

Towards a verifiable topology of data

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Schedule

- Mathematical background.
- Persistent homology over
 - a field.
 - an elementary divisor ring.
- Formalisation in ACL2.

Mathematical background

Simplicial Complex

V ordered set

Simplicial Complex (over V)

- $K \subseteq P_{\text{fin}}(V)$, K closed under the subset relation
- $S_n(K)$, “simplices whose cardinality is $n+1$ ”
- $\sigma_n^i : S_n(K) \rightarrow S_{n-1}(K)$,
$$\sigma_n^i(v_0, \dots, v_n) = (v_0, \dots, v_{i-1}, v_{i+1}, \dots, v_n)$$

Mathematical background
Chain Complex

R ring

Chain Complex (over R)

$$\mathbf{C}_* = \{(C_n, d_n) / n \in \mathbb{Z}\}$$

- C_n , R-module
- $d_n: C_n \rightarrow C_{n-1}$, morphism of R-modules, $\forall n$
- $d_n d_{n+1} = 0$, $\forall n$

Mathematical background

Chain Complex associated with a simplicial complex K

K simplicial complex.

$C_*(K)$ is given by:

- $C_n(K)$, free R -module over $S_n(K)$
- $d_n = \sum_{i=0}^n (-1)^i \sigma_i^n$

Mathematical background
Homology

K simplicial complex.

The **n -th homology group of K**

$$H_n(K) = \frac{Z_n}{B_n}$$

($d^2 = 0$ implies $\text{Im } d_{n+1} = B_n \leq Z_n = \text{Ker } d_n$)

Mathematical background

Filtration

- **Filtration** of a simplicial complex K

$$K^1 \subseteq K^2 \subseteq \dots \subseteq K^d = K$$

(K is a **filtered simplicial complex**)

- **Filtration** of chain complexes

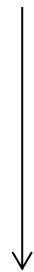
$$C_*^0 \xrightarrow{f_0} C_*^1 \xrightarrow{f_1} C_*^2 \rightarrow \dots$$

(f_i is a morphism of chain complexes)

Mathematical background
Induced filtration

A filtration of simplicial complexes

$$K^0 \rightarrow K^1 \rightarrow K^2 \dots \rightarrow K^d = K$$



induces

$$C_*(K^0) \rightarrow C_*(K^1) \rightarrow \dots \rightarrow C_*(K^d)$$

a filtration of chain complexes

(arrows are inclusion maps)

Mathematical background
Persistent homology

A filtration of simplicial complexes

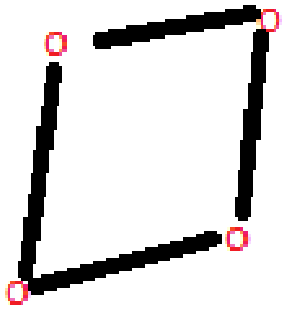
$$K^0 \rightarrow K^1 \rightarrow K^2 \dots \rightarrow K^d = K$$

The **p-persistent n-dimensional homology group of K^j**
is given by

$$H_n^{j,p} = Z_n^j / B_n^p \cap Z_n^j \quad (j \leq p)$$

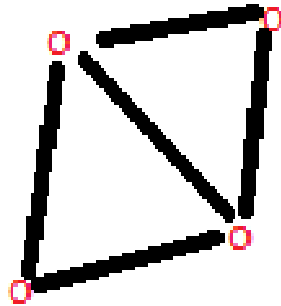
$$(Z_n^j \leq Z_n^p \leq C_n(K^p) \text{ and } B_n^p \leq C_n(K^p))$$

An example



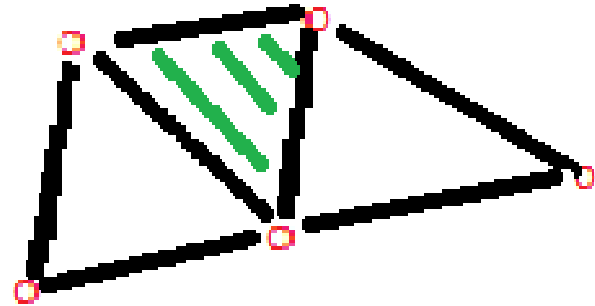
K^3

$$H_1^{4,4} = \mathbb{Z} + \mathbb{Z}$$



K^4

$$H_1^{4,5} = \mathbb{Z}$$



K^5

$$H_1^{5,5} = \mathbb{Z} + \mathbb{Z}$$

Computing persistent homology (over a field F)

Mathematical basis

- Artin-Rees Theorem implies
“ F -modules of persistence $\equiv F[x]$ -modules”
(finite type)
- F is a field, $F[x]$ is a PID
- To computing the persistent homology of a filtration of F -modules is equivalent to computing the homology of its (Artin-Rees) associated $F[x]$ -module.

