Phi(u \cdot v).$$

Since $M \models \operatorname{RIE}_0^1$ it must be

$$M(\forall u\Phi'(u)) \subseteq M(\forall u\Phi(u)) \subseteq M(\mathrm{I}\Delta_0)$$

then the formula $\Phi'(x)$ interprets an extension of $I\Delta_0$ in RIE_0^1 . (2) \Longrightarrow (1): Let $\Phi(x)$ the Π_1 -interpretation of an extension T of $I\Delta_0$. If $M \models RIE_0^1$, then

$$M(\forall u\Phi(u)) \subseteq M(\mathrm{I}\Delta_0)$$

and, reasoning as in Theorem 3.17, it follows that the $I\Delta_0$ -kernel is defined by $\Phi(u)$ in every model of RIE_0^1 . Therefore

$$\operatorname{RIE}_0^1 + \forall u \Phi(u) \equiv \mathrm{I}\Delta_0 \quad \Box$$

3.3. Arithmetic of gaps and segments

Several technical results about an extension of the model, where the gaps are considered as elements, are described in this section. A convenient way to do this is to consider them as if they were done with Henkin's cuts. This section, will introduce how arithmetic is in that extended model (see [32,33] for a study with a broader perspective of the idea). On the topology induced by cuts, in the case of PA there are studies on its behavior (see e.g. [31] where indicators are used).

Definition 3.19. Let $\varphi(x, \vec{a})$ and $\psi(x, \vec{b})$ be L(M)-formulas. It is said that $\varphi(x, \vec{a})$ and $\psi(x, \vec{b})$ define the same gap, $M \models \psi \sim \varphi$, if

$$M \models \forall u \forall v (u < \varphi < v \leftrightarrow u < \psi < v)$$

Clearly, '~' is an equivalence relation. It will be denoted by $[\varphi]$ the equivalence class of φ . Whilst it is understood that the definitions in this section are intended for gaps, note that they are also suitable for Δ_0 -definable sets with a first element, by identifying each $a \in M$ with $[\varphi_a]$ where $\varphi_a(x) := a \leq x$. Thus,

If
$$M \models (\mu x)\varphi(x) = a$$
 then $M \models a \sim \varphi$

For convenience it will be written [a] or simply a instead of $[\varphi_a]$.

Given $X \subseteq M$, it is defined

$$\mathcal{G}_k(M, X) := \{ [\varphi] : \varphi \text{ is a } E_k \text{-gap with parameters in } X \}$$

and it will be written $\mathcal{G}_k(M)$ if X = M. It is possible to extend the order \langle_M to $\mathcal{G}_k(M, X)$ as follows:

$$M \models \varphi < \psi$$
 if and only if $M \models \exists x(\varphi(x) \land x < \psi)$

The segment determined by a formula φ is defined as

$$S(\varphi) := \{ a \in M : M \models a < \varphi \}$$

A useful simplification is to work only with *coinitial* gaps. A formula $\varphi(x)$ is coinitial if $S(\varphi)$ agree with $\neg \varphi(M)$. It is obvious that every gap is equivalent to a coinitial gap.

Please note that if $X \not\subseteq S(\varphi)$, then the formula $\varphi < \psi$ is Δ_0 . It is also defined

$$(\varphi, b) := \{ x : M \models \varphi < x \land x < b \}$$

Let $\varphi, \psi \in \mathcal{G}_k(M, X)$, with $1 \le k \le \omega$. The addition and the product of two gaps are defined as

$$(\varphi * \psi)(x) := \exists u, v \le x(\varphi(u) \land \psi(v) \land x = u * v) \qquad (* \in \{+, \cdot\})$$

Note that trivially

$$M \models \varphi_a * \varphi_b \sim \varphi_{a*b}$$

It is easy to see that both operations are well defined on equivalence classes, extending that of operations in M as well as that they are associative and commutative $(1 \le k \le \omega)$. Therefore two semi-group structures are available for endowing $\mathcal{G}_k(M, X)$. We are interested in the *idempotent* gaps.

Definition 3.20. Let $* \in \{+, \cdot\}$. A gap $\varphi \in \mathcal{G}_k(M, X)$ is **idempotent** with respect to * (briefly, is an id(*)-gap) if $M \models \varphi * \varphi \sim \varphi$.

That is, $[\varphi]$ is an idempotent element in the semigroup $(\mathcal{G}_k(M,X),*)$. It will be denoted by $I_k^*(M,X)$ the set of idempotents for * and parameters in X (if X = M it will be written $I_k^*(M)$).