Computing persistent homology (over a field F)

Algorithm

Algorithm

Variant of the gaussian elimination algorithm

- Polynomial time
- Persistence intervals
- Barcodes

Computing persistent homology (over an ED-ring R)

Mathematical basis

Two essential bricks:

- Echelon forms for matrices
(effective Bezout domains)
- Smith normal form
(elementary divisor rings)

Computing persistent homology (over an ED-ring R)

Echelon form

An echelon form

$(A \in M(n \times m, R))$

$n=1 \longrightarrow A \longrightarrow (0 \ 0 \ 0 \dots a), \ a \in R$

$n>1 \longrightarrow A \longrightarrow \begin{array}{|c|c|} \hline A_1 & \\ \hline 0 & \\ \hline \end{array}$

\searrow

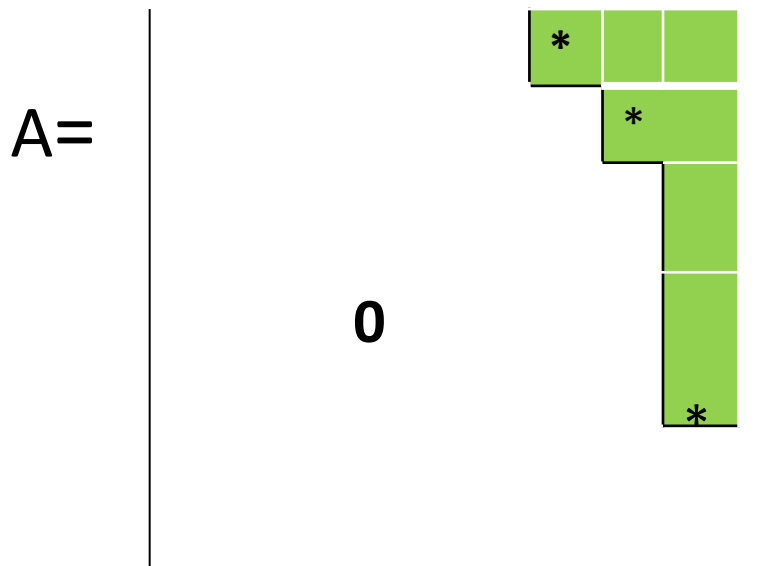
$\begin{array}{|c|c|} \hline A_1 & * \\ \hline 0 & a \\ \hline \end{array}$

$(a \neq 0)$ and $\text{echelonForm}(A_1)$

Computing persistent homology (over an ED-ring R)

EchelonForm

echelonForm(A)



(the width of the steps is 1)

Computing persistent homology (over an ED-ring R)

Existence of echelon form

R is an effective Bezout domain.

Th1: Existence of echelon form ($A \in M(n \times m, R)$)

“There exists P invertible s.t. $\text{echelonForm}(AP)$ ”

(P is a sequence of Bezout and permutation “elementary operations”)

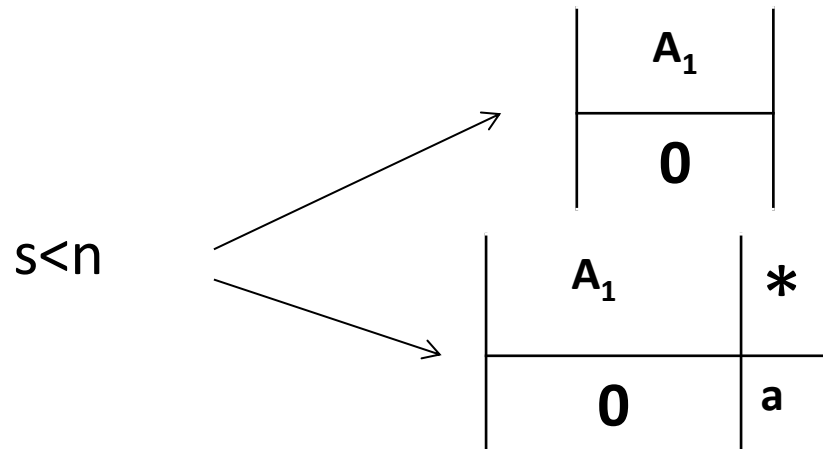
Computing persistent homology (over an ED-ring R)