It is easy to write a formula $Idemp_*(\varphi)$ expressing that φ is a id(*)-gap in PA⁻. In the case of the theory RIE_0^1 , the following formula can be used:

$$Idemp_{\bullet}(\varphi) \equiv \forall x, y(x < \varphi \land y < \varphi \to x \cdot y < \varphi)$$

and, in IE_0^1 , $Idemp_+(\varphi)$ is equivalent to a similar formula, due to the following result:

Proposition 3.21. Let $M \models PA^-$ and let $\varphi \in \Delta_0$.

- If M ⊨ IE¹₀, φ is an id(+)-gap if and only if S(φ) is closed by +.
 If M ⊨ RIE¹₀ then φ is an id(·)-gap if and only if S(φ) ⊨ PA⁻. This fact is not true for models of IE¹₀ in general.

Proof. Proof of (2): Assume $\varphi(x)$ is idempotent and $a, b \in S(\varphi)$ such that $M \models \varphi < a \cdot b$. Then there exists c such that $M \models \varphi(c) \land c < a \cdot b$. Let $d, e \in M$ such that

$$M \models \varphi(d) \land \varphi(e) \land d \cdot e < c$$

Then one of $\{d, e\}$ is less than one of $\{a, b\}$; suppose d < a for example. Thus $\varphi < a$, a contradiction.

20

Conversely, suppose that $\varphi(x)$ and $a \in M$ such that $M \models \varphi < a$. If $S(\varphi) \models PA^-$, then $M \models \varphi < \lfloor \sqrt{a} \rfloor$. Therefore there exists $c \in M$ such that $M \models \varphi(c) \land c < \lfloor \sqrt{a} \rfloor$. Thus

$$M \models (\varphi \cdot \varphi)(c^2) \wedge c^2 < a$$

The proof of (1) is similar, replacing \cdot by +, and $\lfloor \sqrt{a} \rfloor$ by $\lfloor \frac{a}{2} \rfloor$.

Finally it will show that 3.21.(2) fails for a model of IE_0^1 : Let $M_1 \models \text{IE}_0^1$ the model defined in Theorem 3.8.(2), and $\varphi(u) \equiv X \leq u^2$. It holds that $S(\varphi) \models \text{PA}^-$ and

$$(\varphi \cdot \varphi)(u) \equiv \exists u_1 u_2 < u(X \le u_1^2 \land X \le u_2^2 \land u = u_1 \cdot u_2).$$

Since it is verified that $M_1 \models (\varphi \cdot \varphi)(u) \rightarrow u > X$ and $M_1 \models \varphi(X)$, then $M_1 \models \varphi < \varphi \cdot \varphi$. \Box

The following result, essentially due to J. Paris [41], shows that it is possible to restrict ourselves to idempotent gaps.

Theorem 3.22. Let $M \models PA^-$.

- 1. If $M \models IE_0^1$, $M \models I\Delta_0$ if and only if M has no bounded id(+)-gap.
- 2. If $M \models \text{RIE}_0^1$, $M \models \text{I}\Delta_0$ if and only if M has no bounded $id(\cdot)$ -gap.

The proof of 3.22.(1) is based in the scheme

$$\mathbf{I}_{\varphi}^{+} := \varphi(0) \land \varphi(1) \land \forall x, y(\varphi(x) \land \varphi(y) \to \varphi(x+y)) \to \forall x \varphi(x)$$

It verifies that $I\Delta_0 \equiv I^+\Delta_0$ ([41], lemma 2). The fact 3.22.(2) is proved like th. 3 of [41] (using for example the set

$$\{2^a : M \models \exists y 2^a = y\}$$

instead of the enumeration of the roots of the equation $x^2 + y^2 - 2axy - 1 = 0$).

Corollary 3.23. If $M \models \text{RIE}_0^1$ is k-high, then $M(\text{IE}_k)$ is Δ_0 -definable.

Proof. Suppose that $a^{\frac{1}{\omega}} = M(\mathrm{IE}_k)$. Since $a^{\omega} \not\models \mathrm{IE}_k$, by 3.22 there exists $\varphi \in \mathrm{E}_k$ an $id(\cdot)$ -gap such that

$$M(\mathrm{IE}_k) < \varphi < a^{\omega}$$

By 3.21 it has $S(\varphi) \models PA^-$. Therefore the only possibility is that $S(\varphi) = M(IE_k)$. Thus the formula

$$\psi(x) \equiv x < \varphi$$

defines $M(IE_k)$. \Box

Using Corollary 3.23 it is possible to improve Theorem 3.10:

Corollary 3.24. If there exists $M \models \text{RIE}_0^1$ with $M(\text{IE}_k)$ not Δ_0 -definable, then $\text{IE}_k \nvDash \text{I}\Delta_0$.

Proof. If M were such a model, then it would be k-low by the previous corollary. So, by Theorem 3.10 it follows that $IE_k \nvDash I\Delta_0$. \Box

Proposition 3.25. Let $[\varphi], [\psi] \in I_k^*(M)$, and $M \models \operatorname{IE}_0^1$ if * = +, or $M \models \operatorname{RIE}_0^1$ if $* = \cdot$. Then

 $[\varphi]*[\psi]=max\{[\psi],[\varphi]\} \qquad *\in\{+,\cdot\}$

Above result contrasts with an old unsolved question on $\beta \mathbb{N}$ that clearly exposes the complexity of arithmetic in $\beta \mathbb{N}$: Are there $p, q, r, s \in \beta \mathbb{N}^*$ such that $p + q = r \cdot s$? ([48], Chapter 7, Question 39. See also [1] for other similar problems).

Both operations present similar properties under the proper base theory $(IE_1 \text{ or } RIE_0^1)$, so only the results for the product will be demonstrated. We only enunciate the corresponding ones for the addition.

Consider the structure $(M[\mathcal{G}_k], +, \cdot, \leq)$ whose universe is the set of equivalence classes of E_k -definable sets (identifying the element a with φ_a), and the symbols are interpreted as before. Note that $M[\mathcal{G}_k] \not\models \operatorname{PA}^$ in general, because $M[\mathcal{G}_k] \models \exists x(x = x + 1)$ when $M \not\models \operatorname{IE}_k$ (by Proposition 3.25). In fact, the relationship between $M[\mathcal{G}_k]$ and M can be determined by relativizing the formulas to M via the natural map $\varphi \mapsto \varphi'$ that relativizes to elements x such that $x \neq x + 1$.

Since $\mathcal{G}_k(M)$ is defined in $M[\mathcal{G}_k]$ by the formula 'x = x + 1', then for every L(M)-sentence φ ,

 $M \models \varphi$ if and only if $M[\mathcal{G}_k] \models \varphi'$

On the other hand, since $\mathbb{N} \subset_e M[\mathcal{G}_k]$, it is verified that $M[\mathcal{G}_k] \models \Sigma_1(\mathbb{N})$ (the true Σ_1 -sentences in \mathbb{N}). In general, in IE_k-models, if $Th(M[\mathcal{G}_k]) \vdash x \neq x + 1$ then $M \models IE_{k+1}$.

4. Part II: algebraic combinatorics in IE_k

The results of the first part will now allow us to both enunciate and prove properties in weak induction models that relate combinatorial properties and sets without a first element (gaps). The task will then be carried out concerning the existence of relative composite numbers in some sets, following the idea described in the introduction.

4.1. Combinatorics in gaps

The previous section has provided two semigroup structures on $M[\mathcal{G}_k]$. Therefore, one is in a position to apply the (combinatorial) results on topological semigroups to the new structures, to later describe the results in terms of the starting model M.

The next aim is to state several combinatorial properties about the *upper neighborhood* of an idempotent gap (that is, intervals like (φ, a)) in the semigroup $(M[\mathcal{G}_k], *)$ ($* \in \{+, \cdot\}$), formalizing the idea about prime numbers relative to a set that was described in the introduction.

For the study it is convenient to split the interval (φ, ∞) into $\varphi(M)$ and its complement within the same interval:

Definition 4.1. Let $\varphi(x)$ be an E_k -gap. The complement of φ is the formula $\overline{\varphi}(x) := \neg \varphi(x) \land \varphi < x$.