Generalized echelon form

echelonForm(A,s)

$A \in M(n \times m, R)$

$s \geq n \longrightarrow \text{true}$

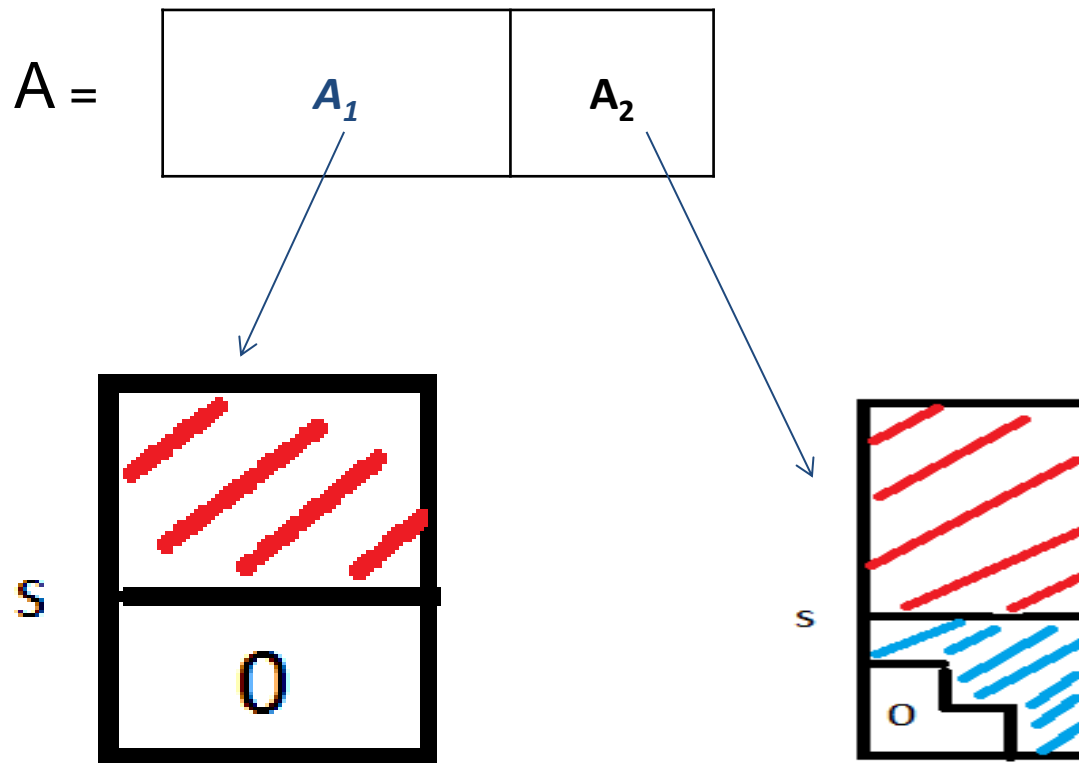


$(a \neq 0)$ and $\text{echelonForm}(A_1, s)$

Computing persistent homology (over an ED-ring R)

Generalized echelon form

echelonForm (A, s)