The definition of composite number related to a set (*relative composite number*) would be defined as follows:

Definition 4.2. Let $\varphi \in E_k$ and $M \models PA^-$. An element $a \in M$ is φ -composite with respect to $* \in \{+, \cdot\}$ if $M \models \varphi_{comp*}(a)$, where

$$\varphi_{comp*}(x) := \varphi(x) \land \exists u_1, u_2 \le x(x = u_1 * u_2 \land \varphi(u_1) \land \varphi(u_2))$$

That is

$$\varphi_{comp*}(M) = (\varphi(M) * \varphi(M)) \cap \varphi(M)$$

The translation of arithmetical properties (e.g. prime, irreducible elements) to Stone-Čech compactification is not direct and it raises very interesting questions [45]. Please observe that, in general,

$$M \not\models \varphi_{comp*} \sim \varphi * \varphi$$

If $0, 1 \notin \varphi(M)$ and $\varphi(M)$ has a first element, then the set $\varphi_{comp*}(M)$ is not coinitial with φ (nor $\overline{\varphi}$)composite elements, because the first element of $\varphi(M)$ does not satisfy φ_{comp*} , that is, is a prime relative. Analogously with $\overline{\varphi}_{comp*}$. Therefore, the following translation of the arithmetical property (\diamond) discussed in the introduction is satisfied.

Proposition 4.3. Assume $M \models IE_0^1$ and φ is an L(M)-formula. Then

$$M \models \neg \varphi(0) \land \neg \varphi(1) \land \exists x \varphi(x) \to (\mathcal{L}_{\varphi} \to \varphi_{comp*} > \varphi \land (\overline{\varphi})_{comp*} > \varphi) \qquad (**)$$

Proof. Trivial.

Please observe that for proving the reciprocal of Proposition 4.3 (namely, if M verifies $A \approx_S A^*$ for any nontrivial Δ_0 -definable set A, then is a model of $I\Delta_0$) can not be straightforwardly proved. The existence of relative composite elements arbitrarily close to the gap needs to be demonstrated (actually it is just needed to state this feature for idempotent Δ_0 -gaps). The property can be demonstrated by showing that $\varphi(M)$ satisfies some Ramsey type regularity for a certain kind of gap φ .

4.2. FP and FS-sets within gaps

The regularity feature required to prove the reciprocal of Proposition 4.3 is stated in Theorem 4.4, which is a version of the finite sums theorem (Theorem 4.4). Its proof is inspired by that of Galvin-Glatzer's proof of the classic result on FP-sets, namely by exploiting the algebraic structure of βM . Specifically, it is shown that given φ a gap, this or its complement contains FP-sets arbitrarily close to $[\varphi]$.

Theorem 4.4. Let $M \models IE_0^1 + \neg IE_k$, and let φ be an idempotent E_k -gap with respect to \cdot . Then there exists $\psi \in \{\varphi, \overline{\varphi}\}$ such that $\psi(M) \cap (\psi, c)$ is a FP-set for arbitrarily $c > \varphi$.

Proof. Consider the closed (hence compact) nonempty set in βM (recall that, in the Stone-Čech compactification, \overline{A} denotes the closure of the set $A \subseteq M$ in βM)

$$\mathbb{H}_{\varphi} := \bigcap_{\varphi < a} \overline{\{b \ : \ M \models \varphi < b \land b < a\}}$$

We claim that

<u>Claim</u> \mathbb{H}_{φ} is a subsemigroup of βM .

<u>Proof of the Claim</u>: It is necessary to prove that $p \cdot q \in \mathbb{H}_{\varphi}$ for any $p, q \in \mathbb{H}_{\varphi}$. That is, that $(\varphi, a) \in p \cdot q$, for all $a > \varphi$; or equivalently,

$$\{x : \{y : x \cdot y \in (\varphi, a)\} \in q\} \in p$$

Write $x^{-1}(\varphi, a) = \{y : x \cdot y \in (\varphi, a)\}$. Let us observe that $x^{-1}(\varphi, a) \in q$ if and only if $\varphi < \lfloor \frac{a}{r} \rfloor$:

If $\varphi < \lfloor \frac{a}{x} \rfloor$, then $(\varphi, \lfloor \frac{a}{x} \rfloor) \in q$ hence $x^{-1}(\varphi, a) \in q$. Otherwise $\lfloor \frac{a}{x} \rfloor < \varphi$, then

$$(\varphi, a) \cap x^{-1}(\varphi, a) = \emptyset$$

hence $x^{-1}(\varphi, a) \notin q$.