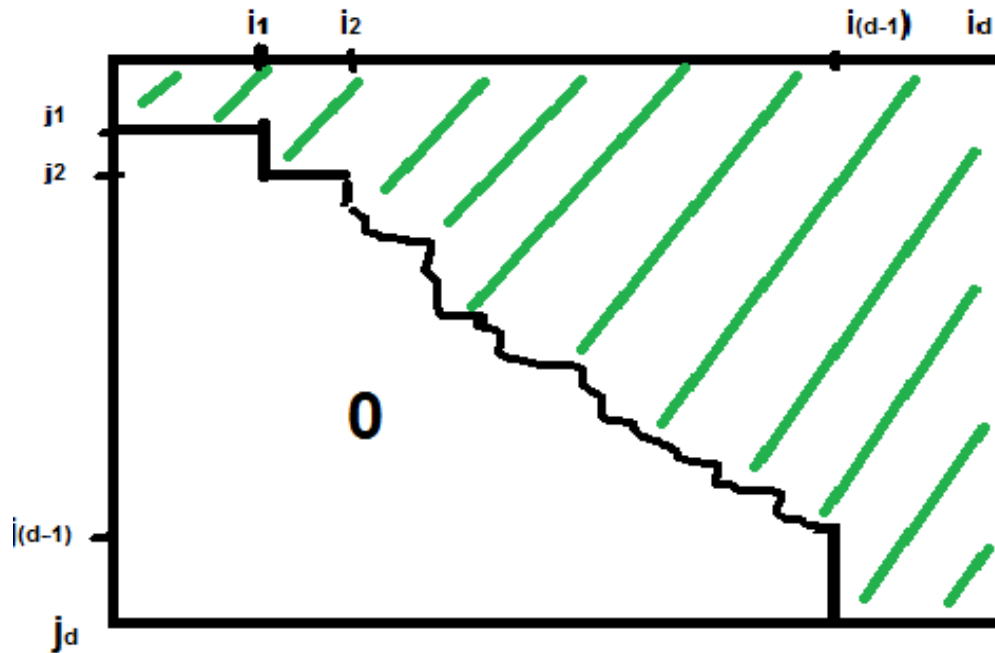


Computing persistent homology (over an ED-ring R)

Graded matrices

Graded matrix

$A \in M(n \times m, R)$



$$0 \leq j_1 \leq j_2 \leq \dots \leq j_d = n$$
$$0 \leq i_1 \leq i_2 \leq \dots \leq i_d = m$$

Computing persistent homology (over an ED-ring R)

Graded matrices

A filtration of simplicial complexes

$$K^0 \rightarrow K^1 \rightarrow K^2 \dots \rightarrow K^d = K$$

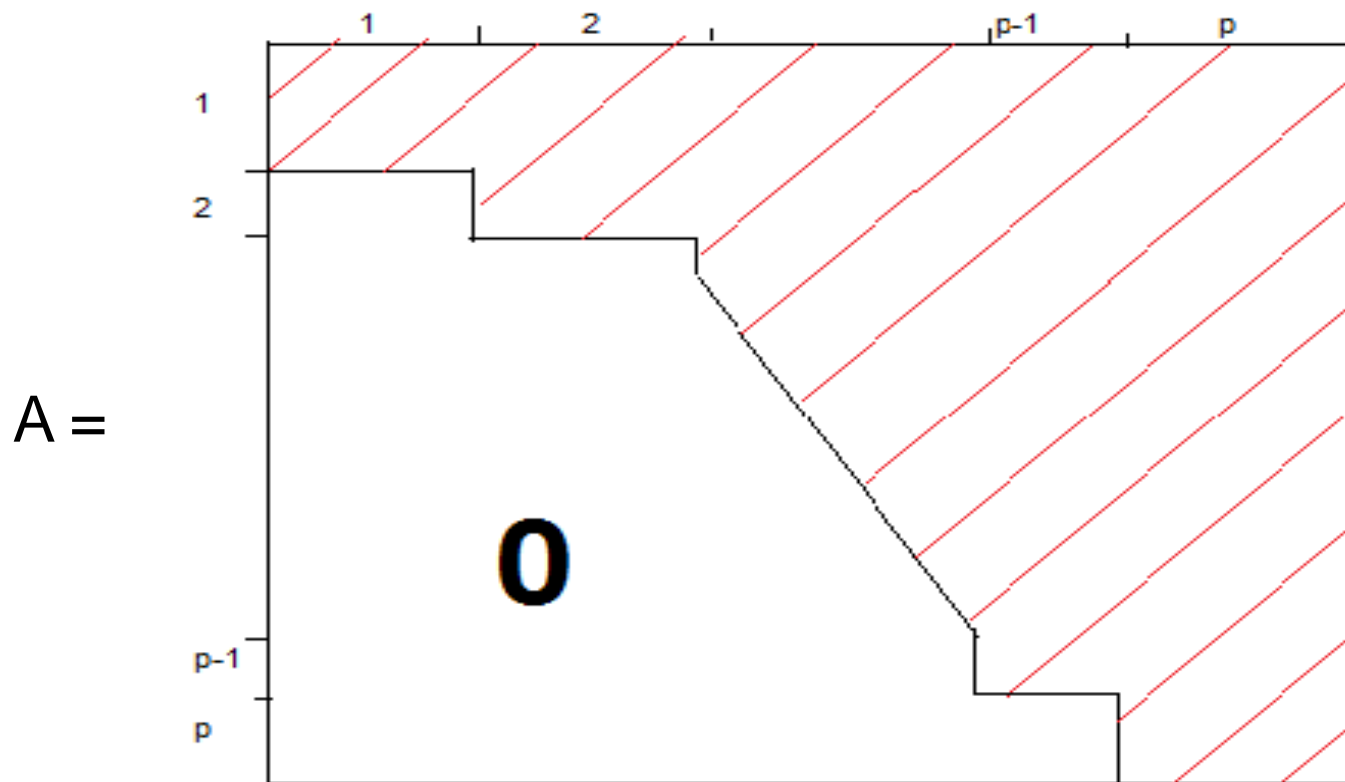
The **standard matrix representation** of the differential map