Therefore, it is sufficient to show

$$\{x \ : \ \varphi < \lfloor \frac{a}{x} \rfloor\} \in p$$

Since the set is convex and $p \in \mathbb{H}_{\varphi}$, the condition above is equivalent to the existence of $x > \varphi$ such that $\varphi < \lfloor \frac{a}{x} \rfloor$. Let us see that such an element exists: Since $M \models \varphi \sim \varphi \cdot \varphi$, there is d < a such that

$$M \models (\varphi \cdot \varphi)(d) \land d < a$$

so $M \models \varphi(d_1) \land \varphi(d_2) \land d_1 \cdot d_2 = d$ for some d_1, d_2 . If $x = d_1$, then $\lfloor \frac{a}{x} \rfloor > \lfloor \frac{d}{x} \rfloor = d_2$ and $M \models d_2 > \varphi$, because $M \models \varphi(d_2)$. end of claim proof

According to Result 2.4, it is known that \mathbb{H}_{φ} contains an idempotent element p. Given $c > \varphi$, since

$$\{(\varphi, c) \cap \varphi(M), (\varphi, c) \cap \overline{\varphi}(M)\}\$$

is a partition of $(\varphi, c) \in p$, there is $\psi \in \{\varphi, \overline{\varphi}\}$ such that

$$\psi(M) \cap (\psi, c) \in p$$

hence this set is a FP-set.

It only remains to prove that the choice of ψ is valid for any element $c' > \varphi$, (that we can suppose c' < c). That fact is true, because necessarily $(\varphi, c') \cap \psi(M) \in p$ (otherwise we would have $(\varphi, c') \cap \overline{\psi}(M) \in p$, but this set would disjoint with $\psi(M) \cap (\varphi, c)$, that is not possible). \Box

Please observe that it is easy to see that $\pi : \beta M \to M[\mathcal{G}]$ of Fig. 1 verifies $\pi[\mathbb{H}_{\varphi}] = \{[\varphi]\}$, as well as the converse of the above result: If $\varphi(M)$ contains FP-sets arbitrarily near the gap, then it is an $id(\cdot)$ -gap. Without any essential modification, a similar result would be obtained for id(+)-gaps can be proved:

Theorem 4.5. Let $M \models IE_0^1 + \neg IE_k$, and let $\varphi(x) \in E_k$ be an id(+)-gap. There is $\psi \in \{\varphi, \overline{\varphi}\}$ such that $\psi(M)$ contains a FS-set arbitrarily near the gap; that is, for any $c > \varphi$, $\psi(M) \cap (\psi, c)$ is a FP-set.

A consequence of the above theorems was remarked above (directly provable by Theorem 4.4, by using Theorem 3.22):

Corollary 4.6. Suppose that φ be a Δ_0 -formula in the language L(M) such that

$$M \models Idemp_*(\varphi) \land \neg \varphi(0) \land \neg \varphi(1)$$

and $M \models \operatorname{IE}_0^1 if * = + or M \models \operatorname{RIE}_0^1 if * = \cdot$. If $M \models \neg \operatorname{L}_{\varphi}$, then

$$M \models \varphi_{comp*} \sim \varphi \lor (\overline{\varphi})_{comp*} \sim \varphi.$$

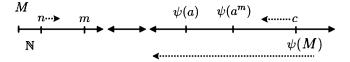


Fig. 3. Existence of arbitrary large powers in gaps (Theorem 4.10).

Therefore, by instantiating the result into a scheme that reflects that property alternative axiomatizations of $I\Delta_0$ are available. The following corollary shows an axiomatic characterization of $I\Delta_0$ using only the arithmetic properties that have been introduced: an idempotent set that is coinitial with its relative composite numbers contains either 0 or 1.

Corollary 4.7. The following theories are equivalent to $I\Delta_0$:

$$\begin{aligned} \operatorname{RIE}_{0}^{1} + \left\{ \exists x \varphi(x) \land Idemp_{\bullet}(\varphi) \land (\varphi_{comp} \sim \varphi \lor (\overline{\varphi})_{comp} \sim \varphi) \rightarrow (\varphi(0) \lor \varphi(1)) \right\}_{\varphi \in \Delta_{0}} \\ \operatorname{IE}_{0}^{1} + \left\{ \exists x \varphi(x) \land Idemp_{+}(\varphi) \land (\varphi_{comp+} \sim \varphi \lor (\overline{\varphi})_{comp+} \sim \varphi) \rightarrow (\varphi(0) \lor \varphi(1)) \right\}_{\varphi \in \Delta_{0}} \end{aligned}$$

Exploiting the ideas introduced in this subsection, richer combinatorial properties can be proved if we allow the model to fulfill basic additional properties, such as overspill for its standard part.

4.3. Combinatorics in presence of weak overspill on \mathbb{N}

This section will show some results from semi-group properties \mathbb{H}_{φ} (introduced in Theorem 4.4) for models which satisfies a weak form of overspill (see [15]):

Definition 4.8. Let $M \models PA^-$ and Γ a class of formulas. It is said that M has overspill Γ on \mathbb{N} , $M \models \mathbb{N} - O(\Gamma)$, if for all $\varphi(x, \vec{u}) \in \Gamma$, $\vec{c} \in M$ and $\beta > \omega$, if

$$\forall n \in \omega \exists m > n[M \models \varphi(m, \vec{c})]$$

then there is $\delta < \beta$ nonstandard such that $M \models \varphi(\delta, \vec{c})$.

This is a weak notion of overspill; any Δ_0 -recursively saturated model of PA⁻ has $\mathbb{N} - O(\Delta_0)$. Nevertheless, sometimes it is a good substitute for Δ_0 -recursive saturation; for example 3.8.(2) is true for models of IE_{ke} + $\mathbb{N} - O(\Delta_0)$, thus several results of Fig. 2 could be refined for these kinds of models.

Theorem 4.10 will show that, in models satisfying $\mathbb{N} - O(\Gamma)$, the regularity property can be refined by replacing the FP-sets with (standard) powers of an element. This will allow another alternative axiom of $I\Delta_0$ using the formula *exp* (Corollary 4.14).

Definition 4.9. A set $A \subseteq M \models PA^-$ is said to be **dense for standard powers** if

$$\forall c \ \forall n \in \omega \left[\left[0, c \right] \cap A \neq \emptyset \quad \Longrightarrow \quad \exists a < c \ \exists m > n(a, a^m \in A) \right]$$

In the following theorem the density for standard powers for gaps (Fig. 3) is stated:

Theorem 4.10. Let $M \models \text{RIE}_0^1 + \neg \text{IE}_k$, and $\varphi(x) \in E_k$ be an $id(\cdot)$ -gap. Then there exists $\psi \in \{\varphi, \overline{\varphi}\}$ such that $\psi(M)$ is dense in standard powers.

Proof. The function $f(x) := x^2$ has a continuous extension $\beta f : \beta M \to \beta M$. That function has not any fixed point in $\beta M \setminus \{0, 1\}$ by Result 2.6.

Since $f[(\varphi, \lfloor \sqrt{b} \rfloor)] \subseteq (\varphi, b)$, it follows that $(\varphi, b) \in \beta f(p)$, for all $p \in \mathbb{H}_{\varphi}$ and $b > \varphi$, so it defines the restriction

$$\beta f \upharpoonright_{\mathbb{H}_{\varphi}} : \mathbb{H}_{\varphi} \to \mathbb{H}_{\varphi}$$

Therefore it can be considered the dynamic system:

$$(\mathbb{H}_{\varphi}, \{(\beta f)^n\}_{n \in \omega^+})$$

on the compact space \mathbb{H}_{φ} ($\omega^+ = \omega \setminus \{0\}$).

By Result 2.7 there exist $p \in \mathbb{H}_{\varphi}$ and $r \in \beta(\omega^+)$ such that $(\beta f)^r(p) = p$.

That means that for all $A \in p$ that is, for all \overline{A} neighborhood of p

$$\{n \in \omega^+ : (\beta f)^n(p) \in \overline{A}\} \in r.$$

Since $(\beta f)^n = \beta(f^n)$ and $f^n[A] \in \beta(f^n)(p)$, it follows that

$$\{n \in \omega^+ : f^n[A] \cap A \neq \emptyset\} \in r.$$

Since r is not principal (because $0, 1 \notin \mathbb{H}_{\varphi}$),

$$\forall n \exists m > n[f^m[A] \cap A \neq \emptyset] \qquad (\dagger)$$

Let $\psi \in \{\varphi, \overline{\varphi}\}$ such that the set $A = \psi(M) \cap (\varphi, c)$ belongs in p. Then, by (\dagger) ,

$$\forall n \exists m > n \text{ such that } M \models \exists x < c[\psi(x) \land \psi(x^{2^m})]$$

 $(x^{2^m}$ denotes the term $x \stackrel{(2^m)}{\cdots} x$). Finally observe that the choice of ψ is true for all $c > \varphi$ (like in 4.4). \Box

We will use again the index k_e (see 2.3) to show, in models of $IE_{k_e} + \mathbb{N} - O(E_k)$, a nonlinear relation near the gap.