$$d_n = \{d_n^i : C_n(K^i) \rightarrow C_{n-1}(K^i)\}$$

is a **graded matrix**.

Computing persistent homology (over an ED-ring R)

Graded matrices



$i=1,2,\dots,p$ represent the filtration index

Computing persistent homology (over an ED-ring R)

Graded echelon form

R is an effective Bezout domain.

Th2: “Existence of graded echelon form”

Input: (A, s, t)

- $A \in M(n \times m, R)$, graded matrix
- $0 \leq s \leq t \leq d$

Computing persistent homology (over an ED-ring R)

Graded echelon form

R is an effective Bezout domain.

Th2: “Existence of graded echelon form”

Input: (A, s, t)

Output: $(P_s, P_{s+1}, \dots, P_t)$

- P_i are invertible matrices
- $\text{echelonForm}(E_i, s)$, where $E_i = A_i P_i$, and...

Computing persistent homology (over an ED-ring R)

Graded echelon form

$$E_i = \begin{array}{|c|c|} \hline E_i^1 & E_i^2 \\ \hline \end{array}$$

$$E_i = A_i P_i$$
$$E_i^1 \in M(_x b_i, R)$$

$$P_i = \begin{array}{|c|c|} \hline P_{i-1}^{1..b(i-1)} & * \\ \hline 0 & \\ \hline \end{array}$$

Computing persistent homology (over an ED-ring R)

Smith normal form

R is an elementary divisor ring.

Th3: “Existence of Smith normal form”

Input: A ($A \in M(n \times m, R)$)

Output: $((d_1, \dots, d_r), P, Q)$

- P, Q invertible matrices
- $d_i \mid d_{i+1}$ ($i=1..r-1$)
- $PAQ = \text{DiagonalMatrix}(d_1, \dots, d_r, 0, \dots, 0)$

Computing persistent homology (over an ED-ring R).

Algorithm

Algorithm

Input:

- $K^0 \rightarrow K^1 \rightarrow K^2 \dots \rightarrow K^d = K$
(a filtered complex K)
- $n \geq 0$
- r, t such that $0 \leq r \leq t \leq d$

Output:

- $H_n^{r,j}(K) \quad r \leq j \leq t$

Computing persistent homology (over an ED-ring R).

Algorithm (Step 1)

Let A be the standard matrix representation of

$$d_n = \{d_n^i: C_n(K^i) \rightarrow C_{n-1}(K^i) \mid i=1..d\}$$

Step 1: `gradedEchelonForm (A, 0, d)`



invertible matrices (P_1, \dots, P_d)

Computing persistent homology (over an ED-ring R).