Definition 4.11. A set $A \subseteq M \models PA^-$ is said to be **dense for powers** if

$$\forall c \left[[0, c] \cap A \neq \emptyset \quad \Longrightarrow \quad \exists a \in A \exists \delta > \omega(a, a^{\delta} \in A) \right]$$

Corollary 4.12. Under the conditions of Theorem 4.10, if $M \models IE_{k_e} + \mathbb{N} \cdot O(E_k)$, then there is $\psi \in \{\varphi, \overline{\varphi}\}$ such that $\psi(M)$ is dense for powers.

Proof. Let $k > k_e$ Let $x^y = z$ be the formula defined in Sect. 2. It has that, for all $n \in \omega$,

$$\operatorname{IE}_{k_e} \vdash x \stackrel{(n)}{\cdots} x = y \to x^n = y$$

By 4.10, we have that for all $n \in \omega$ there is m > n such that

$$M \models \exists x, y < c(x^m = y \land \psi(x) \land \psi(y))$$

Since the formula is E_k $(1 \le k_e \le k)$, by applying $\mathbb{N} - O(E_k)$, the result is concluded. \Box

Corollary 4.13. If $M \models IE_{k_e} + \neg IE_k + \mathbb{N} - O(E_k)$, then M is k-low.

Proof. Suppose $M(IE_k) \subseteq a^{\frac{1}{\omega}}$. Let $\varphi \in E_k$ be a $id(\cdot)$ -gap such that

$$M(\mathrm{IE}_k) < \varphi < a$$

Reasoning like 4.12, there is $\psi \in \{\varphi, \overline{\varphi}\}$ such that

$$M\models \exists x < a(x^{2^m} < a \wedge \psi(x) \wedge \psi(x^{2^m}))$$

By $\mathbb{N} - O(\mathbb{E}_k)$, there is $b \in M$ and $\delta > \omega$ such that $M \models \psi(b) \land \psi(b^{\delta}) \land b^{\delta} < a$. Thus, $M(\mathbb{I}\mathbb{E}_k) < b < a^{\frac{1}{\omega}}$. \Box

The above results allow obtaining other alternative axiomatizations of $I\Delta_0$. Following the idea of the relative composite numbers, given $\varphi(x)$ it is defined the *exponential relative to* φ as the formula

$$\varphi_{exp}(x) := \varphi(x) \land \exists y, z < x[z \neq 1 \land y^z = x \land \varphi(y)]$$

Corollary 4.14. $I\Delta_0$ is equivalent to

$$\mathrm{IE}_{k_e} + \{ Idemp_{\bullet}(\varphi) \land \neg \varphi(0) \land \neg \varphi(1) \land \exists x \varphi(x) \to \varphi < \varphi_{exp} \land \varphi < (\overline{\varphi})_{exp} \}_{\varphi \in \Delta_{\mathbf{C}}}$$

If it is replaced \cdot by + and the map $x \mapsto 2x$ is used instead of $x \mapsto x^2$, then -by a similar argument of above proofs- an axiomatization similar to that of Corollary 4.14 can be obtained, because it has the following result:

Theorem 4.15. Let $M \models IE_0^1 + \neg IE_k$, and $\varphi(x) \in \Delta_0$ be an id(+)-gap. Then there exists $\psi \in \{\varphi, \overline{\varphi}\}$ such that for all $c > \varphi$ and $n \in \omega$, there exist m > n and $a \in (\varphi, c)$ such that

$$M \models \psi(a) \land \psi(2^m \cdot a)$$

Moreover, if $M \models IE_{k_e} + \mathbb{N} \cdot O(E_k)$, then there exists $\psi \in \{\varphi, \overline{\varphi}\}$ such that for all $c > \psi$ there exists $a \in (\varphi, c)$ and δ nonstandard such that

$$M \models \psi(a) \land \psi(2^{\delta} \cdot a)$$

5. Closing remarks

The starting thesis of the article was to advance in the study of models of very weak arithmetic theories by considering different methods from the classics that need more resources (those based on recursiontheoretic methods, definability of satisfaction, etc.). On the one hand, it has been proven that we can reduce some problems to the existence of a certain model or of some Δ_0 -definable set. Axiom schemes describing properties of algebro-combinatoric nature for weak induction have also been provided, thus complementing related results for other theories without Σ_1 induction (see e.g. [18,52]).

In the paper it has been shown that the use of ultrafilters of the set $\mathbb{S}_{\varphi} = \pi^{-1}[\{[\varphi]\}]$ seems to be a good substitute to the use of types in recursive saturation: it proves the existence of an element (e.g. a φ -composite element) without realizing Δ_0 -types. However, is not used the full ultrafilter, actually we work with $p \cap \nabla_k(M)$ instead of $p \in \beta M$ (where $\nabla_k(M)$ is the class of both E_k and U_k definable sets in M), which may be seen as a ∇_k -type.