Algorithm (Step 1)

invertible matrices (P_1, \dots, P_d) such that

$$A P_d = \begin{array}{|c|c|} \hline & \begin{array}{c} \text{id} \quad m \\ \hline \end{array} \\ \hline \begin{array}{c} 0 \\ \hline \end{array} & \begin{array}{c} \text{red staircase} \\ \hline 0 \end{array} \\ \hline \end{array}$$

Columns of P 's matrices provide us a family of compatible basis of the kernels

$$\begin{aligned} B_1 &\leq B_2 \leq \dots \leq B_d \\ Z_1 &\leq Z_2 \leq \dots \leq Z_d \end{aligned}$$

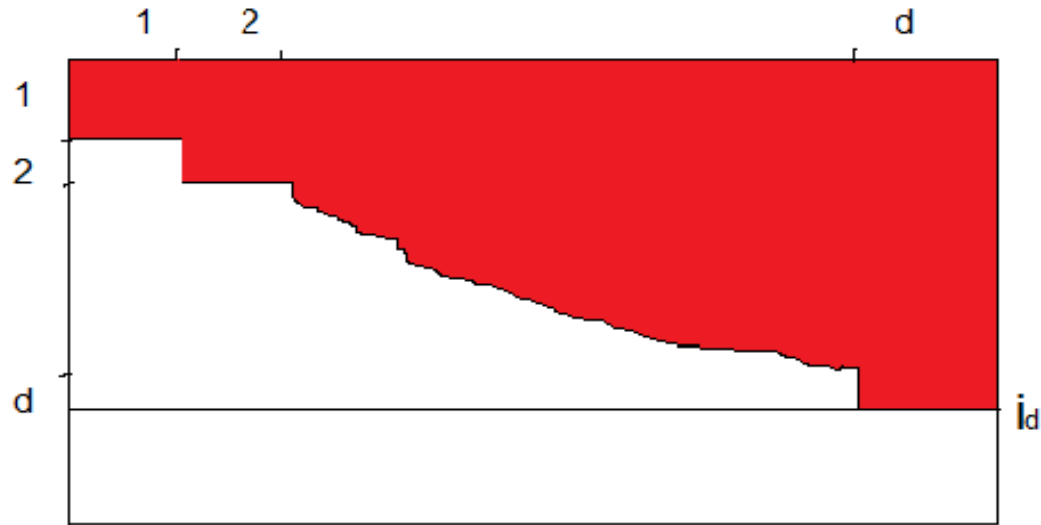
Computing persistent homology (over an ED-ring R).

Algorithm (Step 2)

Let B be the standard matrix representation of d_{n+1}

$$d_{n+1} = \{d_{n+1}^i: C_{n+1}(K^i) \rightarrow C_n(K^i) \mid i=1..d\}$$

Step 2: $P_d^{-1} B =$

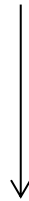


$(\text{Im } d_{n+1}^i \leq \text{Ker } d_n^i = Z_i \Rightarrow \text{last rows of } P_d^{-1} B \text{ are zero})$

Computing persistent homology (over an ED-ring R).

Algorithm (Step 3)

Step 3: $\text{gradedEchelonForm}(P_d^{-1} B, r, t)$



(Q_{r+1}, \dots, Q_t)

$\text{echelonForm}((P_d^{-1} B)_j Q_j, r) \quad (r \leq j \leq t)$

Computing persistent homology (over an ED-ring R).

Algorithm (Step 4)

$$\text{Let } E_j = (P_d^{-1} B)_j Q_j = [E_j^1 \mid E_j^2]$$

Step 4: SmithForm (E_j^1)



$$H_n^{r,j}(K)$$

An Smith computation for each filtration index j !

Formalisation in ACL2

- **Our goal:** implement and formally verify this algorithm in the ACL2 Theorem Prover
- **This means** to implement and verify
 - Echelon form
 - Graded echelon form
 - Smith normal form

Formalisation in ACL2

- One of our main concerns is efficiency
Advantage of using ACL2: we verify Common Lisp code.
- But ACL2 is an applicative subset of CL. In principle, we have not CL arrays, only lists.
- Fortunately, we can use stobjs (single-threaded objects), which allows destructive updates and constant-time accesses, without losing the applicative semantic.

What we have now:

- An array-based version of echelon form algorithm (“efficient” but unverified).
- A list-based version of echelon form and Smith form algorithms (still unverified).
- Infrastructure proving operational equivalence between both approaches.

Our goal: reasoning using the second approach and executing using the first one.