Nonetheless, it should be pointed out that the restriction to Δ_0 -definability limits the algebraic methods designed for $\beta \mathbb{N}$. A closer look reveals that the product has a complex behavior on the set of idempotents of \mathbb{H}_{φ} :

Proposition 5.1. If $M \models RIE_0^1$ is countable and $\psi < \varphi$ are gaps, then the left translation defined by an idempotent $q \in \mathbb{S}_{\psi}$ is not trivial on \mathbb{S}_{φ} (although we know that $[\psi] \cdot [\varphi] = [\varphi]$ for gaps): there is $p \in \mathbb{S}_{\varphi}$ idempotent such that $q \cdot p \neq p$ for all idempotent $q \in \bigcup_{\psi < \varphi} \mathbb{S}_{\psi}$.

Proof. Let us consider the semigroup $M_1 = M \setminus \{0\}$. Since \mathbb{S}_{φ} is a G_{δ} -set in βM_1 that contains idempotents, we have that there exists $p \in \mathbb{S}_{\varphi}$ such that for every idempotent $q \in \beta M_1 \setminus M_1$,

$$q \cdot p = p \implies p \cdot q = q$$

([24], th. 9.7). If $q \in \bigcup_{\psi < \varphi} \mathbb{S}_{\psi}$, then $\pi(p \cdot q) = \pi(p) \cdot \pi(q) = [\varphi] \cdot [\psi] = [\varphi]$ (by 3.25). Thus $q \cdot p \neq p$. \Box

Result 5.1 is also true for \mathbb{H}_{φ} , and it suggests the hardness of defining a semigroup embedding

$$i: M[\mathcal{G}_k] \hookrightarrow \beta M$$

such that $i \upharpoonright_M = Id_M$ with interesting features. Results like above illustrate that the algebraic combinatorics of βM have interest beyond the topic of the paper (Weak Arithmetics). For example, the availability of such an embedding can be useful for the end-extension problem for $I\Delta_0$. Indeed we hope to follow this research line.

Nevertheless, we also remark that it is possible to find an algebraic and topological representation of βM (being M countable) in terms of $\beta \mathbb{N}$: for each $* \in \{+, \cdot\}$ there exists a surjective continuous homomorphism

$$\Phi_*: (\bigcap_{n>0} \overline{2^n \mathbb{N}}^{\beta \mathbb{N}}, +) \to (\beta M, *)$$

(cf. th. 6.4 of [24]). Therefore, if \sim_{Φ_*} is the congruence induced by Φ_* , then

$$\bigcap_{n>0} \overline{2^n \mathbb{N}}^{\beta \mathbb{N}} / \sim_{\Phi_*} \cong \beta M$$

Similar results in that line were already available (see for example [34,35]) from seminal works by Skolem and Tennenbaum, in which they relate the nature of certain ultrafilters used to build the quotient (in our case, the construction would be by means of congruences instead). The fact motivates as future work to study the congruences in βM that produce models of a certain arithmetic theory, particularly: **Open problem:** To characterize the congruences on

$$\bigcap_{n>0} \overline{2^n \mathbb{N}}^{\beta \mathbb{N}}$$

that induce models of $I\Delta_0$.

Results on that line already exist. For example, in [53] Bohr compactification of a group is characterized by a congruence that makes a certain class of ultrafilters equivalent to the identity element. This result applies to the group of integers \mathbb{Z}_M induced by $M \models PA^-$.

Finally, an interesting question arises by considering the idea behind Corollary 4.12. The combinatorial apparatus built to prove Theorem 4.10 could be applied to other functions with similar behavior, for instance to those defined in the given theory by $\exists E_k$ -formulas without parameters such that $M \models x < f(x)$. It seems that 4.10 could be generalized when the formula

$$f(\varphi)(x) := \exists u < x(f(u) = x \land \varphi(u))$$

defines the same gap as φ . From there one could try to obtain results like 4.12, using some Δ_0 -definition of the iteration of f (if this is possible), mainly for gaps φ such that $f(\varphi) \sim \varphi$. There are reasons to expect that, in many cases, one might obtain axiom schemes between IE_n and IE_{n+1}, suggesting the future study of the class of recursive functions with *axiomatizable combinatorics* (like Corollaries 4.6 or 4.14) in a theory weaker than I Δ_0 .

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